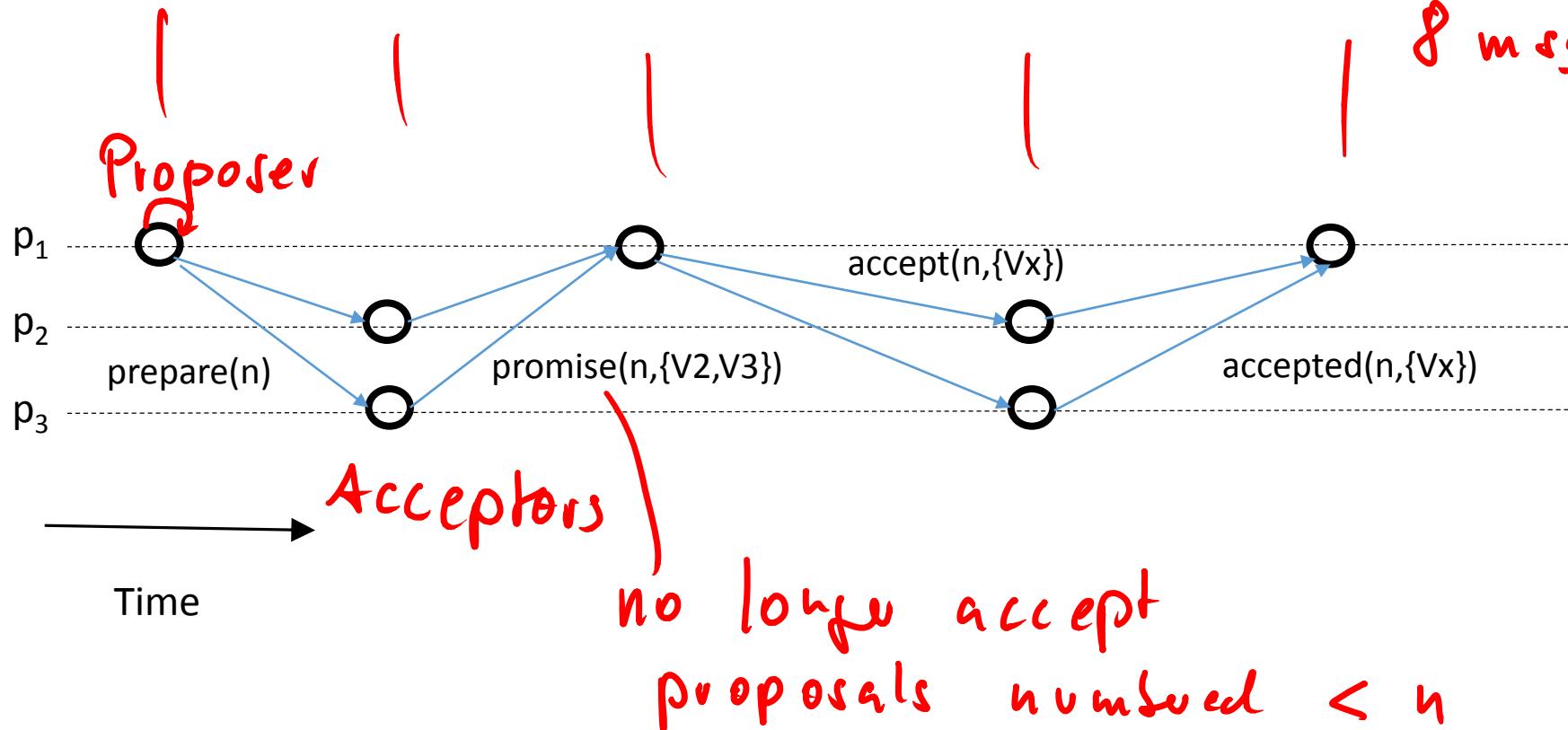


# Distributed Systems

Exercise 5

a) p<sub>1</sub> wants to publish a value and no failures occur.

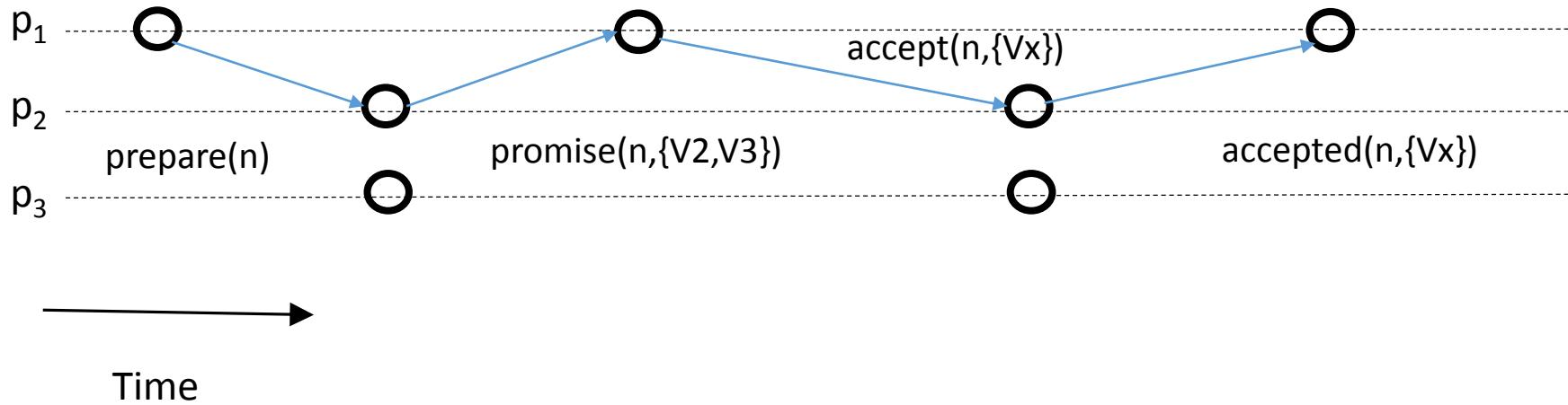
4 rounds  
8 msg



Acceptors made promise before  $\rightarrow$  send most recent value to prop / Prop chosen most recent

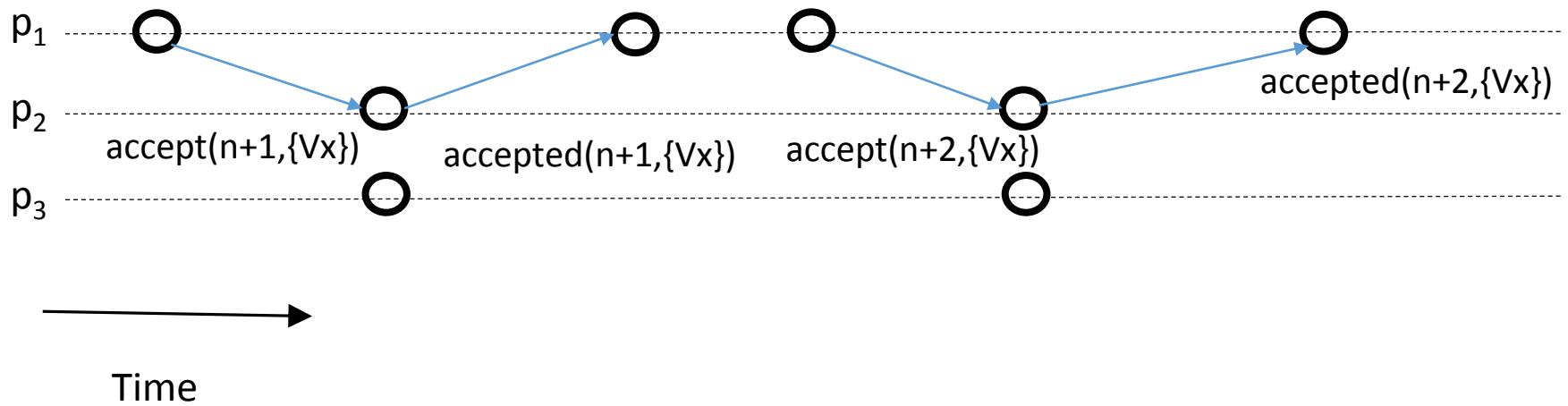
b) p1 uses an optimized version, where it only communicates with a quorum, again without failures!

4 round trips  
4 messages



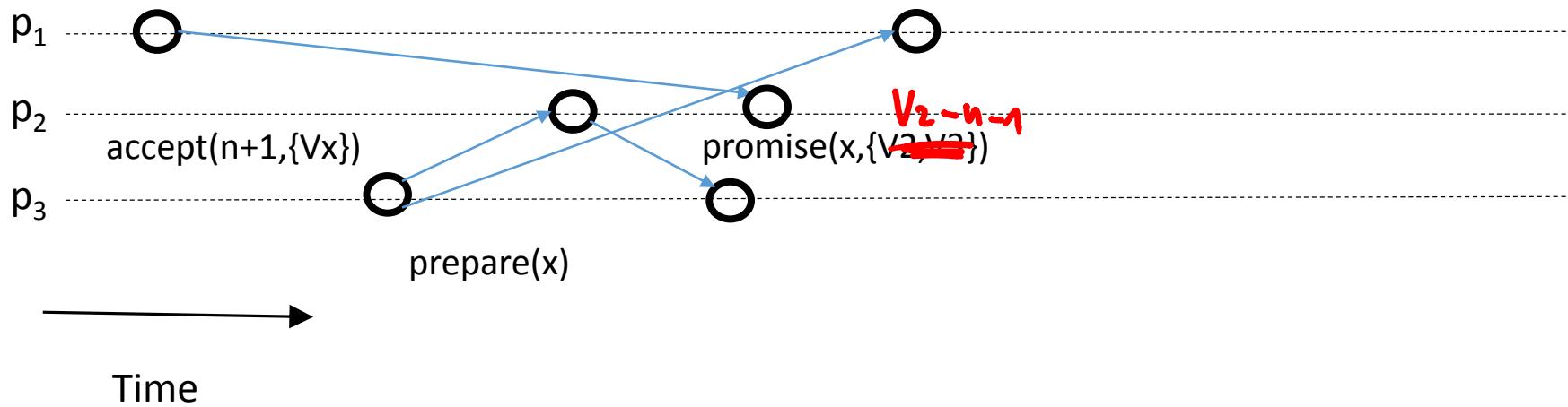
On failure of  $p_2$ , start a new proposal with all acceptors  $\rightarrow p_3 \rightarrow$  quorum reached!

c) Another iteration of b. optimized by using Multi-Paxos  
(System is already in a steady state).



~~c) Another iteration of b. optimized by using Multi-Paxos  
(System is already in a steady state).~~

d)



$x$  smaller : deny (or ignore)

$x$  greater :  $p_n$  promises + sends value  
 $p_1$  fails acceptance  $\Rightarrow$  regular paxos

P<sub>1</sub> one dead in Bulk - Paxos



if P<sub>2/3</sub> - prepare(x) small than existing  
proposal number, keep proposing with  
increasing number

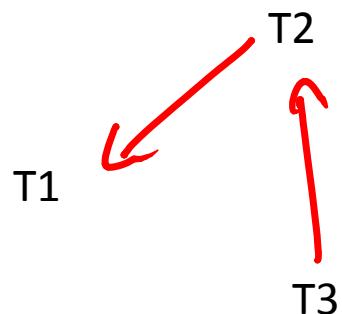
$S_1: R_3X R_2Y W_2Y R_1Y W_1Y R_2X W_2X R_1X W_1X W_3Z.$

$X: \bar{T}_3 \rightarrow \bar{T}_2 \rightarrow \bar{T}_1$

$Y: \bar{T}_2 \rightarrow \bar{T}_1$

$Z:$

yes



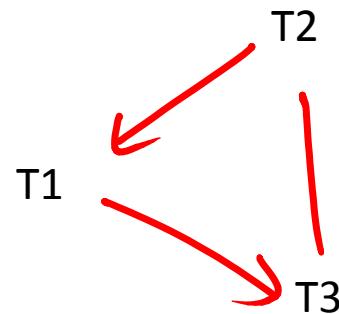
$\bar{T}_3 \rightarrow \bar{T}_2 \rightarrow \bar{T}_1$

equivalent

$S_2: R_3X R_2Y W_2Y R_1Y \underline{W_1Y} R_2X W_2X R_1X W_1X \underline{W_3Y}$

$X: T_3 \rightarrow T_2 \rightarrow T_1$

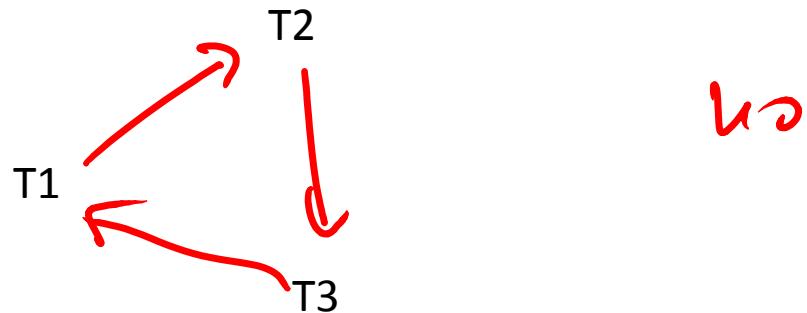
$\gamma T_2 \rightarrow T_1 \rightarrow T_3$



no  $\rightarrow$   
cycle

$S_3: R_1Y W_1Y R_2Y W_2Y R_2X W_2X R_3Z W_3X R_1X W_1X.$

$$\begin{matrix} X: & \overline{T}_2 \rightarrow \overline{T}_3 \rightarrow \overline{T}_1 \\ Y: & \overline{T}_1 \rightarrow \overline{T}_2 \end{matrix}$$



$T_1$ : RA WA

$T_2$ : RA WA

$T_3$ : RA WA

- a) What is the number of serial schedules?

6 - all permutations

- b) Which schedules exists that are **serializable**, but not **serial**?

0 same data access order

$R_i R_j \quad W_i W_j$

$R_i R_j \quad U_j W_i \rightarrow$  not possible

due to  
conflict

T1 : RA WC

T2 : RB WA

T3 : RC WD

- a) How many schedules exist that are not serializable?

0 no cycle possible

- b) Does using 2PL permit all serializable schedules?

$S = R_1 R_2 R_3 A \mid W_2 A \mid W_1 C \quad R_3 C \quad W_3 D$

$L_1 A \quad U_1 A$

not possible

$T_1 \rightarrow T_2 \rightarrow T_3$