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## Combinatorial Optimization

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### Exercise 1 (Optimizing vs. Finding Feasible Solutions)

(1) Consider the LPs

$$\max\{c^\top x : A'x \leq b', A''x \leq b'', x \geq 0\} \quad (*)$$

where  $b' \geq 0$  and  $b'' < 0$  and

$$\min\{(1^\top A'')x + 1^\top y : A'x \leq b', A''x + y \geq b'', x, y \geq 0\}. \quad (**)$$

from the lecture. Observe that  $(x, y)^\top = (0, 0)^\top$  is a feasible vertex for  $(**)$  and we may hence run SIMPLEX on that LP. Show that the LP  $(*)$  is feasible if and only if the optimum value for  $(**)$  is exactly  $1^\top b''$ .

This shows that if we are able to optimize (with a given initial vertex) then we can decide if a certain LP of interest is feasible.

(2) Now show that the converse is also true. Suppose that you are interested in the LP

$$\max\{c^\top x : Ax \leq b\}. \quad (***)$$

Assume that you can decide if any given LP is feasible and derive a feasible solution, if so. Use this knowledge to show that you can derive the optimum solution of  $(***)$  or infer that it is infeasible.

### Exercise 2 (Vertices with Capacities)

Recall the definition of networks from the lecture and note that only the edges (but not the vertices) have capacities. Now we want to extend this. Let  $d : V \rightarrow \mathbb{R}^+$  be a function, called the *capacity* of a vertex and we define that a function  $f : E \rightarrow \mathbb{R}^+$  is a *feasible* flow if it satisfies the flow-conditions from the lecture and has the additional property

$$\sum_{e \in \delta^-(v)} f(e) \leq d(v)$$

for all vertices  $v$ . That is, the incoming flow does not exceed the vertex-capacity. Formulate the maximum flow problem with vertex-capacities as an ordinary maximum flow problem.

### **Exercise 3 (Isolate the Commander)**

A commander is located at a vertex in an undirected communication network and his soldiers are located at vertices denoted by a set  $S$ . Let  $u_{ij}$  be the effort required to eliminate the edge  $ij$  from the network. How can you determine the minimal effort needed so that the commander can not communicate with any of his soldiers?

### **Exercise 4 (Guest Shuffle)**

You are organizing a dinner and lay  $n$  tables. You invite  $m$  families to join the dinner and family  $i$  has  $a_i$  members. Furthermore table  $j$  has  $b_j$  seats. In order to boost the inter-family-communication you want to make sure that no two members of the same family are at the same table (if this is possible). Formulate this seating arrangement problem as a maximum flow problem.