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Algorithms Theory

Exercise Sheet 9

Due: Tuesday, 19th of January 2021, 4 pm

Exercise 1: Matching vs Vertex Cover

Given an undirected Graph G = (V, E), a vertex cover of G is a set of nodes $S \subseteq V$ such that for all $\{u, v\} \in E$, we have $\{u, v\} \cap S \neq \emptyset$. A minimum vertex cover is a vertex cover of minimum size.

a) Show that for a maximum matching M^* and a minimum vertex cover S^* we have $|M^*| \le |S^*|$. (2 Points)

Next we want to show that in bipartite graphs, it also holds $|S^*| \leq |M^*|$.

b) Recall that we can solve the maximum bipartite matching problem by reduction to maximum flow. Also recall that if we are given a maximum matching M^* (and thus a maximum flow of the corresponding flow network), we can find a minimum *s*-*t* cut by considering the residual graph. Describe how such a minimum cut looks like.

Hint: Consider the set of all nodes which can be reached from an unmatched node on the left side via an alternating path. (3 Points)

- c) Use the above description to show that any bipartite graph G has a vertex cover S^* of size $|M^*|$. (3 Points)
- d) Show that the same thing is not true for general graphs by showing that for every $\varepsilon > 0$, there exists a graph G = (V, E) for which $|S^*| \ge (2 \varepsilon)|M^*|$.

Hint: First try to find any graph for which $|S^*| > |M^*|$.

Exercise 2: Contention Resolution

Show that for the randomized algorithm for contention resolution from the lecture, the expected time until all processes have been successful is $O(n \log n)$.

(10 Points)

(10 Points)

(2 Points)