



# Algorithms Theory

## Exercise Sheet 5

Due: Tuesday, 8th of December 2020, 4 pm

### Exercise 1: Amortized Analysis

(10 Points)

Consider a binary min-heap data structure that supports the two operations `insert` and `delete-min`. The heap is initially empty and we assume that its number of elements never exceeds  $n$ .

- (a) Use the *accounting* method to show that we can consider the amortized cost of `insert` to be  $O(\log n)$  and the amortized cost of `delete-min` to be  $O(1)$ . (3 Points)
- (b) Show the statement from part (a), this time using the *potential function* method. (5 Points)
- (c) We would like to amortize the costs differently such that the amortized cost of `insert` is  $O(1)$  and the amortized cost of `delete-min` is  $O(\log n)$ . Either define a feasible *potential function* that yields these amortized costs or argue why this is not possible. (2 Points)

### Exercise 2: Union Find - Linked List Implementation

(8 Points)

In the lecture, we have seen a linked list implementation where each linked list has a pointer to the first *and* last element. Describe an alternative implementation that uses only *one* of these pointers. Your scheme should still allow for the union-by-size heuristic and should not increase the asymptotic running time of the operations.

### Exercise 3: Union Find - Disjoint-Set Forests

(8 Points)

- (a) Give a sequence of  $m$  `make-set`, `union`, and `find` operations,  $n$  of which are `make-set` operations, that takes  $\Omega(m \log n)$  time when we use union by rank only. (3 Points)
- (b) Suppose that we wish to add the operation `print-set`, which is given a node  $x$  and prints all the members of  $x$ 's set, in any order. Show how to add this feature to the disjoint-set forest implementation such that `print-set` takes time linear in the number of members of  $x$ 's set and the asymptotic running times of the other operations are unchanged. Assume that we can print each member of the set in  $O(1)$  time. (5 Points)