
Combinatorial Optimization

Exercise 1 (Set Cover)

Construct an instance for which the approximation-factor of H_n for the GREEDY algorithm for SET COVER is tight.

Exercise 2 (Set Multicover)

The SET MULTICOVER problem is a generalization of SET COVER, where each element e needs to be covered by a specified integer number r_e of times. The objective is to cover all elements up to their coverage requirements at minimum cost. You can assume that the cost of picking a set S k times is $kc(S)$.

Extend the GREEDY algorithm for SET COVER to handle this problem. Show that it is a H_n -approximation algorithm.

Hint. One way to do it is to apply the dual fitting method.

Exercise 3 (Maximum Coverage)

The MAXIMUM COVERAGE problem is the following: We are given a universe U of n elements with non-negative weights specified, a collection of subsets of U , S_1, \dots, S_ℓ and an integer k . Our task is to pick k sets as to maximize the total weight of the elements covered.

Show that the obvious algorithm of greedily picking the best set in each iteration until k sets are picked achieves an approximation factor of

$$1 - \left(1 - \frac{1}{k}\right)^k > 1 - \frac{1}{e}.$$