

# Theoretical Computer Science - Bridging Course

## Summer Term 2019

### Exercise Sheet 7

for getting feedback submit electronically by 06:00 am, Friday, June 21st, 2019

#### Exercise 1: Decidability (7 Points)

Let  $\Sigma$  be a fixed finite alphabet. Show that the language of deterministic finite automata (DFAs) on  $\Sigma$  that accept every word is decidable. Formally, show that

$$L = \{\langle A \rangle \mid A \text{ is a deterministic finite automaton with } L(A) = \Sigma^*\}$$

is a decidable language.

*Remark: You can use that it is not difficult to construct a Turing machine which tests whether an input is the well formed encoding of a deterministic finite automaton.*

#### Exercise 2: Landau Notation (2+2+3 Points)

The set  $\mathcal{O}(f)$  contains all functions that are asymptotically not growing faster than the function  $f$  (when additive or multiplicative constants are neglected). That is:

$$g \in \mathcal{O}(f) \iff \exists c \geq 0, \exists M \in \mathbb{N}, \forall n \geq M : g(n) \leq c \cdot f(n)$$

For the following pairs of functions, check whether  $f \in \mathcal{O}(g)$  or  $g \in \mathcal{O}(f)$  or both. Proof your claims (you do not have to prove a negative result  $\notin$ , though).

(a)  $f(n) = 100n$ ,  $g(n) = 0.1 \cdot n^2$

(b)  $f(n) = \sqrt[3]{n^2}$ ,  $g(n) = \sqrt{n}$

(c)  $f(n) = \log_2(2^n \cdot n^3)$ ,  $g(n) = 3n$

**Hint:** You may use that  $\log_2 n \leq n$  for all  $n \in \mathbb{N}$ .

#### Exercise 3: Sorting Functions by Asymptotic Growth (6 Points)

Sort the following functions by asymptotic growth using the  $\mathcal{O}$ -notation. Write  $g <_{\mathcal{O}} f$  if  $g \in \mathcal{O}(f)$  and  $f \notin \mathcal{O}(g)$ . Write  $g =_{\mathcal{O}} f$  if  $f \in \mathcal{O}(g)$  and  $g \in \mathcal{O}(f)$ .

$n^2$	$\sqrt{n}$	$2^n$	$\log(n^2)$
$3^n$	$n^{100}$	$\log(\sqrt{n})$	$(\log n)^2$
$\log n$	$10^{100}n$	$n!$	$n \log n$
$n \cdot 2^n$	$n^n$	$\sqrt{\log n}$	$n$