Chapter 6
Graph Algorithms

Algorithm Theory
WS 2016/17

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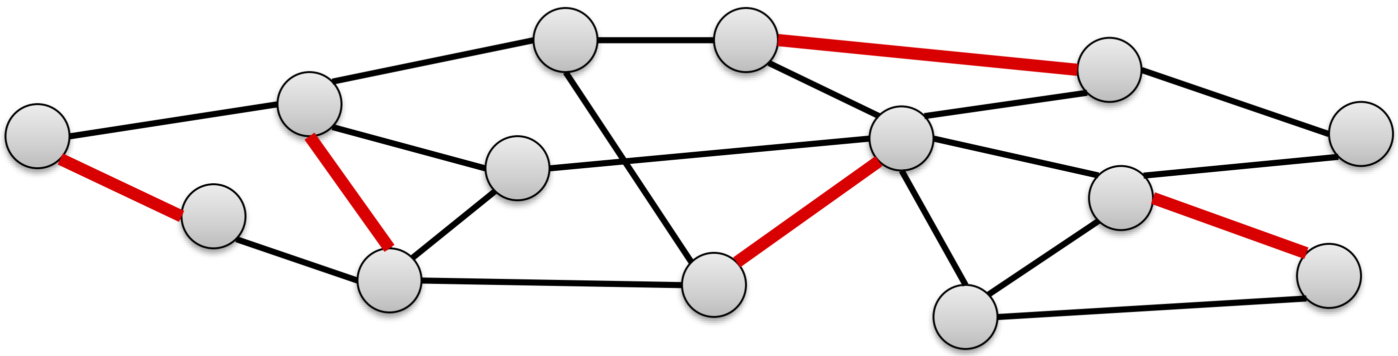
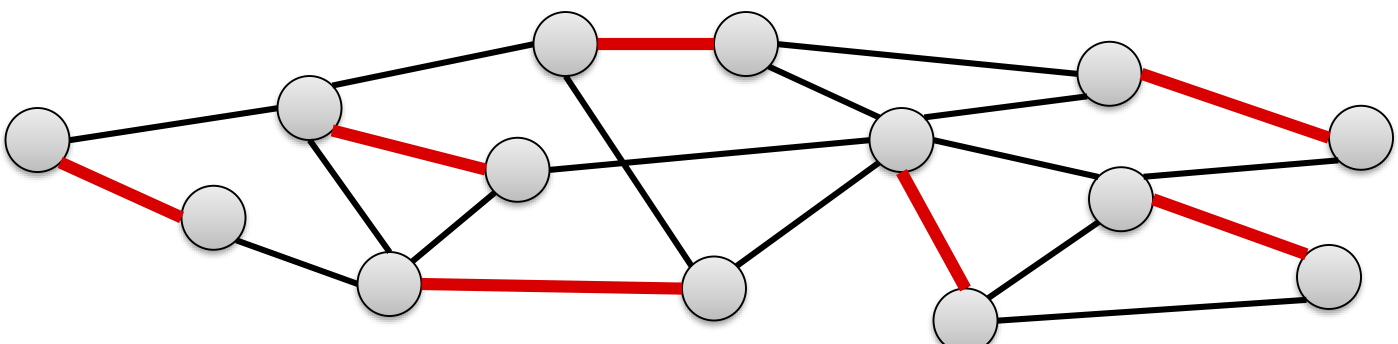
Matching

Matching: Set of pairwise non-incident edges



Maximal Matching: A matching s.t. no more edges can be added

Maximum Matching: A matching of maximum possible size



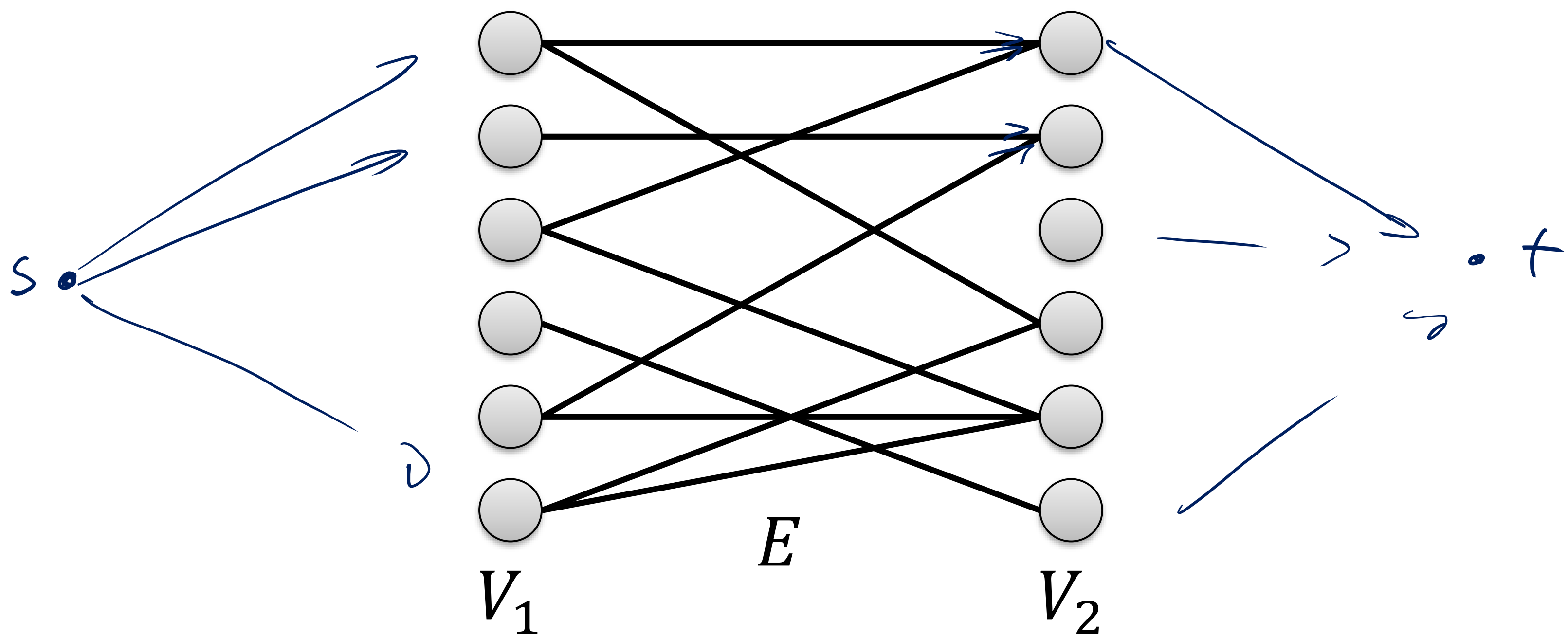
Perfect Matching: Matching of size $n/_2$ (every node is matched)

Bipartite Graph

Definition: A graph $G = (V, E)$ is called bipartite iff its node set can be partitioned into two parts $V = V_1 \dot{\cup} V_2$ such that for each edge $\{\mathrm{u}, \mathrm{v}\} \in E$,

$$|\{u, v\} \cap V_1| = 1.$$

- Thus, edges are only between the two parts



Hall's Marriage Theorem

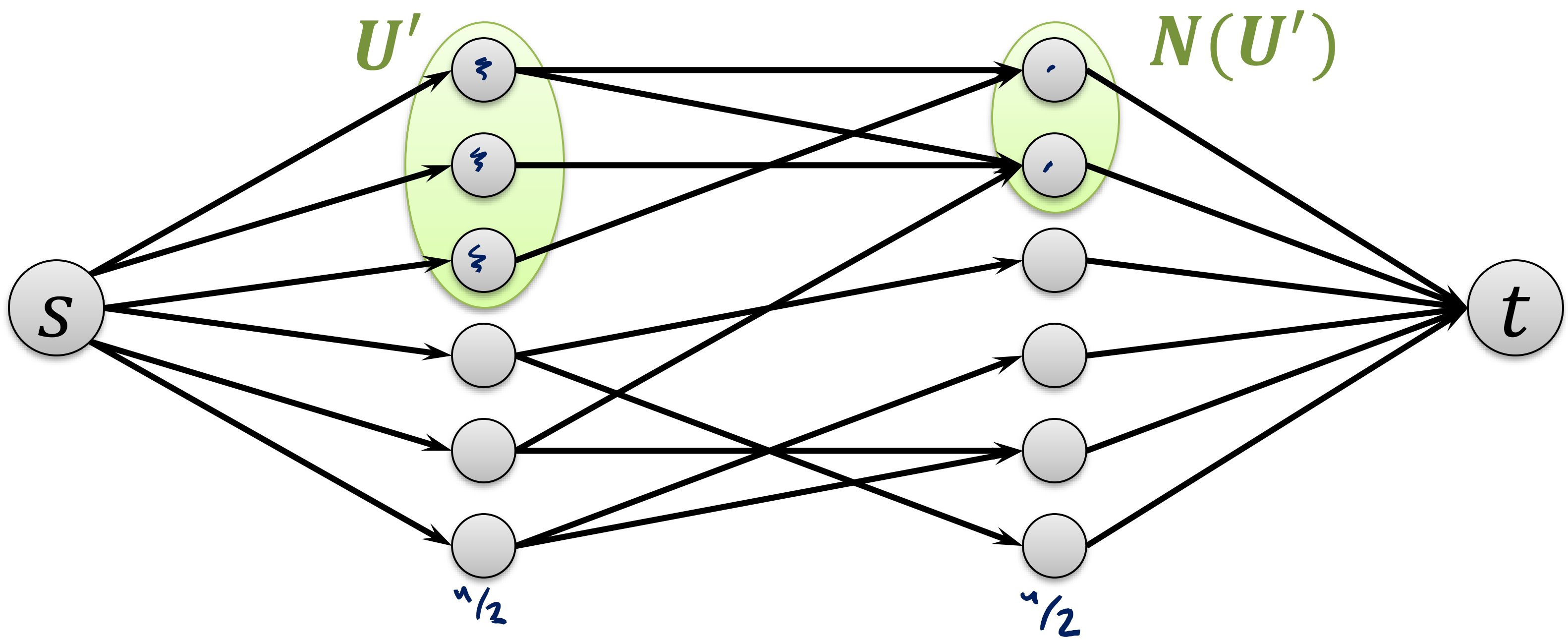
Theorem: A bipartite graph $G = (U \cup V, E)$ for which $|U| = |V|$ has a perfect matching if and only if

$$\forall \boldsymbol{U'} \subseteq \boldsymbol{U} : |\boldsymbol{N}(\boldsymbol{U'})| \geq |\boldsymbol{U'}|,$$

where $N(U') \subseteq V$ is the set of neighbors of nodes in U'.

Proof: No perfect matching $\Leftrightarrow$ some s-t cut has capacity $< n/2$

1. Assume there is U' for which $|N(U')| < |U'|$:

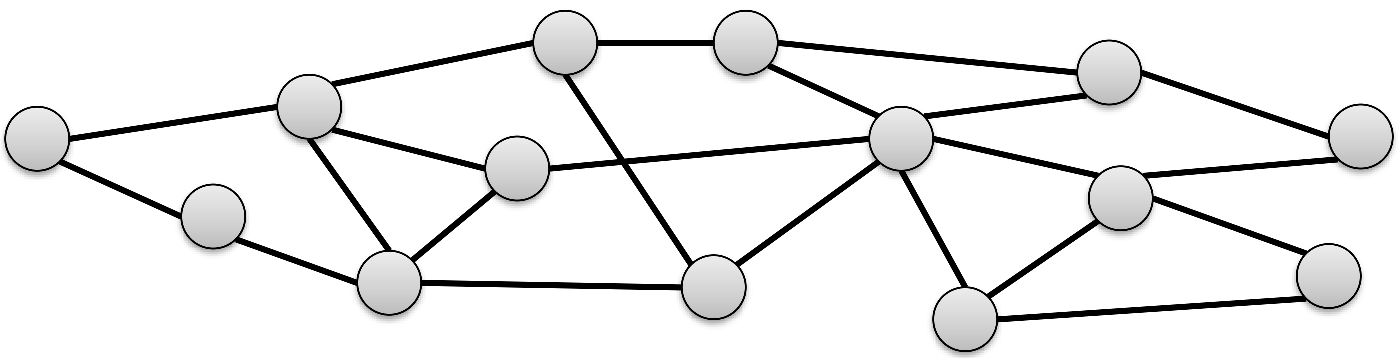


What About General Graphs

- Can we efficiently compute a maximum matching if G is not bipartite?

- How good is a maximal matching?
  - A matching that cannot be extended…

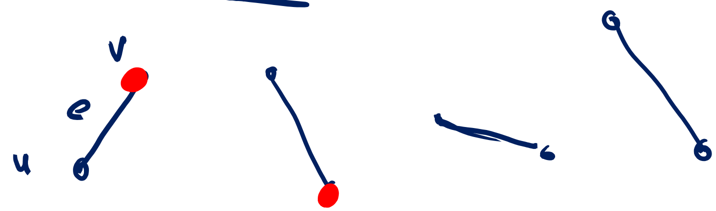
- **Vertex Cover:** set $S \subseteq V$ of nodes such that
$$\forall \{u, v\} \in E, \qquad \{u, v\} \cap S \neq \emptyset.$$



- A vertex cover covers all edges by incident nodes

Vertex Cover vs Matching

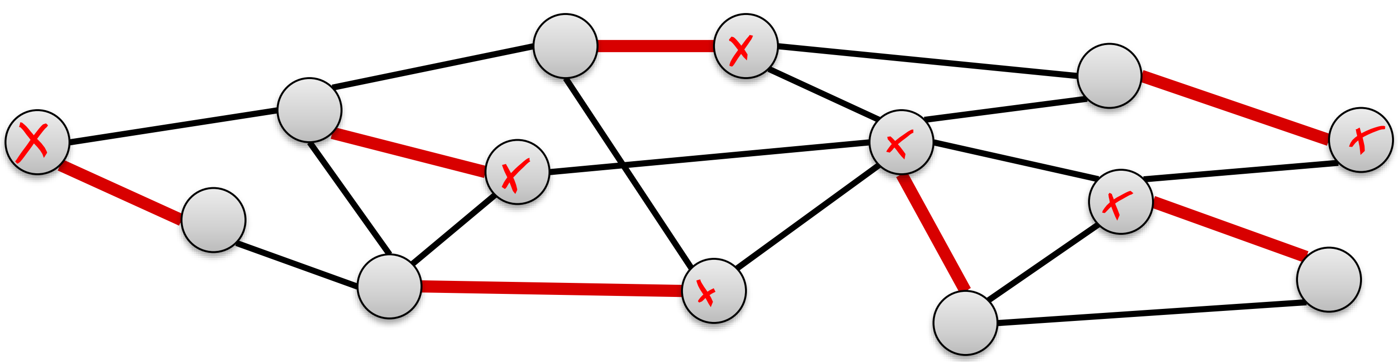
Consider a matching M and a vertex cover S



Claim: $|M| \leq |S|$

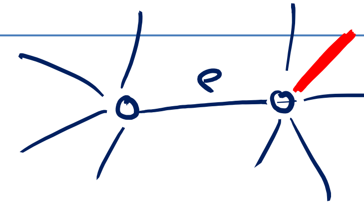
Proof:

- At least one node of every edge $\{u, v\} \in M$ is in S
- Needs to be a different node for different edges from M



Vertex Cover vs Matching

Consider a matching M and a vertex cover S

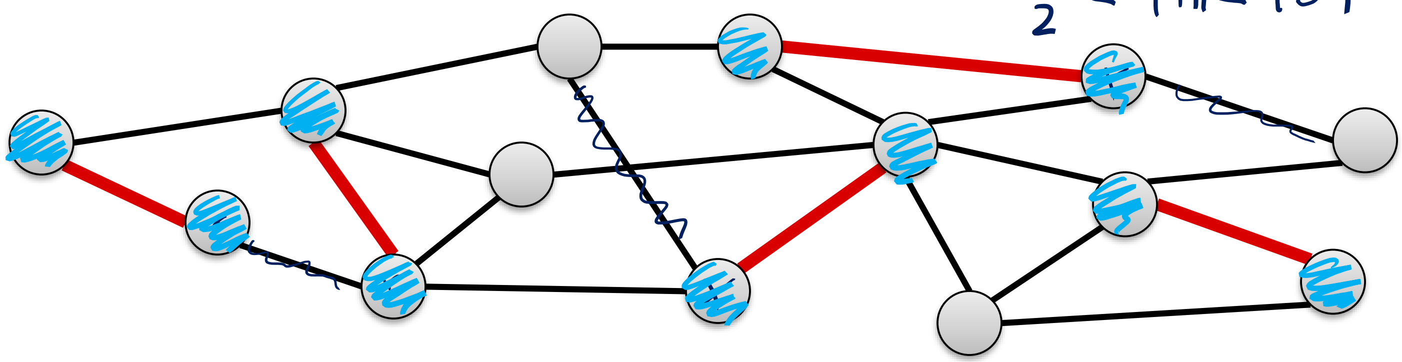


Claim: If M is maximal and S is minimum, $|S| \leq 2|M|$

Proof:

- M is maximal: for every edge $\{u, v\} \in E$, either u or v (or both) are matched

$\frac{|S|}{2} \leq |M| \leq |S^*|$



- Every edge $e \in E$ is "covered" by at least one matching edge
- Thus, the set of the nodes of all matching edges gives a vertex cover S of size $|S| = 2|M|$.

Maximal Matching Approximation

Theorem: For any maximal matching M and any maximum matching M^*, it holds that

$$|M| \geq \frac{|M^*|}{2}.$$

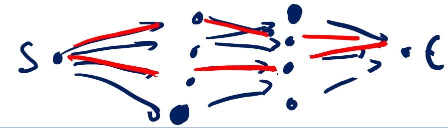
Proof:

S^*: opt. vertex cover

$|M^*| \leq |S^*| \leq 2|M|$

Theorem: The set of all matched nodes of a maximal matching M is a vertex cover of size at most twice the size of a min. vertex cover.

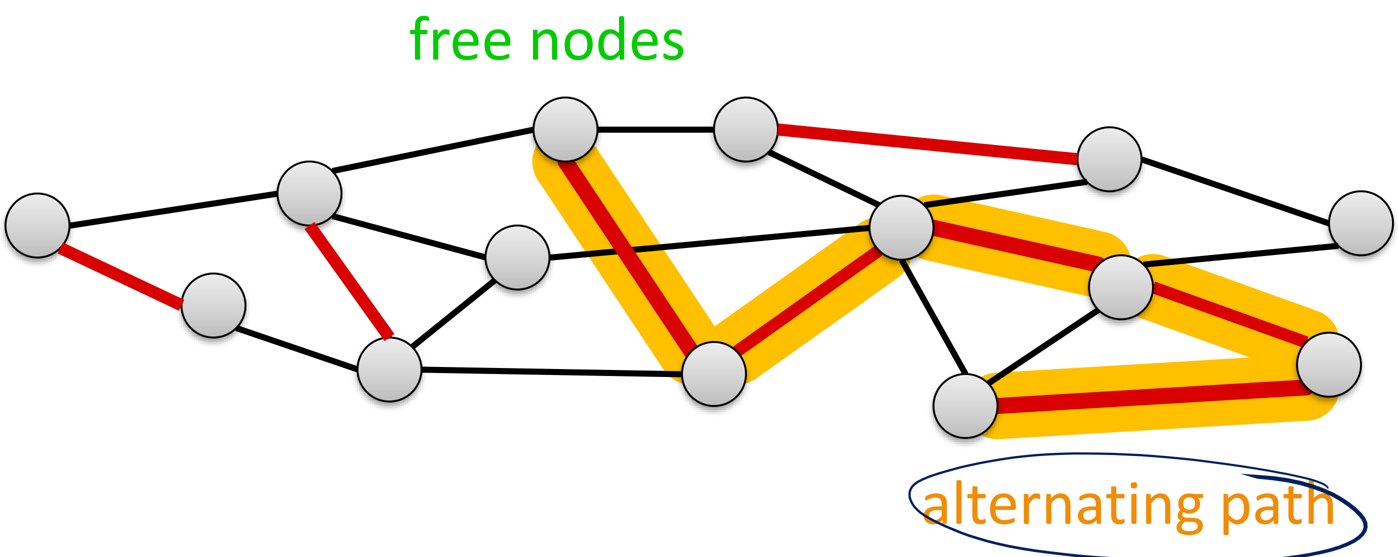
Augmenting Paths



Consider a matching M of a graph $G = (V, E)$:

- A node $v \in V$ is called **free** iff it is not matched

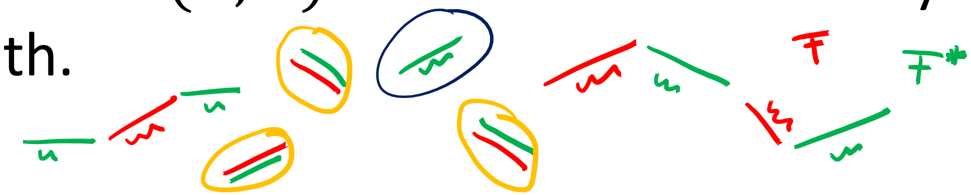
Augmenting Path: A (odd-length) path that starts and ends at a free node and visits edges in $E \setminus M$ and edges in M alternatingly.



- Matching M can be improved using an augmenting path by switching the role of each edge along the path

Augmenting Paths

Theorem: A matching M of $G = (V, E)$ is maximum if and only if there is no augmenting path.

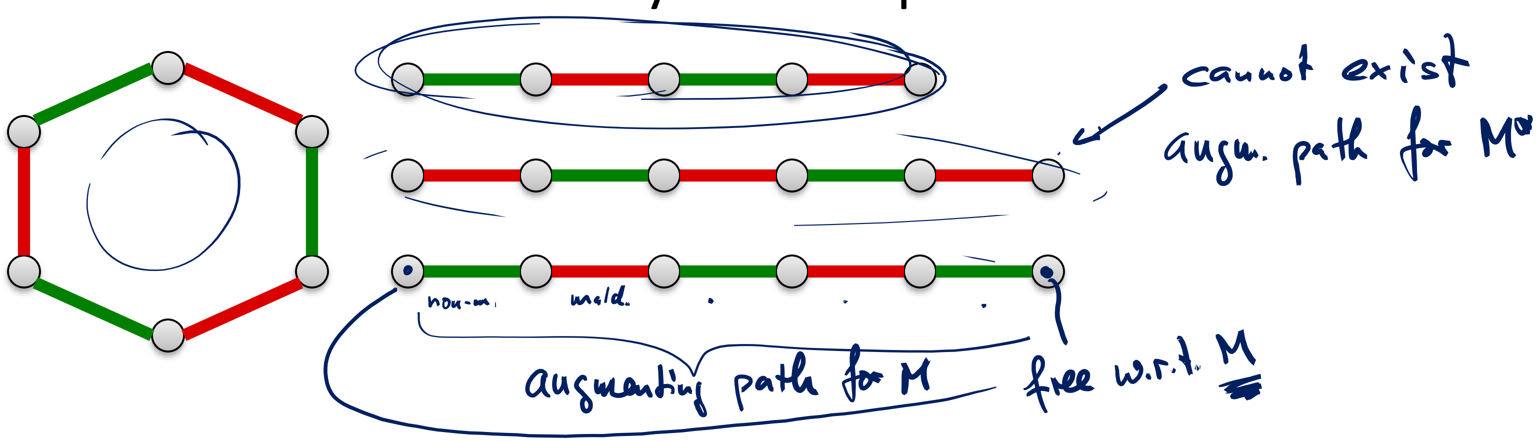


Proof:

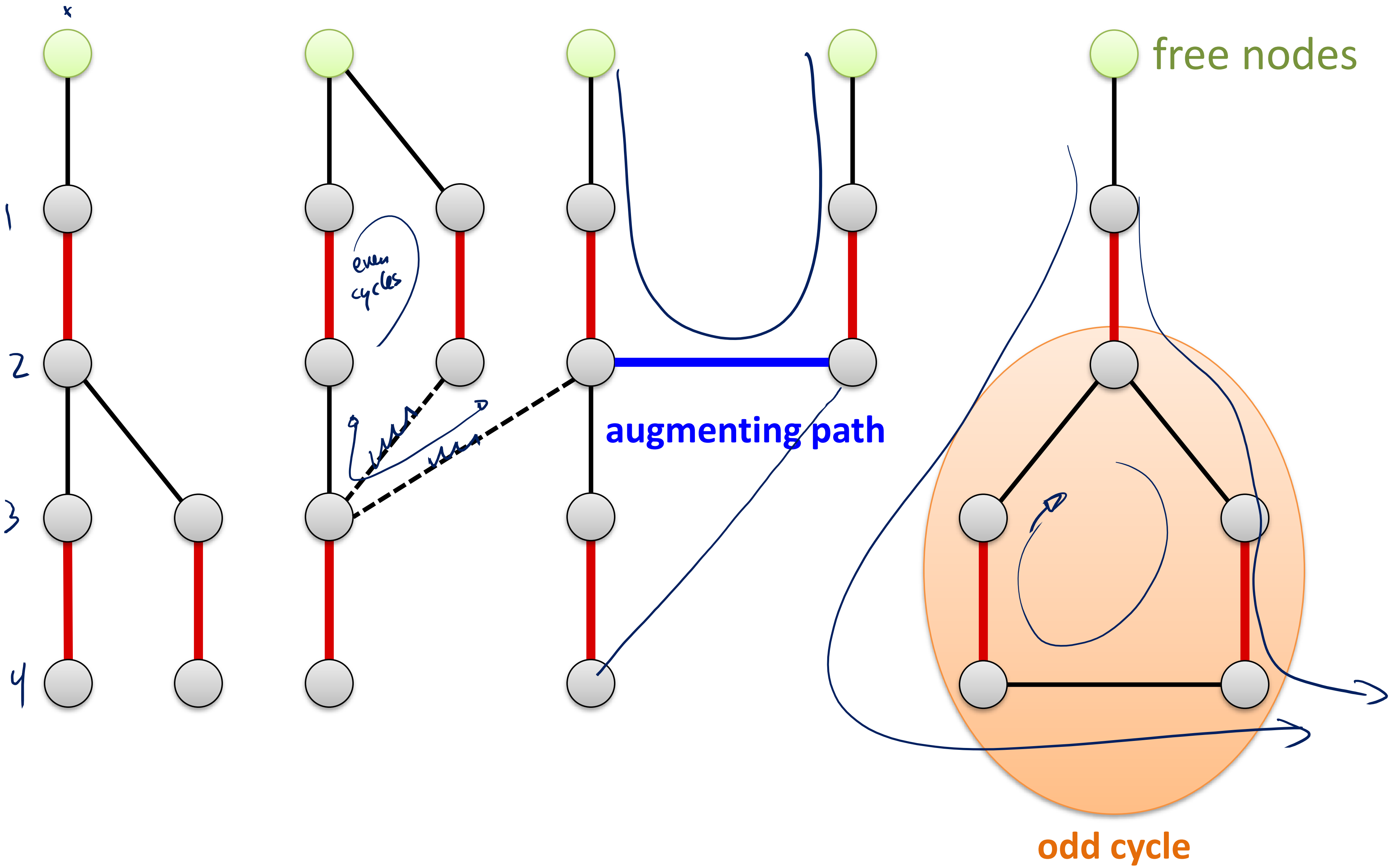
- Consider non-max. matching M and max. matching M^* and define

$$F := M \setminus M^*, \qquad F^* := M^* \setminus M$$

- Note that $F \cap F^* = \emptyset$ and $|F| < |F^*|$
- Each node $v \in V$ is incident to at most one edge in both F and F^*
- $F \cup F^*$ induces even cycles and paths

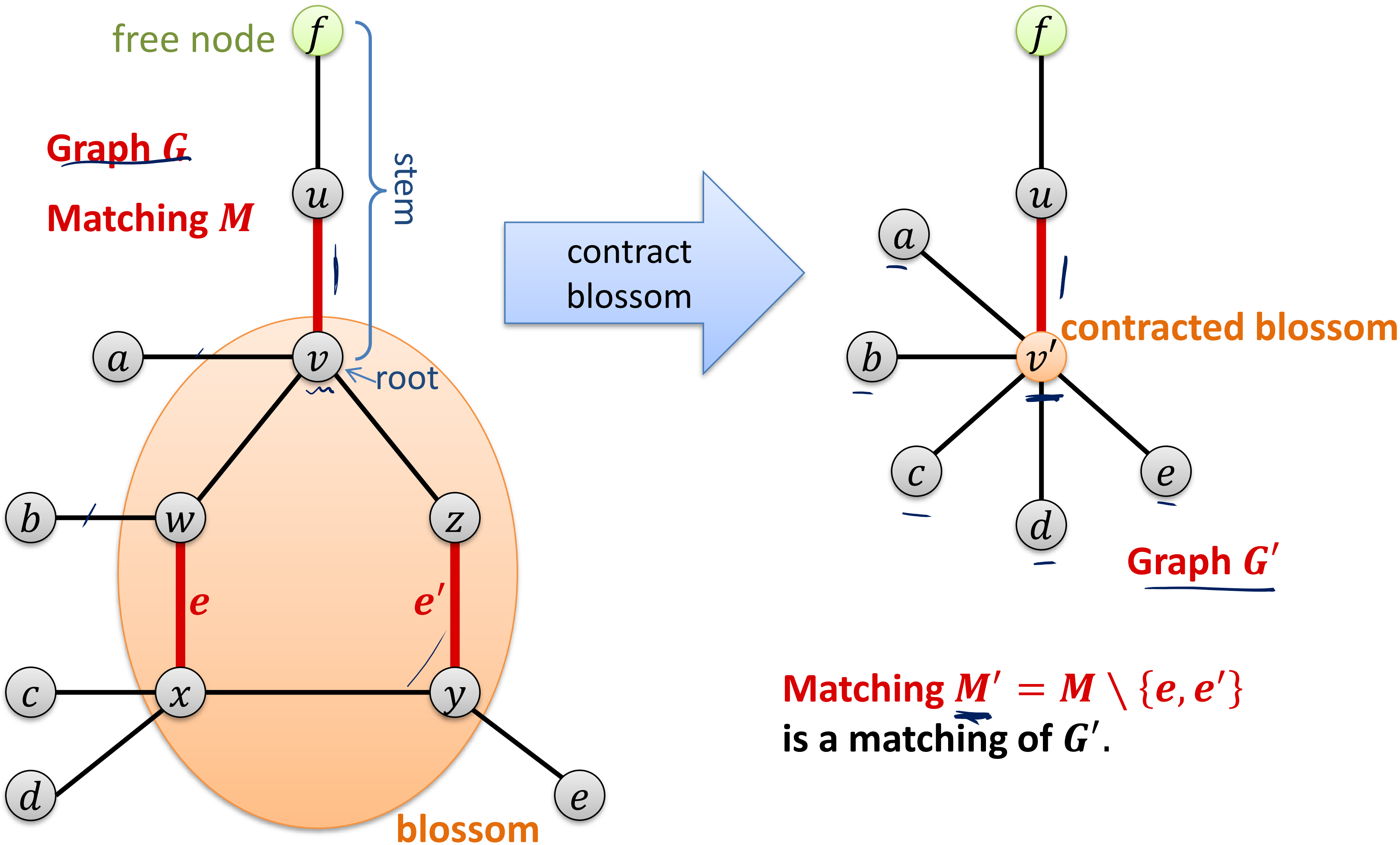


Finding Augmenting Paths



Blossoms

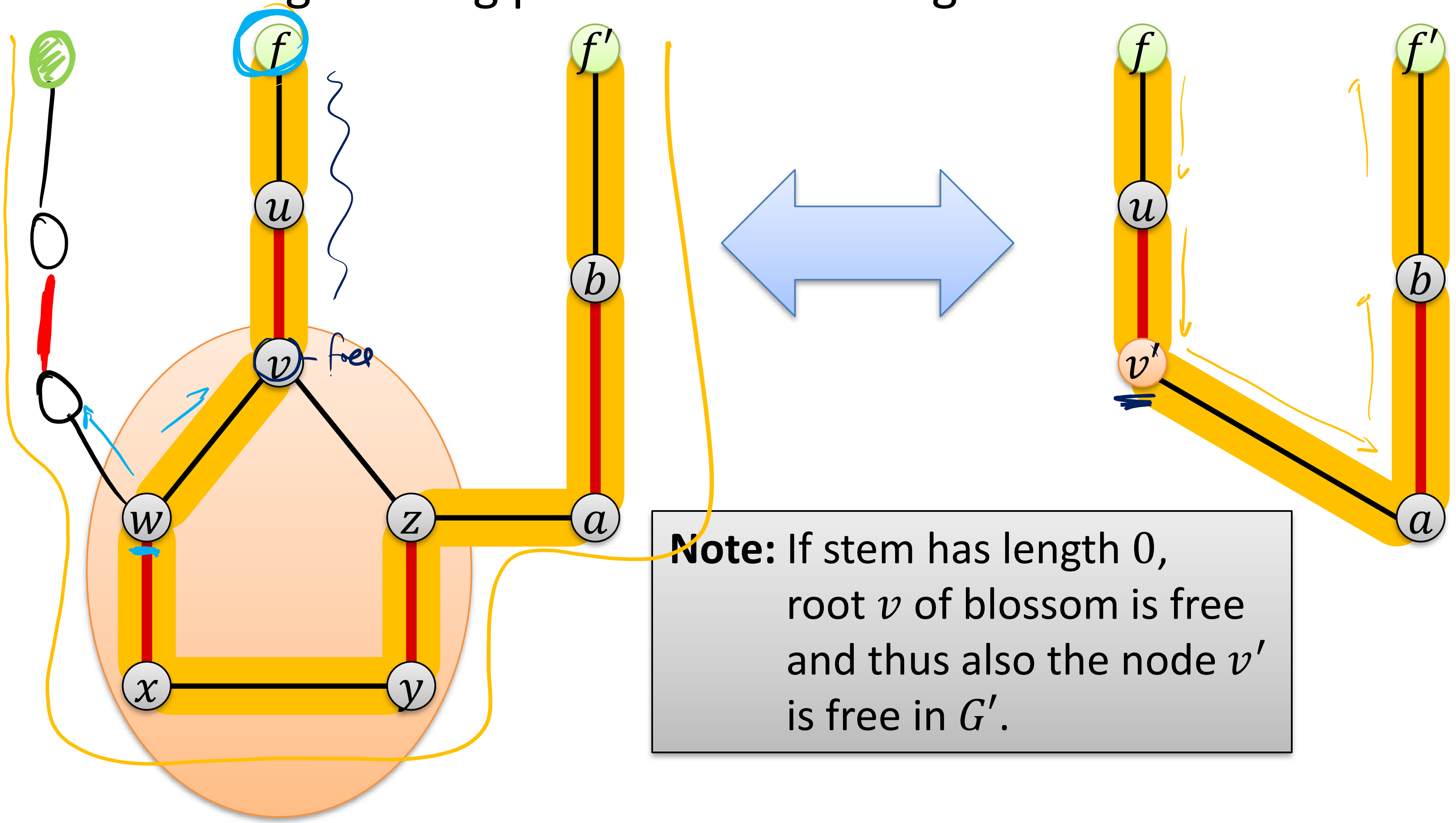
- If we find an odd cycle…



Matching $M' = M \setminus \{e, e'\}$
is a matching of G'.

Contracting Blossoms

Lemma: Graph G has an augmenting path w.r.t. matching M iff G' has an augmenting path w.r.t. matching M'



Note: If stem has length 0, root v of blossom is free and thus also the node v' is free in G'.

Also: The matching M can be computed efficiently from M'.

Edmond's Blossom Algorithm

Algorithm Sketch:

1. Build a tree for each free node
2. Starting from an explored node u at even distance from a free node f in the tree of f, explore some unexplored edge $\{u, v\}$:

    1. If v is an unexplored node, v is matched to some neighbor w:
    add w to the tree (w is now explored)

    2. If v is explored and in the same tree:
    at odd distance from root → ignore and move on
    at even distance from root → **blossom found** → smaller graph

    3. If v is explored and in another tree
    at odd distance from root → ignore and move on
    at even distance from root → **augmenting path found**

Running Time

Finding a Blossom: Repeat on smaller graph

Finding an Augmenting Path: Improve matching

Theorem: The algorithm can be implemented in time $\underline{O(mn^2)}$.

graph exploration to find augm. path / blossom
$\rightarrow$ DFS traversal : $O(m)$

can contract only $O(n)$ blossoms until we find an augm. path

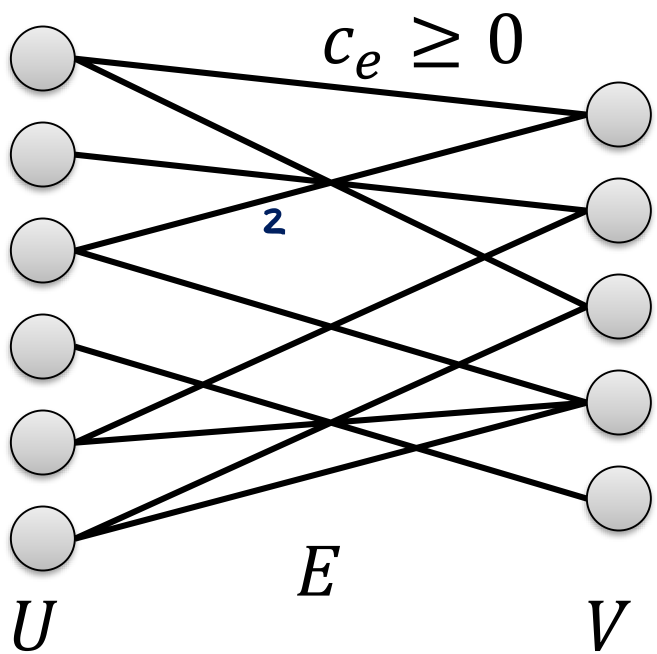
at most $\frac{n}{2}$ augm. paths

Maximum Weight Bipartite Matching

- Let's again go back to bipartite graphs...

Given: Bipartite graph $G = (U \dot\cup V, E)$ with edge weights $c_e \geq 0$

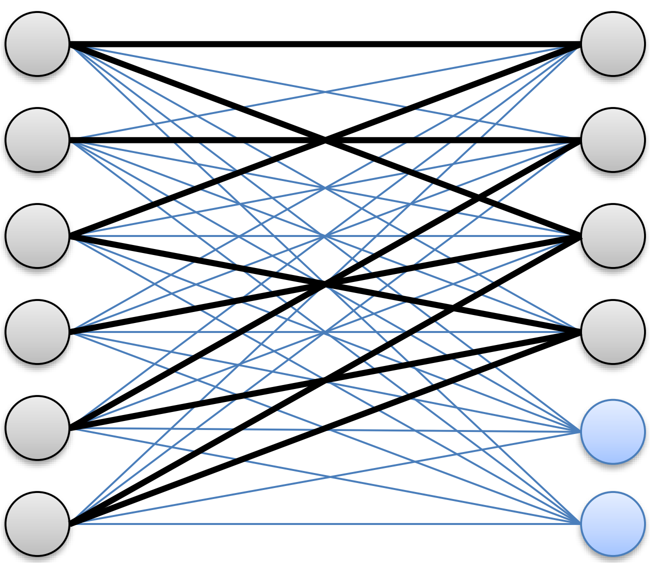
Goal: Find a matching M of maximum total weight



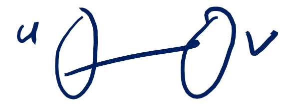
Minimum Weight Perfect Matching

Claim: Max weight bipartite matching is equivalent to finding a minimum weight perfect matching in a complete bipartite graph.

1. Turn into maximum weight perfect matching
   - add dummy nodes to get two equal-sized sides
   - add edges of weight 0 to make graph complete bipartite
2. Replace weights: $c'_e := \max_f \{c_f\} - c_e$



As an Integer Linear Program



- We can formulate the problem as an integer linear program

Var. x_{uv} for every edge $(u,v) \in U \times V$ to encode matching M:

$$x_{uv} = \begin{cases} 1, & \text{if } \{u,v\} \in M \\ 0, & \text{if } \{u,v\} \notin M \end{cases}$$

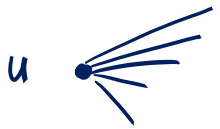
Minimum Weight Perfect Matching

$$\min \sum_{u,v \in U \times V} c_{u,v} \cdot x_{uv}$$

$$\forall u \in U: \quad \sum_{v \in V} x_{uv} = 1$$

$$\forall v \in V: \quad \sum_{u \in U} x_{uv} = 1$$

$$\forall u,v: \quad x_{uv} \in \{0,1\}$$



Linear Programming (LP) Relaxation

Linear Program (LP)

- Continuous optimization problem on multiple variables with a linear objective function and a set of linear side constraints

LP Relaxation of Minimum Weight Perfect Matching

- Weight c_{uv} & variable x_{uv} for ever edge $(u, v) \in U \times V$

$$\min \sum_{(u,v)\in U\times V} c_{uv} \cdot x_{uv}$$

$LP^* \leq ILP^*$

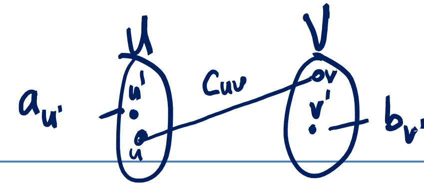
s. t.

$$\forall u \in U: \sum_{v\in V} x_{uv} = 1,$$

$$\forall v \in V: \sum_{u\in U} x_{uv} = 1$$

$$\forall u \in U, \forall v \in V: x_{uv} \geq 0$$

Dual Problem

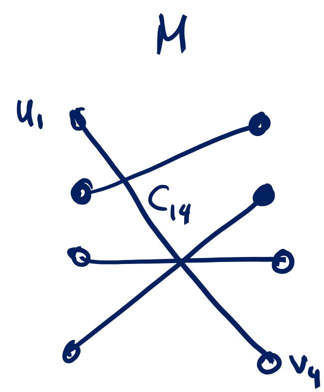


- Every linear program has a dual linear program
  - The dual of a minimization problem is a maximization problem
  - Strong duality: primal LP and dual LP have the same objective value

In the case of the minimum weight perfect matching problem

- Assign a variable $a_u \geq 0$ to each node $u \in U$
  and a variable $b_v \geq 0$ to each node $v \in V$
- **Condition:** for every edge $(\boldsymbol{u}, \boldsymbol{v}) \in \boldsymbol{U} \times \boldsymbol{V}$: $\boldsymbol{a_u + b_v \leq c_{uv}}$
- Given perfect matching M:

$$\sum_{(u,v) \in M} c_{uv} \geq \sum_{u \in U} a_u + \sum_{v \in V} b_v$$



$a_{u_1} + b_{v_4} \overset{=}{\leq} c_{14}$

Dual Linear Program

- Variables $a_u \geq 0$ for $u \in U$ and $b_v \geq 0$ for $v \in V$

$$\max \sum_{u \in U} a_u + \sum_{v \in V} b_v$$

$$s.t.$$

$$\forall u \in U, \forall v \in V: \; a_u + b_v \leq c_{uv}$$

- For every perfect matching M:

$$\sum_{(u,v) \in M} c_{uv} \geq \sum_{u \in U} a_u + \sum_{v \in V} b_v$$

Complementary Slackness

- A perfect matching M is optimal if

$$\sum_{(u,v)\in M} c_{uv} = \sum_{u\in U} a_u + \sum_{v\in V} b_v$$

- In that case, for every $(u, v) \in M$

$$\boldsymbol{w_{uv}} \coloneqq c_{uv} - a_u - b_v = 0$$

  – In this case, M is also an optimal solution to the LP relaxation of the problem
  – Every optimal LP solution can be characterized by such a property, which is then generally referred to as complementary slackness

- **Goal:** Find a dual solution a_u, b_v and a perfect matching such that the complementary slackness condition is satisfied!
  – i.e., for every matching edge (u, v), we want $w_{uv} = 0$
  – We then know that the matching is optimal!

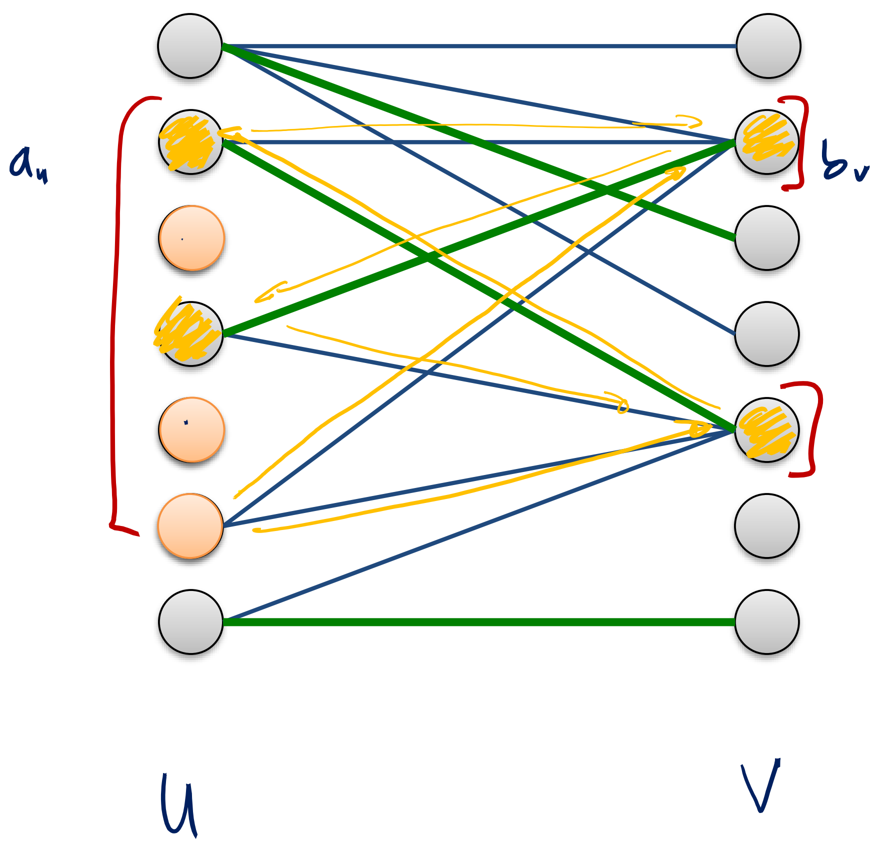
Algorithm Overview

- Start with any feasible dual solution a_u, b_v
  - i.e., solution satisfies that for all (u, v): $c_{uv} \geq a_u + b_v$
- Let E_0 be the edges for which $w_{uv} = 0$
  - Recall that $w_{uv} = c_{uv} - a_u - b_v$
- Compute maximum cardinality matching M of E_0
- All edges (u, v) of M satisfy $w_{uv} = 0$
  - Complementary slackness if satisfied
  - If M is a perfect matching, we are done
- If M is not a perfect matching, dual solution can be improved

Marked Nodes

Define set of marked nodes L:

- Set of nodes which can be reached on alternating paths on edges in E_0 starting from unmatched nodes in U



edges E_0 with $w_{uv} = 0$

optimal matching M

L_0: unmatched nodes in U

L: all nodes that can be reached on alternating paths starting from L_0

Marked Nodes

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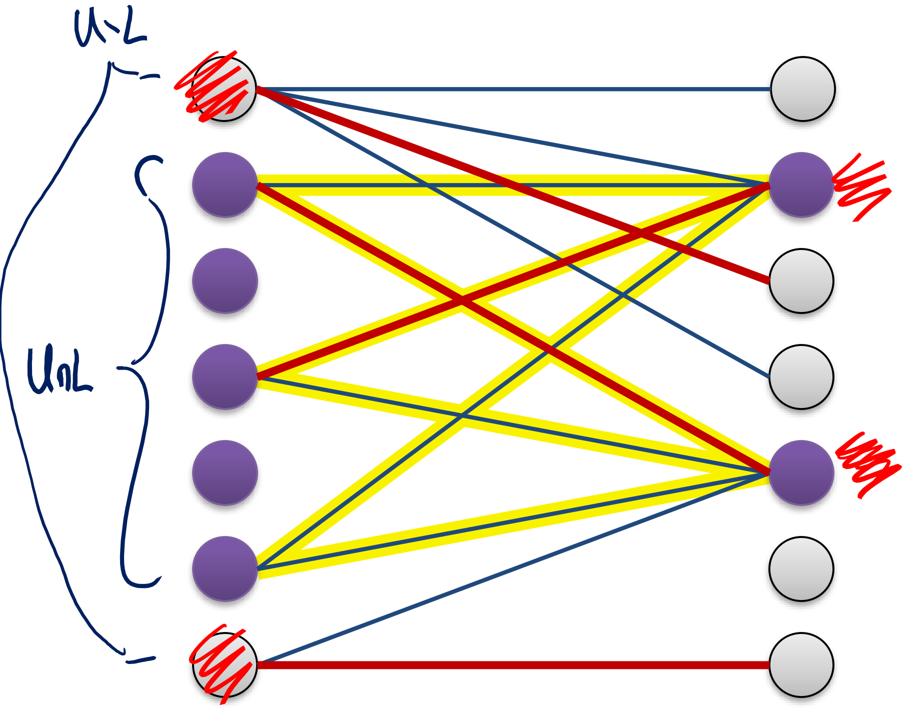

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L: all nodes that can be reached on alternating paths starting from L_0