



# Algorithm Theory

## Exercise Sheet 13 - Bonus\*

Due: Tuesday, 1st of February, 2021, 4 pm

\*Bonus points can be earned normally, but do not increase the threshold for the “Studienleistung”.

### Exercise 1: Online Vertex Cover

(10\* Points)

Let  $G = (V, E)$  be an unweighted undirected graph. Consider the following online version of the *minimum vertex cover* problem. Initially, we are given the set of nodes  $V$  and an empty vertex cover  $S = \emptyset$ . Then, the edges appear one-by-one in an online fashion. When a new edge  $\{u, v\}$  appears, the algorithm needs to guarantee that the edge is covered (i.e., if this is not already the case, at least one of the two nodes  $u$  and  $v$  needs to be added to  $S$ ). Once a node is in  $S$  it cannot be removed from  $S$ .

- (a) Provide a deterministic online algorithm with strict competitive ratio at most 2. (1 Point)
- (b) Show that any deterministic online algorithm for the online vertex cover problem has strict competitive ratio at least 2. (4 Points)
- (c) Use Yao’s principle to show that any randomized online algorithm for the online vertex cover problem has a strict competitive ratio at least  $3/2$ . (5 Points)

### Exercise 2: Maximum Cut

(10\* Points)

Let  $G = (V, E)$  be an unweighted undirected graph. A *maximum cut* of  $G$  is a cut whose size is at least the size of any other cut in  $G$ . We consider an online version of the maximum cut problem, where the nodes  $V$  of a graph  $G = (V, E)$  appear in an online fashion. The algorithm should partition the nodes  $V$  into two sets  $A$  and  $B$  such that induced cut is as large as possible. Whenever a new node appears we also learn its edges to the nodes that have *already appeared before*. An incoming node has to be immediately assigned to either  $A$  or  $B$  and that decision is final.

- (a) Describe a deterministic online maximum cut algorithm with *strict competitive ratio* at least  $1/2$ .  
*Hint: An online algorithm for a maximization problem is said to have strict competitive ratio  $\alpha$  if it guarantees that  $\text{ALG} \geq \alpha \cdot \text{OPT}$ , where ALG and OPT are the solutions of the online algorithm and of an optimal offline algorithm, respectively.* (5 Points)
- (b) Show that no deterministic online algorithm for the online maximum cut problem can have a (constant) strict competitive ratio that is better (larger) than  $1/2$ . (5 Points)