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Algorithms and Data Structures
Summer Term 2019
Exercise Sheet 6

Exercise 1: Master Theorem for Recurrences

Use the *Master Theorem* for recurrences, to fill the following table. That is, in each cell write $\Theta\big(g(n)\big)$, such that $T(n) \in \Theta\big(g(n)\big)$ for the given parameters $a, b, f(n)$. Assume $T(1) \in \Theta(1)$. Additionally, in each cell note the case you used (1st, 2nd or 3rd by the order given in the lecture). We filled out one cell as an example.

$T(n)=aT(\frac{n}{b})+f(n)$	$a = 16, b = 2$	$a = 1, b = 2$	$a = b = 3$
$f(n) = 1$	$\Theta(n^4)$, 1st		
$f(n) = n$			
$f(n) = n^4$			

Exercise 2: Peak Element

You are given an array $A[1 \dots n]$ of n integers and the goal is to find a peak element, which is defined as an element in A that is equal to or bigger than its direct neighbors in the array. Formally, $A[i]$ is a peak element if $A[i-1] \leq A[i] \geq A[i+1]$. To simplify the definition of peak elements on the rims of A, we introduce *sentinel-elements* $A[0] = A[n+1] = -\infty$.

(a) Give an algorithm with runtime $\mathcal{O}(\log n)$ (measured in the number of read operations on the array) which returns the position i of a peak element.

(b) Prove that your algorithm always returns a peak element, give a recurrence relation for the runtime and use it to prove the runtime.