

# Theoretical Computer Science - Bridging Course

## Winter Term 2017

### Exercise Sheet 1

hand in (electronically or hard copy) by 12:15 pm, Monday, October 30th, 2017

#### Exercise 1: Induction

(6 Points)

Prove the following summation formula by induction on  $n$ :

For all natural numbers  $n \geq 1$  it holds:  $\sum_{k=1}^n k^2 = \frac{1}{6}n(n+1)(2n+1)$

#### Sample Solution

For  $n = 1$  we have to show that  $\sum_{k=1}^1 k^2 = \frac{1(1+1)(2+1)}{6}$ , which is obviously true because both sides are equal 1.

Now assume the statement is true for  $n$ . It follows that

$$\sum_{k=1}^{n+1} k^2 = \sum_{k=1}^n k^2 + (n+1)^2 = \frac{1}{6}n(n+1)(2n+1) + (n+1)^2 \quad (1)$$

$$= \frac{1}{6}n(n+1)(2n+1) + n^2 + 2n + 1 = \frac{1}{6}(2n^3 + 3n^2 + n) + n^2 + 2n + 1 \quad (2)$$

$$= \frac{1}{6}(2n^3 + 9n^2 + 13n + 6) \quad (3)$$

$$= \frac{1}{6}(n+1)(n+2)(2n+3). \quad (4)$$

which shows that the statement also holds for  $n+1$ . The second equality in the equation above comes from the assumption for  $n$ .

#### Exercise 2: Sets & More

(8 Points)

Either prove the following statements or show that they are not true.

1.  $\{(x, y) \in \mathbb{Z} \times \mathbb{Z} \mid x + y = 0\} \subseteq \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \mid x \cdot y = 0\}$
2.  $A \setminus (B \cap C) = (A \setminus B) \cup (A \setminus C)$
3.  $A \subsetneq B$  implies  $C \cup A \subsetneq C \cup B$ .

## Sample Solution

1.  $(5, -5) \in \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \mid x + y = 0\}$  but  $(5, -5) \notin \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \mid x + y = 0\}$ .
2.  $\subseteq$ : Let  $x \in A \setminus (B \cap C)$ , which implies that  $x \in A$ . Assume that  $x$  is not contained in the right hand side, i.e.,  $x \notin (A \setminus B) \cup (A \setminus C)$  which implies

$$(1): \quad x \notin (A \setminus B) \text{ and}$$

$$(2): \quad x \notin (A \setminus C).$$

As  $x \in A$  it (1) implies  $x \in B$  and (2) implies  $x \in C$ . That means  $x \in B \cap C$ . But this is a contradiction to  $x \in A \setminus (B \cap C)$ .

$\supseteq$ : Let  $x \in A \setminus B \cup A \setminus C$  which implies  $x \in A$ . Assume that  $x$  is not contained in the right hand side, i.e.,  $x \notin A \setminus (B \cap C)$  which implies  $x \in B \cap C$ . But then  $x \notin A \setminus B$  and  $x \notin A \setminus C$  which implies  $x \notin A \setminus (B \cap C)$ , a contradiction.

3. This statement does not hold. Let  $A = \{1, 2\}$   $B = \{1, 2, 3\}$  and  $C = \{3, 4, 5\}$ . Then  $A \subsetneq B$  but  $A \cup C \subsetneq B \cup C$  is equivalent to  $\{1, 2, 3, 4, 5\} \subsetneq \{1, 2, 3, 4, 5\}$  which is not true.

## Exercise 3: Where is the even degree node?

(6 Points)

A *simple graph* is a graph without self loops, i.e., every edge of the graph is an edge between two distinct nodes. The degree  $d(v)$  of a node  $v \in V$  in an undirected graph  $G = (V, E)$  is the number of its neighbors, i.e.,

$$d(v) = |\{u \in V \mid \{v, u\} \in E\}|.$$

Show that every simple graph with an odd number of nodes contains a node with even degree.

*Hint: Consider the sum  $D = \sum_{v \in V} d(v)$  of all degrees. Is  $D$  odd or even?*

## Sample Solution

Let  $G = (V, E)$  be a graph.

The degree  $d(v)$  of a node  $v$  is the number of incident edges, i.e., edges that contain  $v$ . Let  $D$  be the sum of all degrees.

Every edge contributes two to  $D$ , hence  $D = 2|E|$ , i.e.,  $D$  is even.

We can also write  $D$  as the sum of the degrees of the nodes in  $V$ , i.e.,  $D = \sum_{v \in V} d(v)$ . The sum of an odd number of odd numbers is odd. If all  $d(v)$  were odd also  $D$  would be odd as  $|V|$  is odd. Thus there has to be a node with even degree.