



# Algorithm Theory

## 06 – Amortized Analysis

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- ~~It~~ Technique for the analysis of algorithms and data structures
- Works if a sequence of operations has to be analyzed
- How much does an operation cost "on average"?

# Amortization

- Consider a sequence  $a_1, a_2, \dots, a_n$  of  $n$  operations performed on a data structure  $D$
- $T_i =$  execution time of  $a_i$  *could be high.*
- $T = T_1 + T_2 + \dots + T_n$  total execution time
- The execution time of a single operation can vary within a large range, e.g. in  $1, \dots, n$ , but the worst case does not occur for all operations of the sequence.
- Average execution time of an operation is small, even though a single operation can have a high execution time.

$$\frac{1}{n} \cdot \sum_{i=1}^n T_i \quad \text{Goal: What is the average execution time an operation in a sequence of operations?}$$

# Analysis of algorithms

- **Best case**  $\leadsto$  often too optimistic  
Smallest execution times
- **Worst case**  $\leadsto$  often too pessimistic  
Largest execution times
- **Average case**  $\leadsto$  often realistic distributions unknown, not analyzable often  
Input is drawn from prob. distribution, e.g. all inputs are equally likely
- **Amortized worst case**

What is the average cost of an operation in a worst case sequence of operations?



# Amortization

## Idea:

- Pay more for inexpensive operations
- Use the credit to cover the cost of expensive operations

## Three methods:

1. Aggregate method
2. Accounting method
3. Potential method

# Amortization = Overcharging + Bookkeeping



Pay two ~~cost~~ cost-units per operation  
 Store unused cost-units in a bank account

