



## Algorithms and Data Structures Summer Term 2021 Exercise Sheet 5

### Exercise 1: Bad Hash Functions

Let  $m$  be the size of a hash table and  $M \gg m$  the largest possible key of the elements we want to store in the table. The following “hash functions” are poorly chosen. Explain for each function why it is not a suitable hash function.

- (a)  $h : x \mapsto \lfloor \frac{x}{m} \rfloor \bmod m$
- (b)  $h : x \mapsto (2x + 1) \bmod m$  ( $m$  even).
- (c)  $h : x \mapsto (x \bmod m) + \lfloor \frac{m}{x+1} \rfloor$
- (d) For each calculation of the hash value of  $x$  one chooses for  $h(x)$  a uniform random number from  $\{0, \dots, m-1\}$
- (e)  $h : x \mapsto \lfloor \frac{M}{x \cdot p \bmod M} \rfloor \bmod m$ , where  $p$  is prime and  $\frac{M}{2} < p < M$
- (f) For a set of “good” hash functions  $h_1, \dots, h_\ell$  with  $\ell \in \Theta(\log m)$ , we first compute  $h_1(x)$ , then  $h_2(h_1(x))$  etc. until  $h_\ell(h_{\ell-1}(\dots h_1(x)))$ . That is, the function is  $h : k \mapsto h_\ell(h_{\ell-1}(\dots h_1(x)))$

### Exercise 2: (No) Families of Universal Hash Functions

- (a) Let  $\mathcal{S} = \{0, \dots, M-1\}$  and  $\mathcal{H}_1 := \{h : x \mapsto a \cdot x^2 \bmod m \mid a \in \mathcal{S}\}$ . Show that  $\mathcal{H}_1$  is not  $c$ -universal for *constant*  $c \geq 1$  (that is  $c$  is fixed and must not depend on  $m$ ).
- (b) Let  $m$  be a *prime* and let  $k = \lfloor \log_m M \rfloor$ . We consider the keys  $x \in \mathcal{S}$  in base  $m$  presentation, i.e.,  $x = \sum_{i=0}^k x_i m^i$ . Consider the set of functions from the lecture (week 5, slide 15)

$$\mathcal{H}_2 := \left\{ h : x \mapsto \sum_{i=0}^k a_i x_i \bmod m \mid a_i \in \{0, \dots, m-1\} \right\}.$$

Show that  $\mathcal{H}_2$  is 1-universal.

*Hint: Two keys  $x \neq y$  have to differ at some digit  $x_j \neq y_j$  in their base  $m$  representation.*