

CONGEST model

bandwidth limitations

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Part of the slides are from Jukka Suomela

CONGEST model

- **LOCAL** model: arbitrarily large messages
- **CONGEST** model: $O(\log n)$ -bit messages

CONGEST model

- Any of these can be encoded in $O(\log n)$ -bit messages:
 - node identifier
 - number of nodes
 - number of edges
 - distance between two nodes ...

CONGEST model

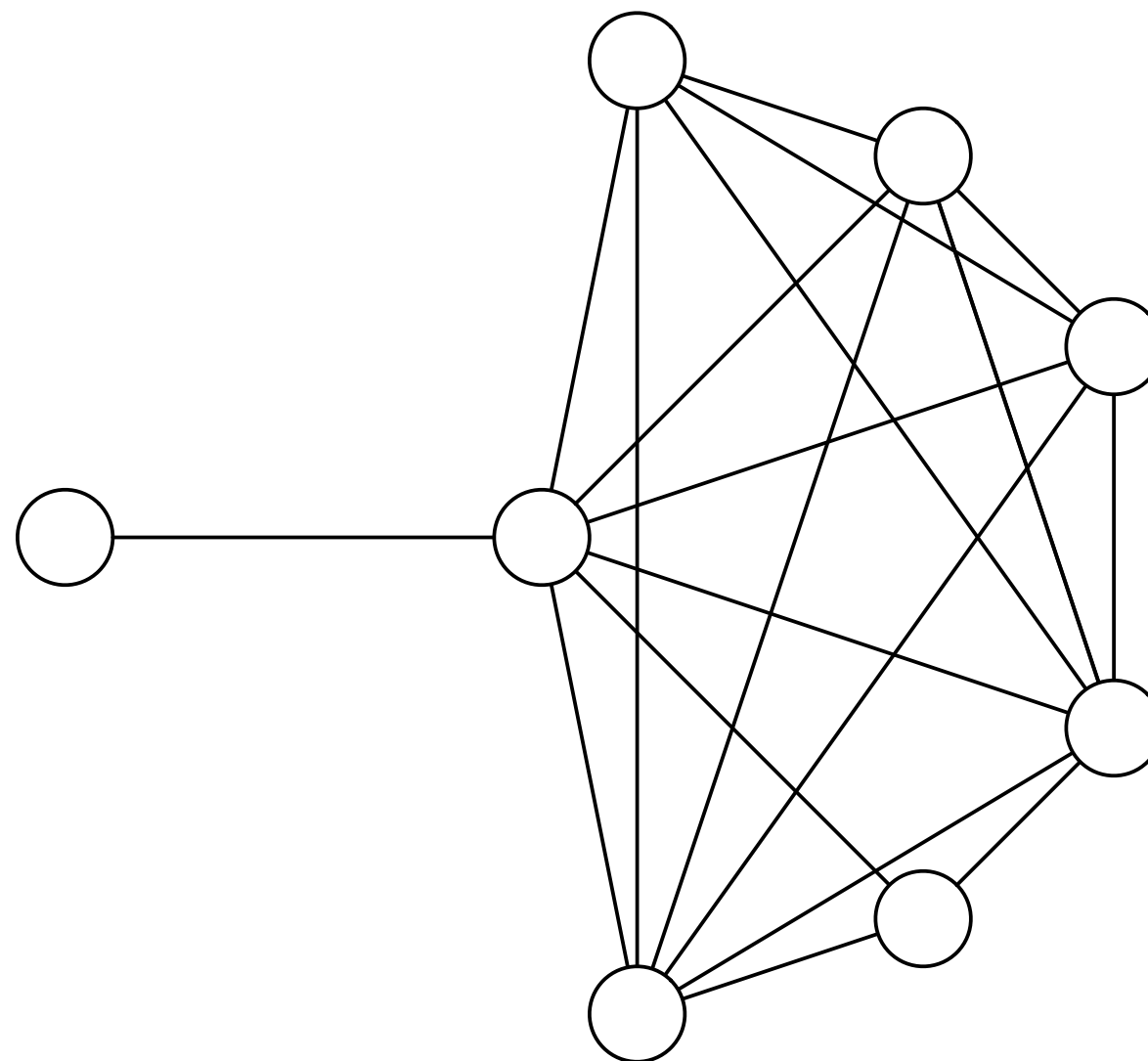
- Many algorithms that we have seen use small messages
 - can be used directly in CONGEST:
 - Example: coloring algorithms seen in the lectures
- There are some **exceptions**

Solving everything in LOCAL

- **Gather** the whole graph + solve the problem locally (e.g., by brute force)
 - $O(\text{diam}(G))$ rounds
 - See animation here:
<https://jukkasuomela.fi/animations/local-horizon.gif>

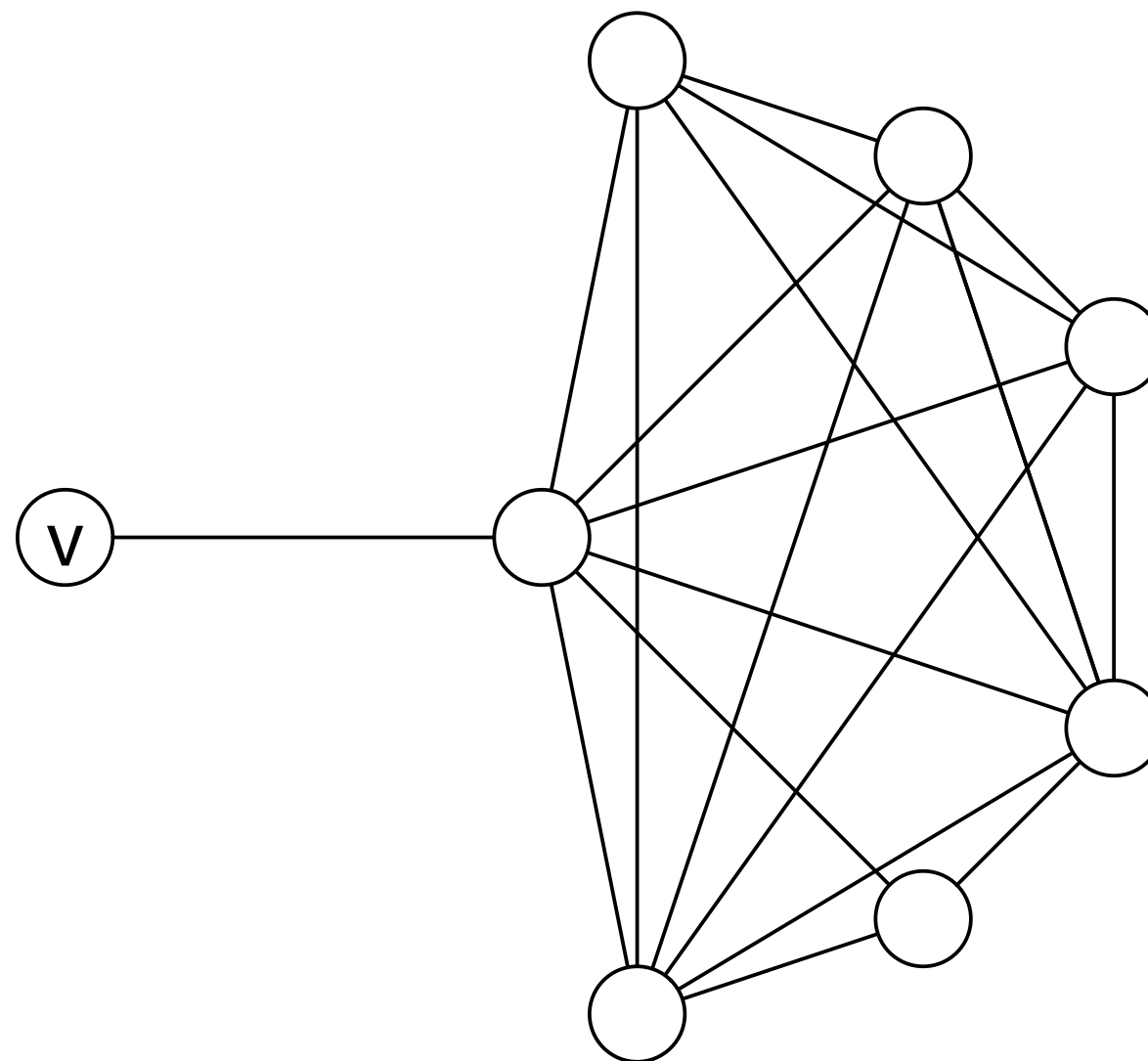
Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
 - Nodes have IDs from 1 to n



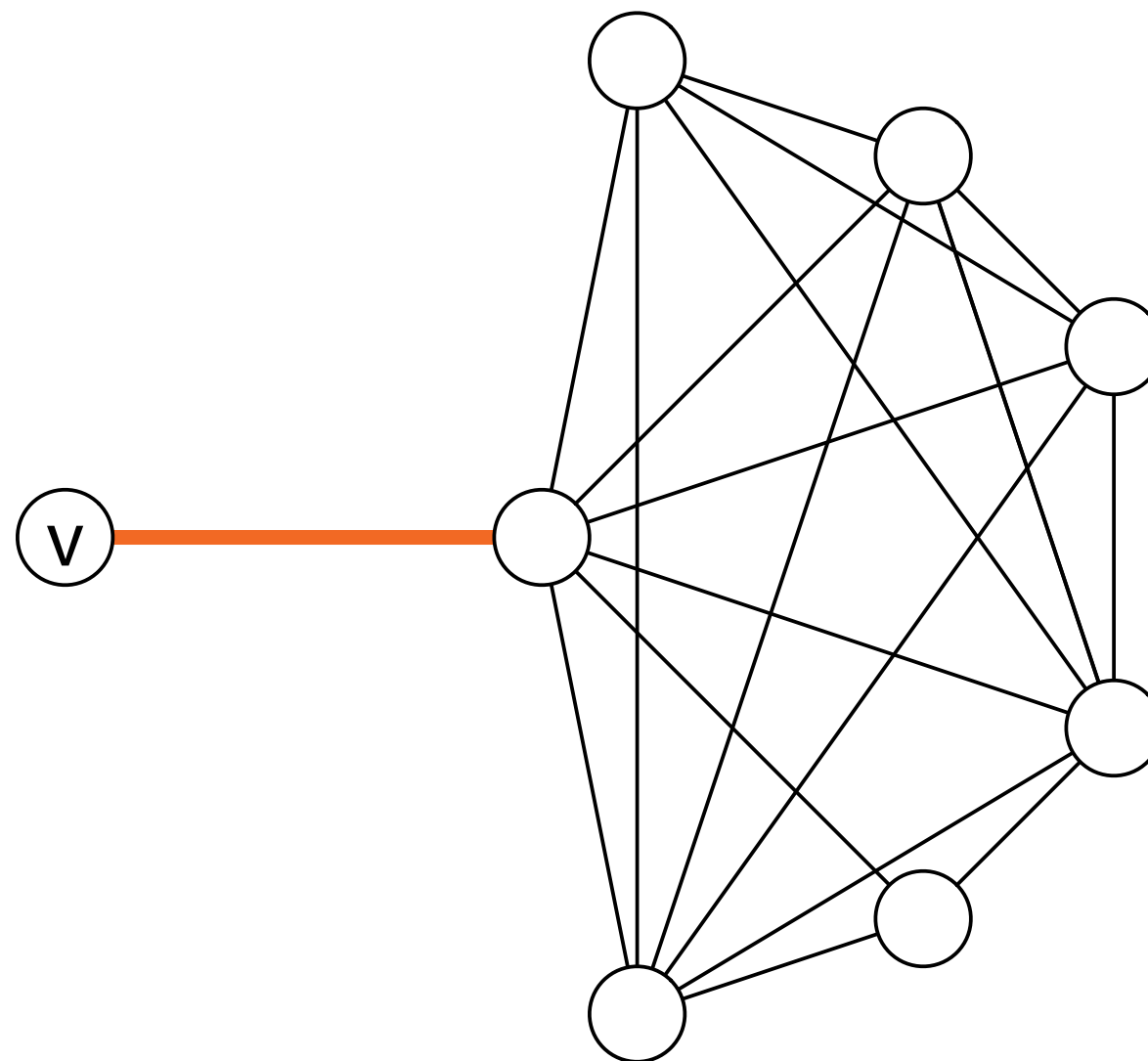
Algorithm Gather

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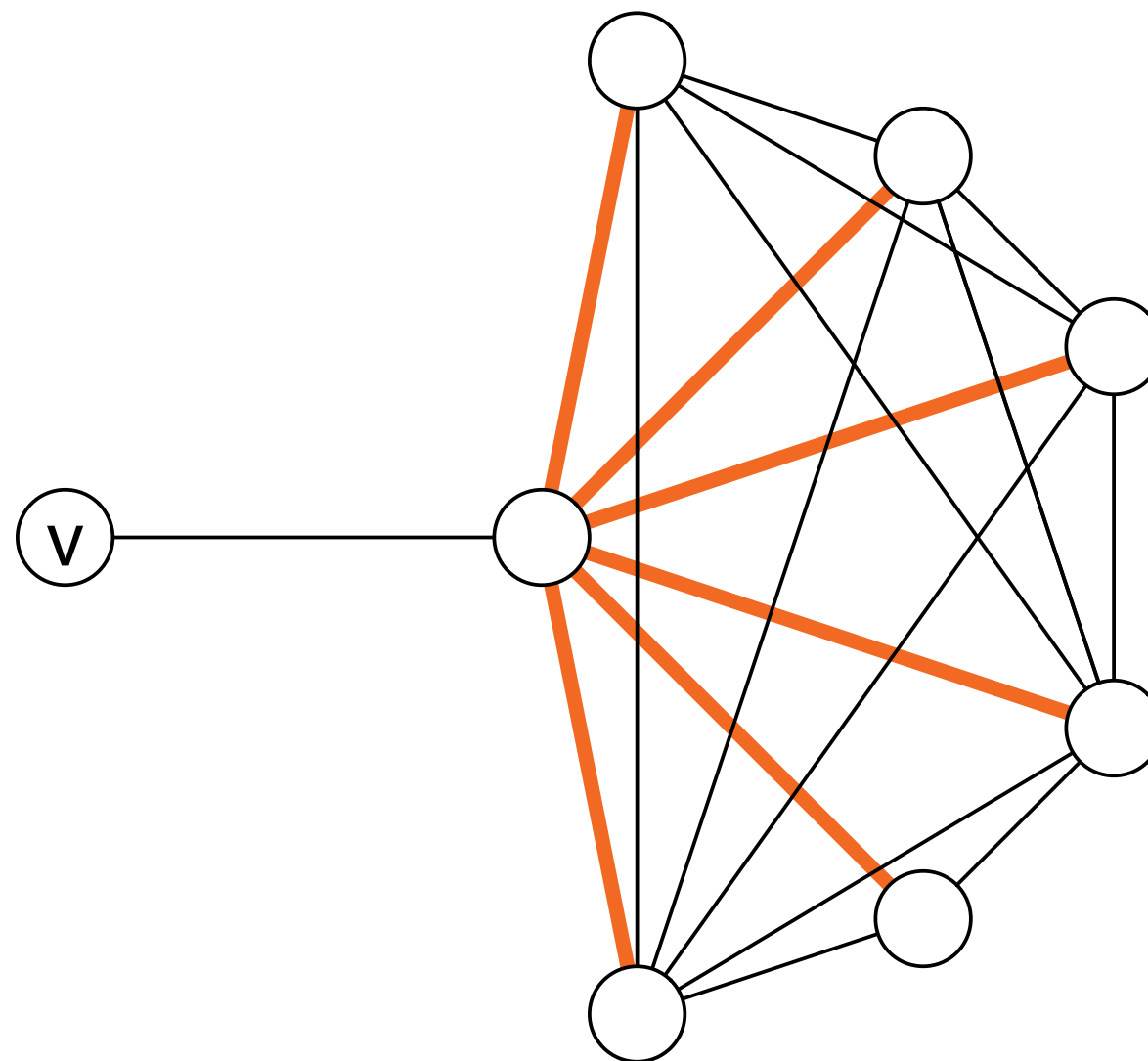
Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
 - Round 1



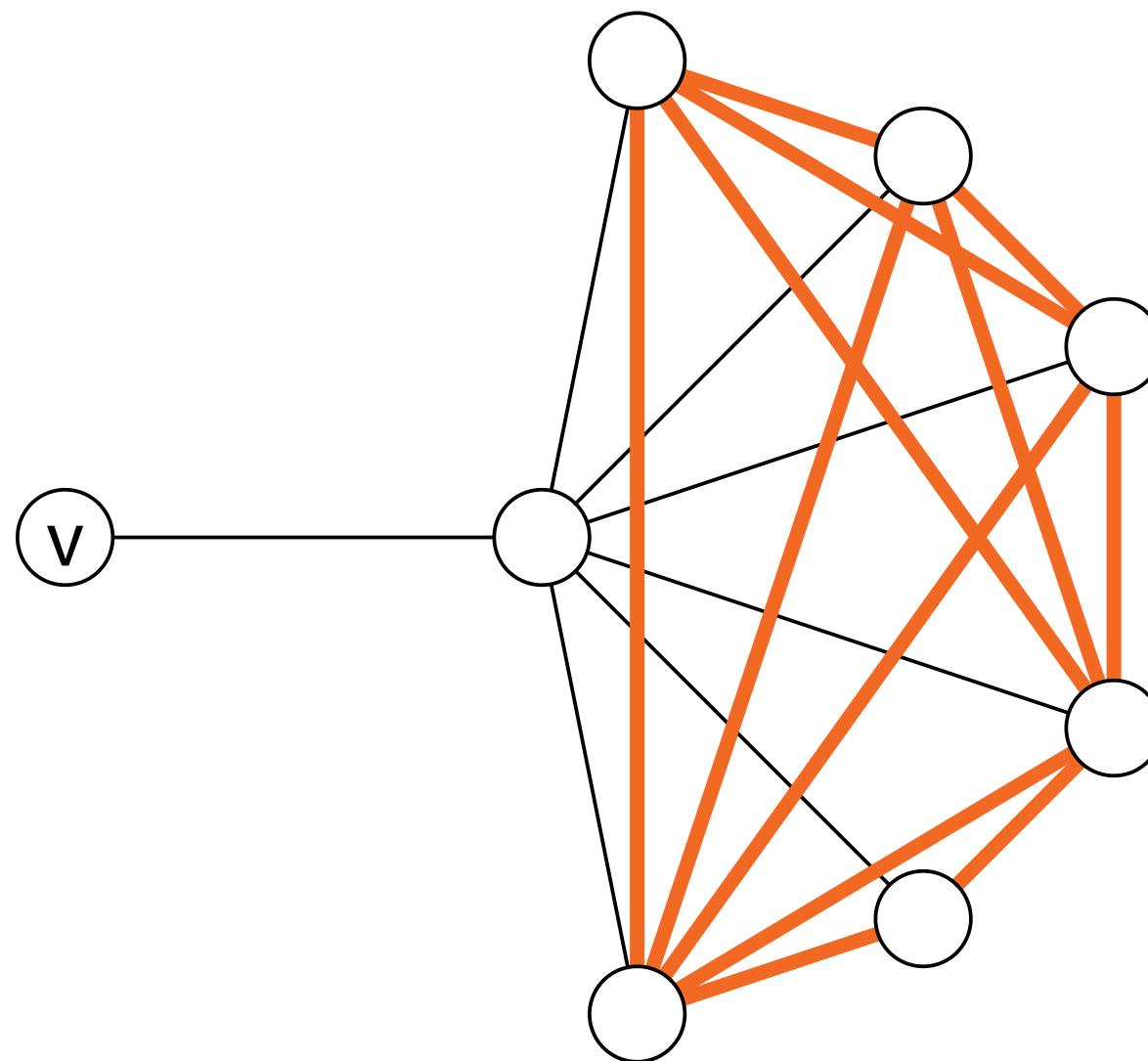
Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
 - Round 2



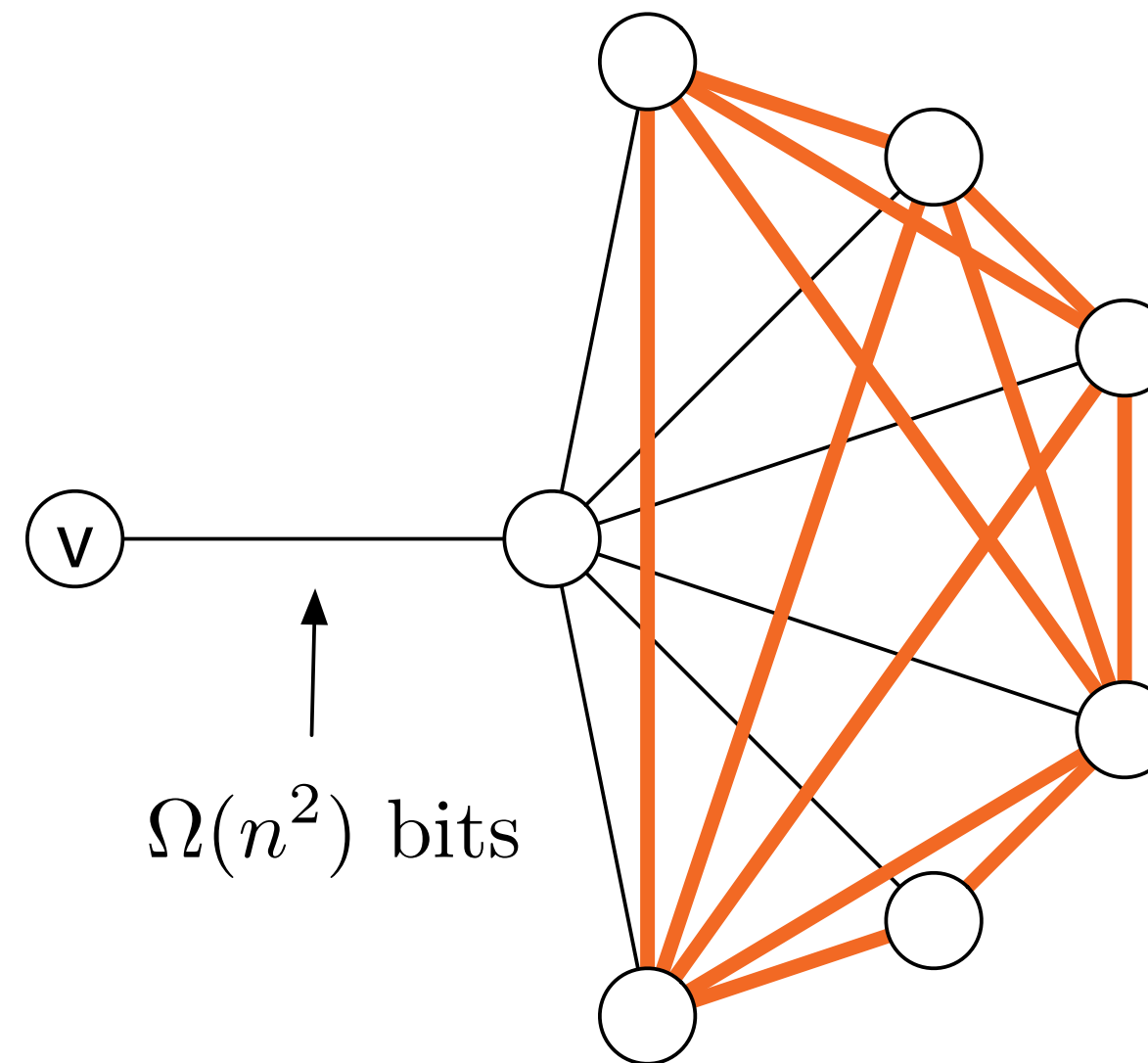
Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
 - Round 3



Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
 - Round 3, send the adjacency matrix



Algorithm Gather

- May need $\Omega(n^2)$ -bit messages
- Cannot directly be used in CONGEST
- **Exercise:** gather all the graph in CONGEST in $O(|E|)$ rounds

CONGEST model

- $O(n)$ time trivial in the LOCAL model
 - brute force approach: Gather + solve locally
- $O(n)$ time **non-trivial** in the **CONGEST** model

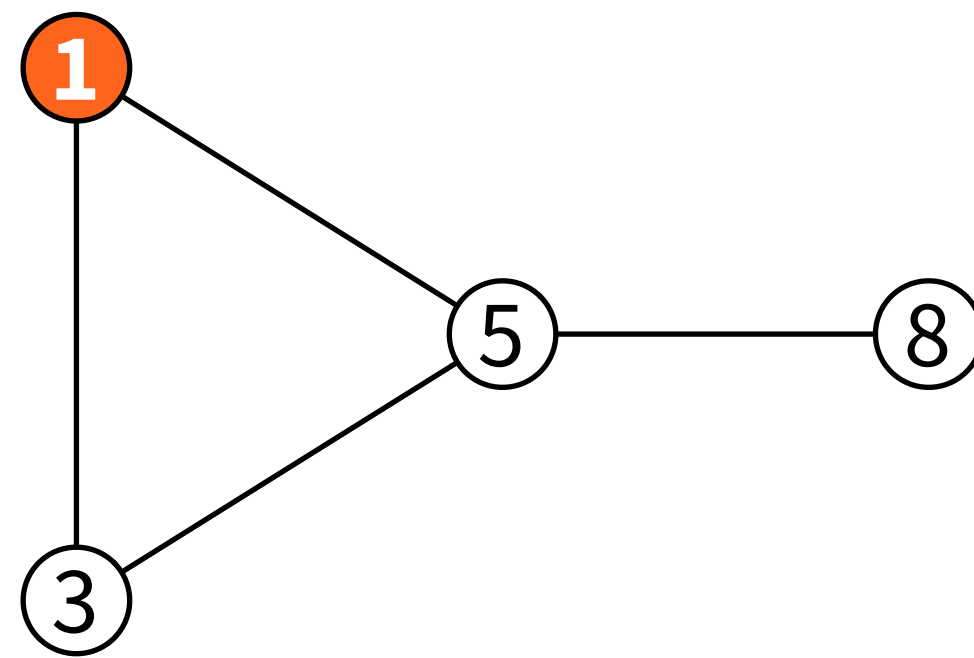
Today

- How to find **all-pairs shortest paths (APSP)** in **$O(n)$ time** in the CONGEST model [Holzer, Wattenhofer]
- Lower bound of **$\Omega(n / \log n)$** rounds for **APSP** [Frischknecht, Holzer, Wattenhofer]

(Complexity of APSP in CONGEST: **$\Theta(n / \log n)$**)

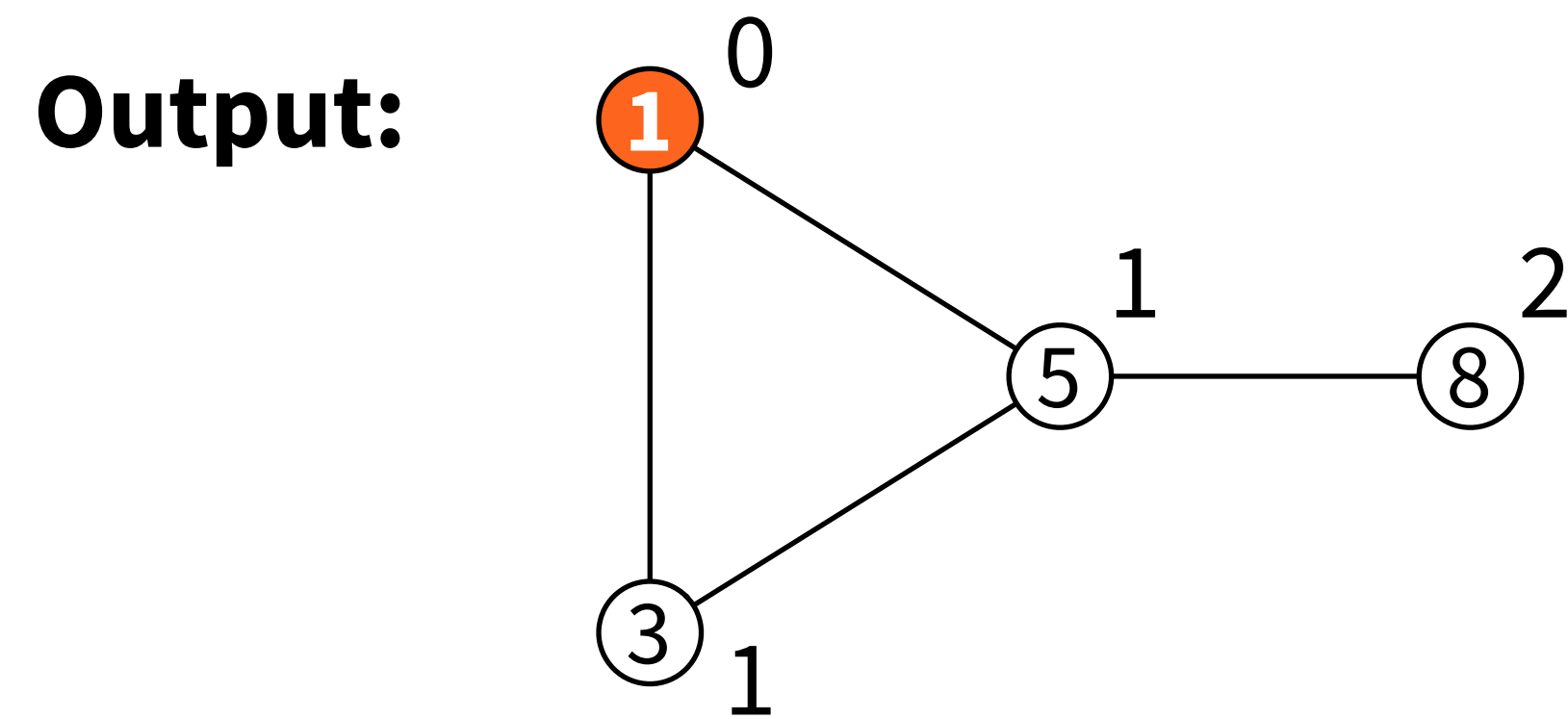
Single-source shortest paths

Input:



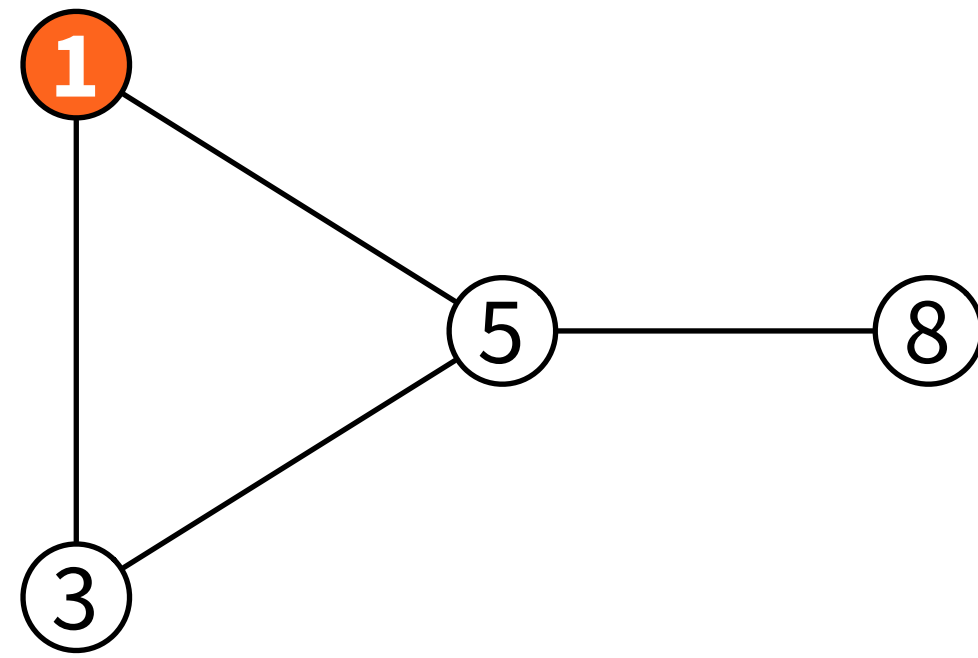
Single-source shortest paths

Distances from s



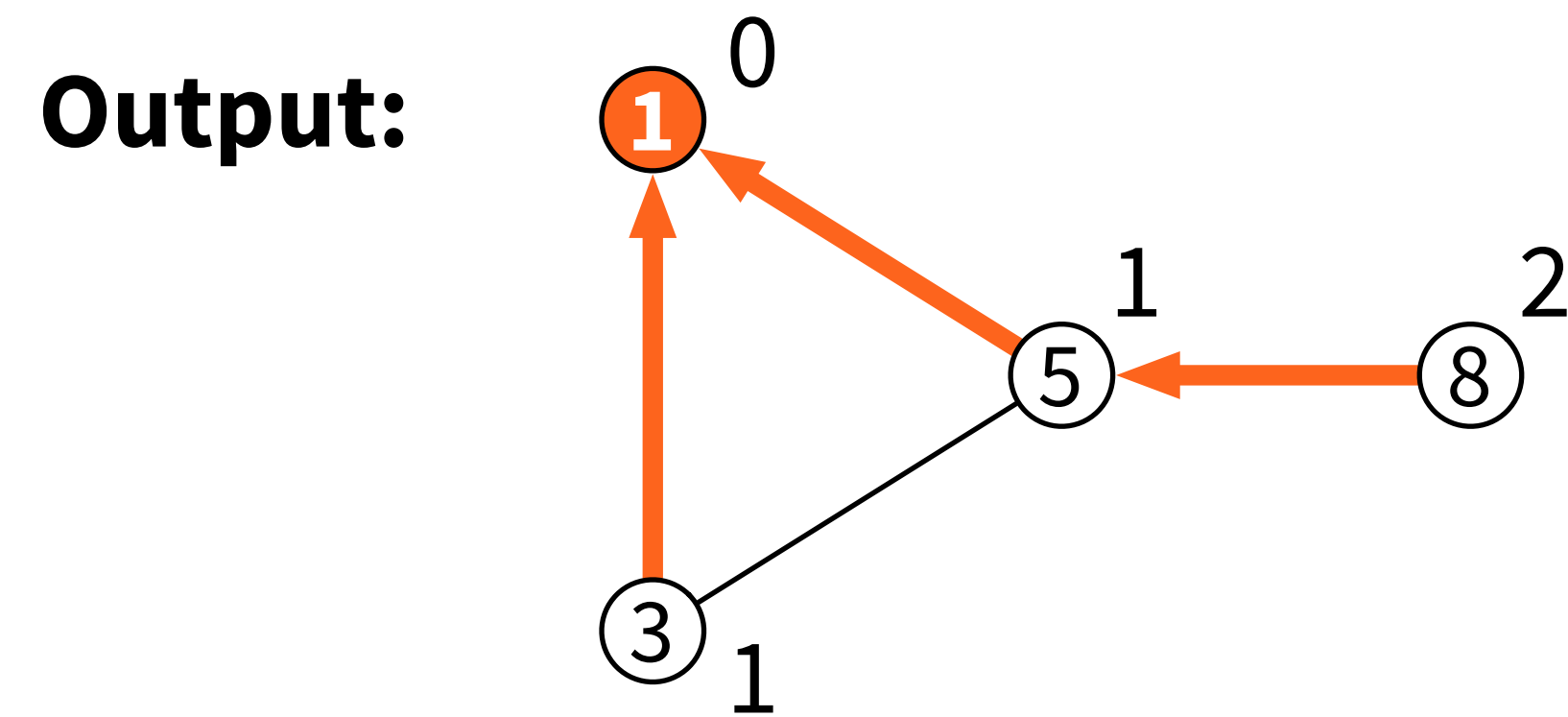
BFS tree

Input:



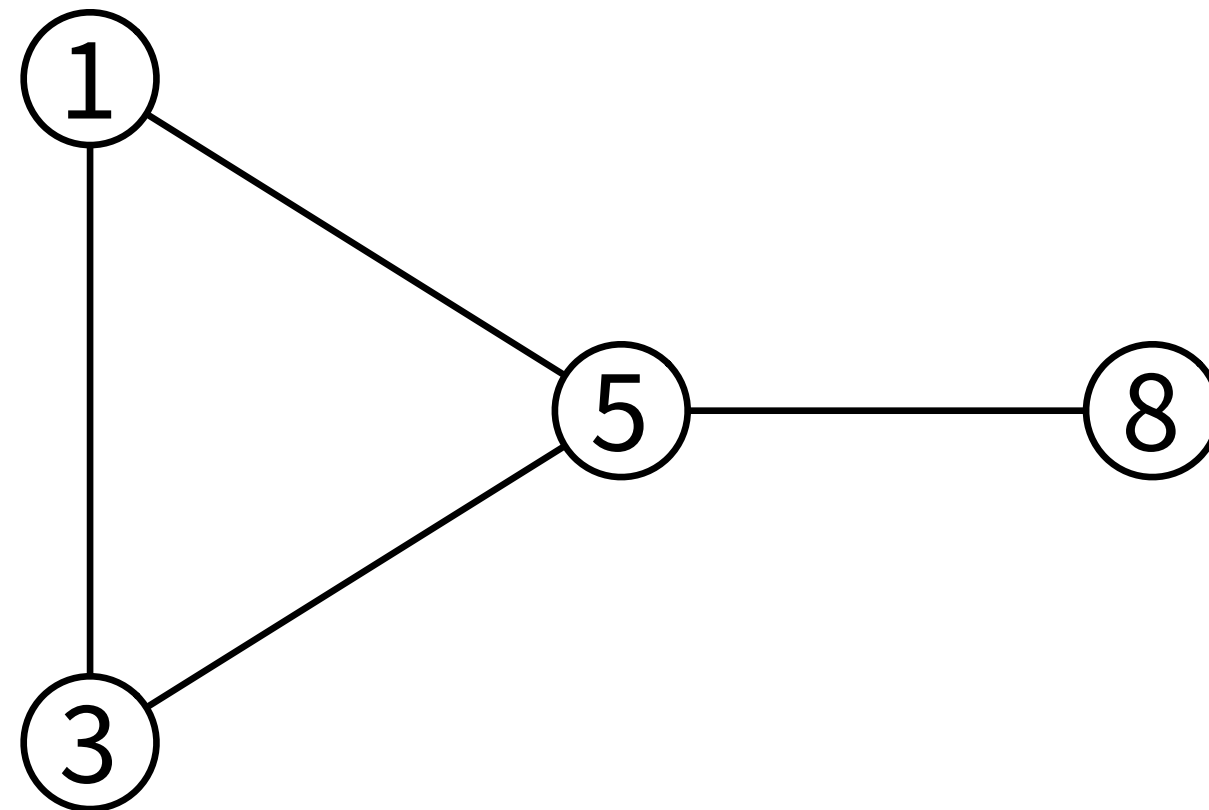
BFS tree

- Distances from s + shortest paths



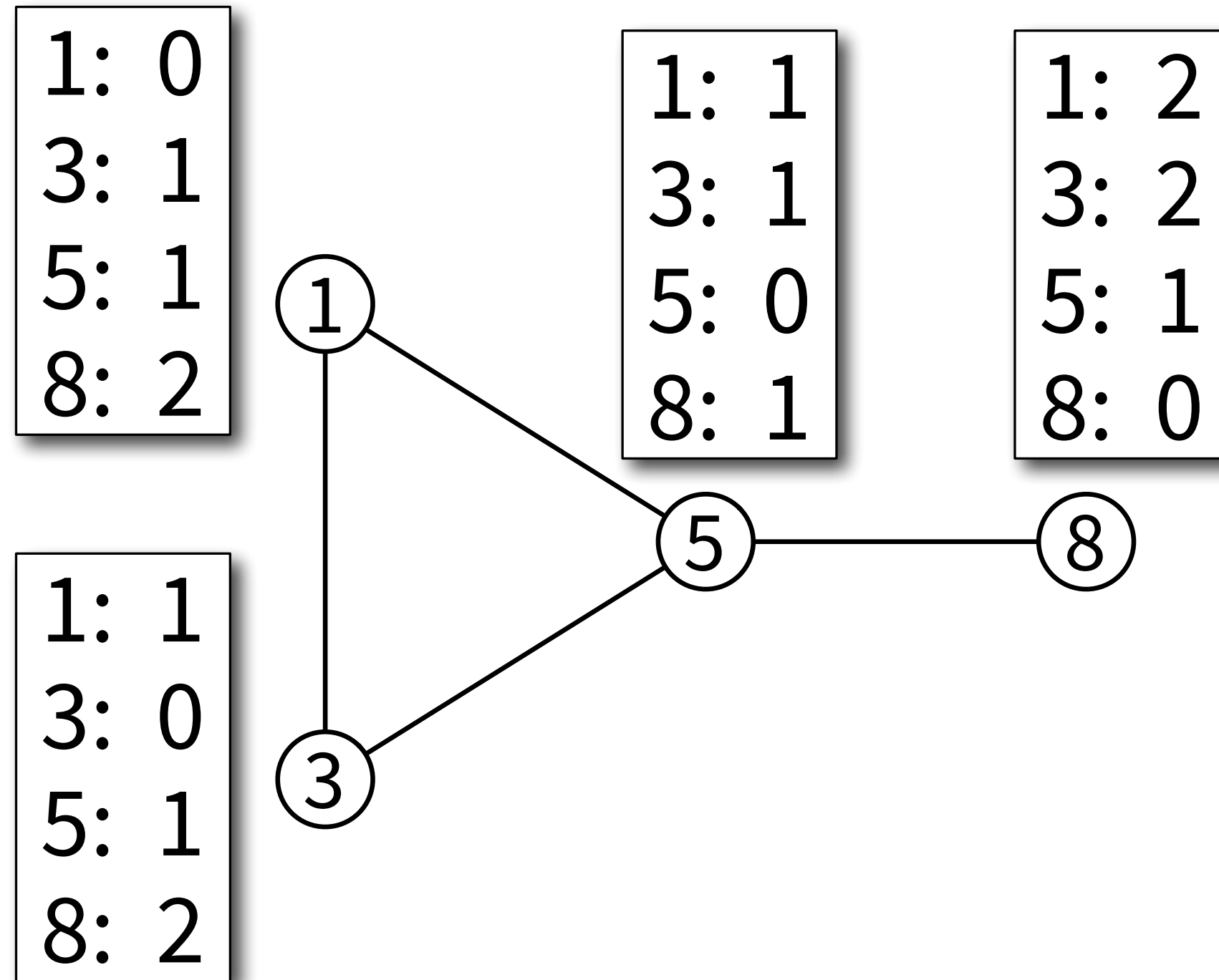
All-pairs shortest paths

Input:



All-pairs shortest paths

Output:

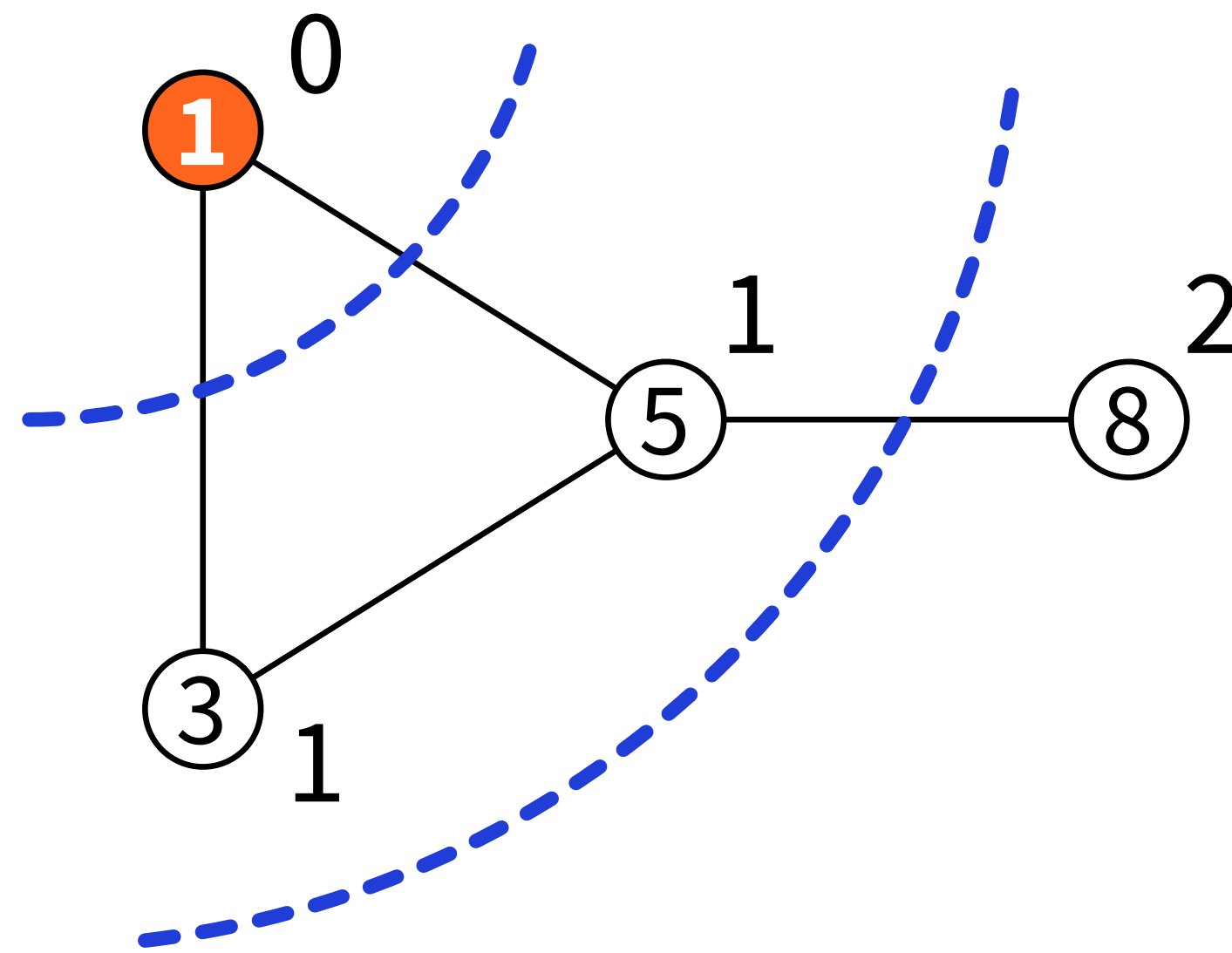


Algorithm Wave

- Solves **single-source shortest paths (SSSP)** in time $O(\text{diam}(G))$
- Leader/source sends a message “wave”, switches to state 0, stops
- Wave received in round t for the first time:
send “wave”, switch to state t , stop
- In time $O(\text{diam}(G))$ all nodes receive the wave

Algorithm Wave

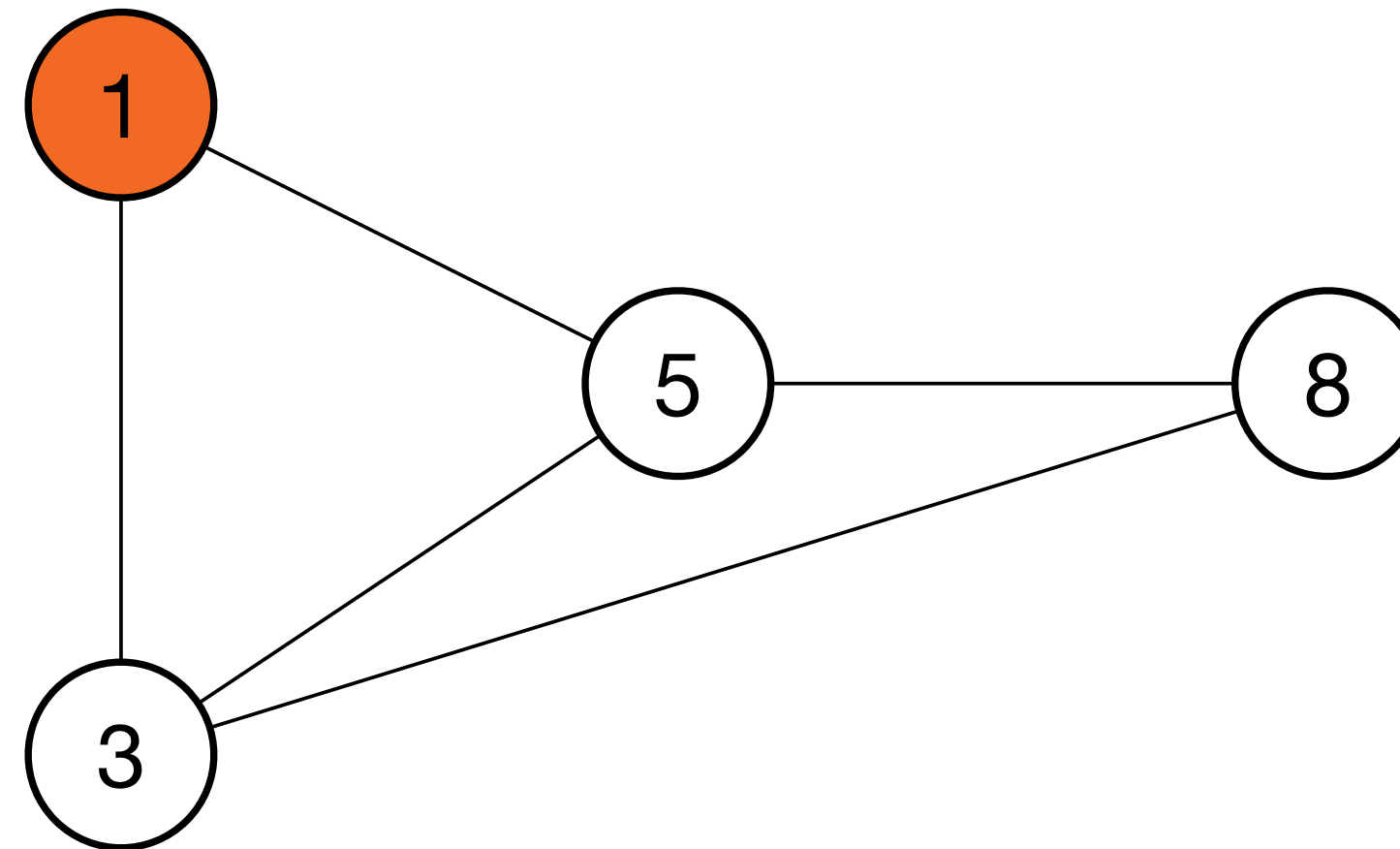
Output:



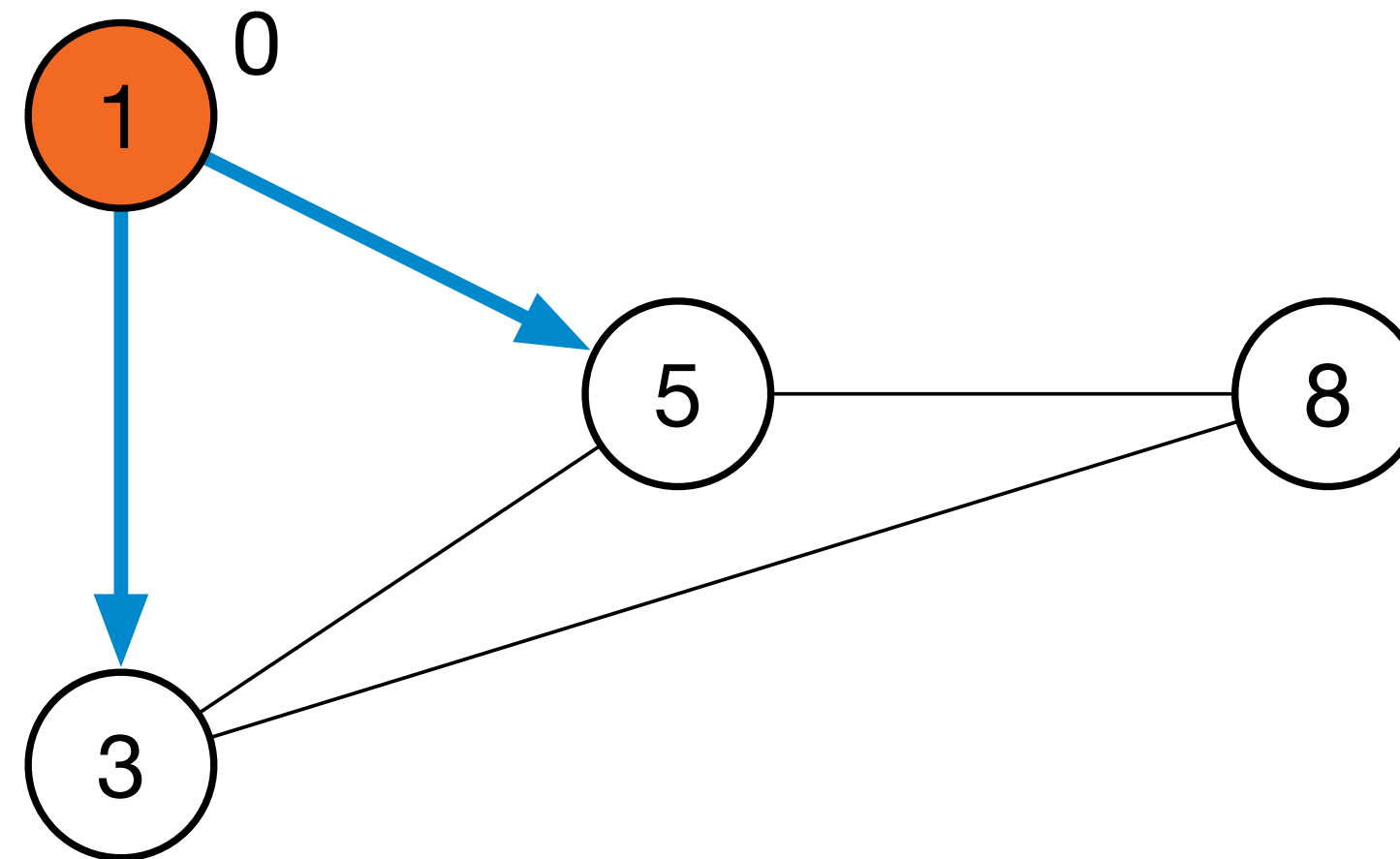
Algorithm BFS

- **Wave + handshakes**
- Tree construction:
 - “proposal” + “accept”
 - everyone knows their parent & children
- Acknowledgements back from leaf nodes

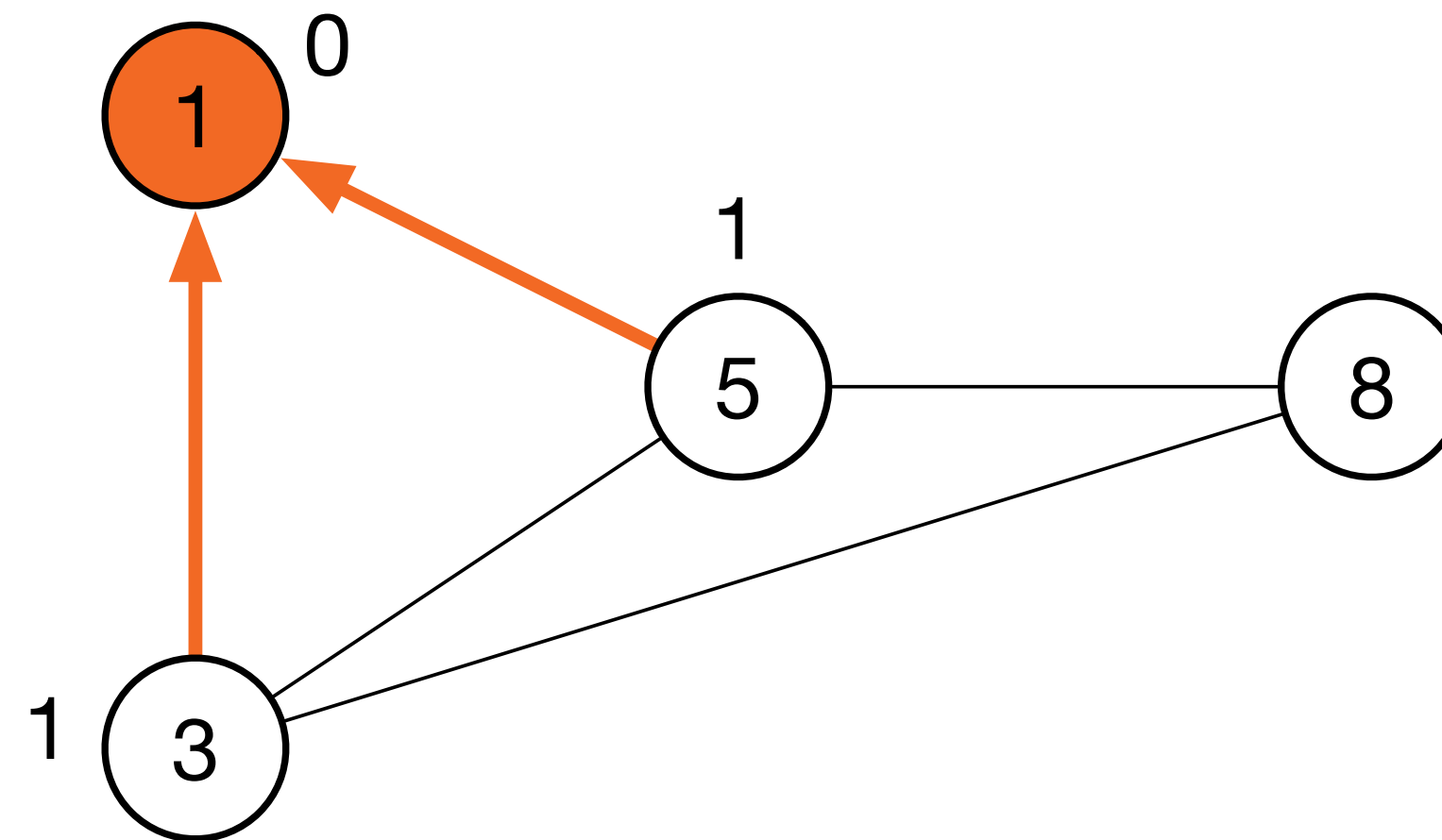
Algorithm BFS



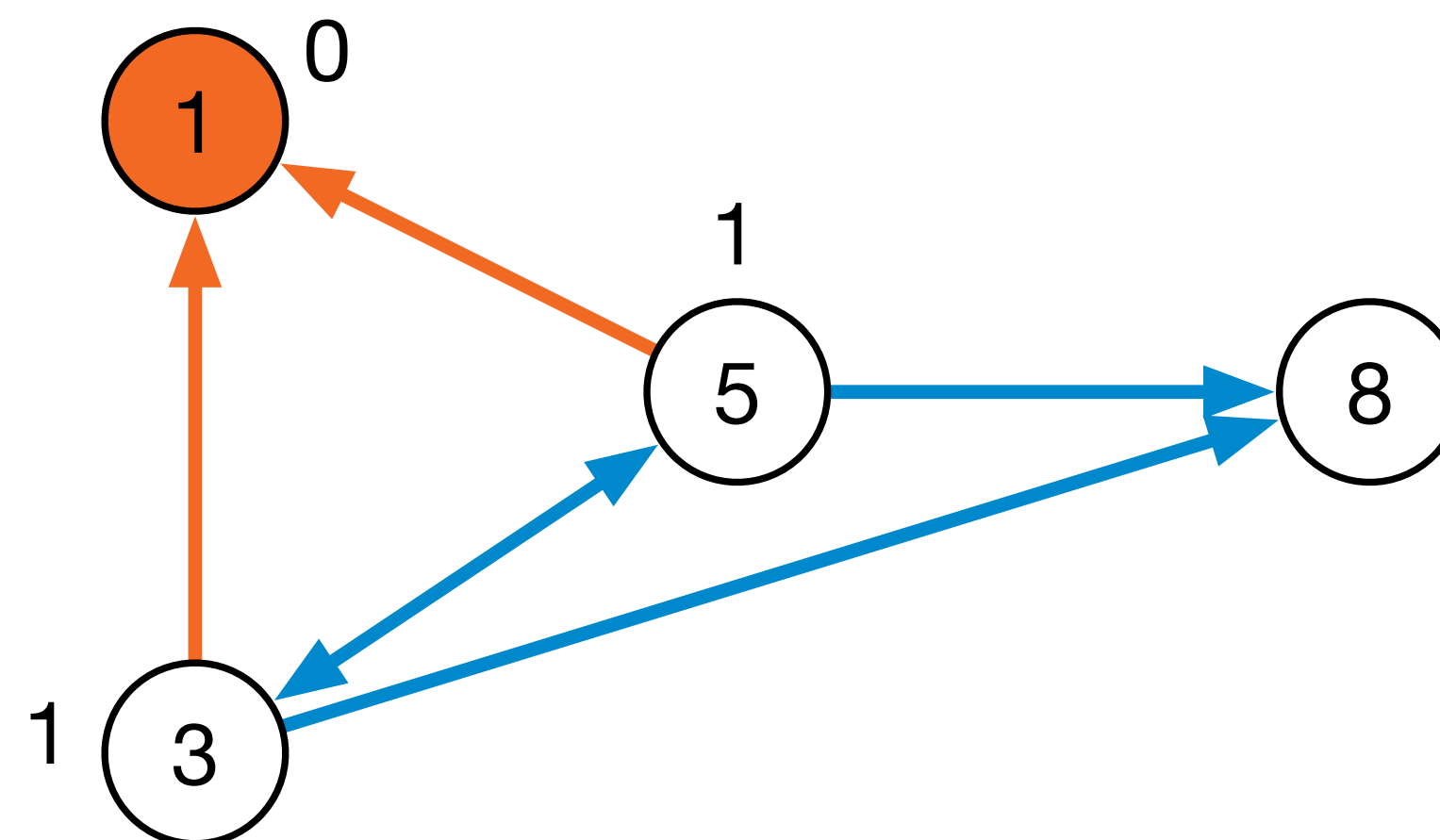
Algorithm BFS



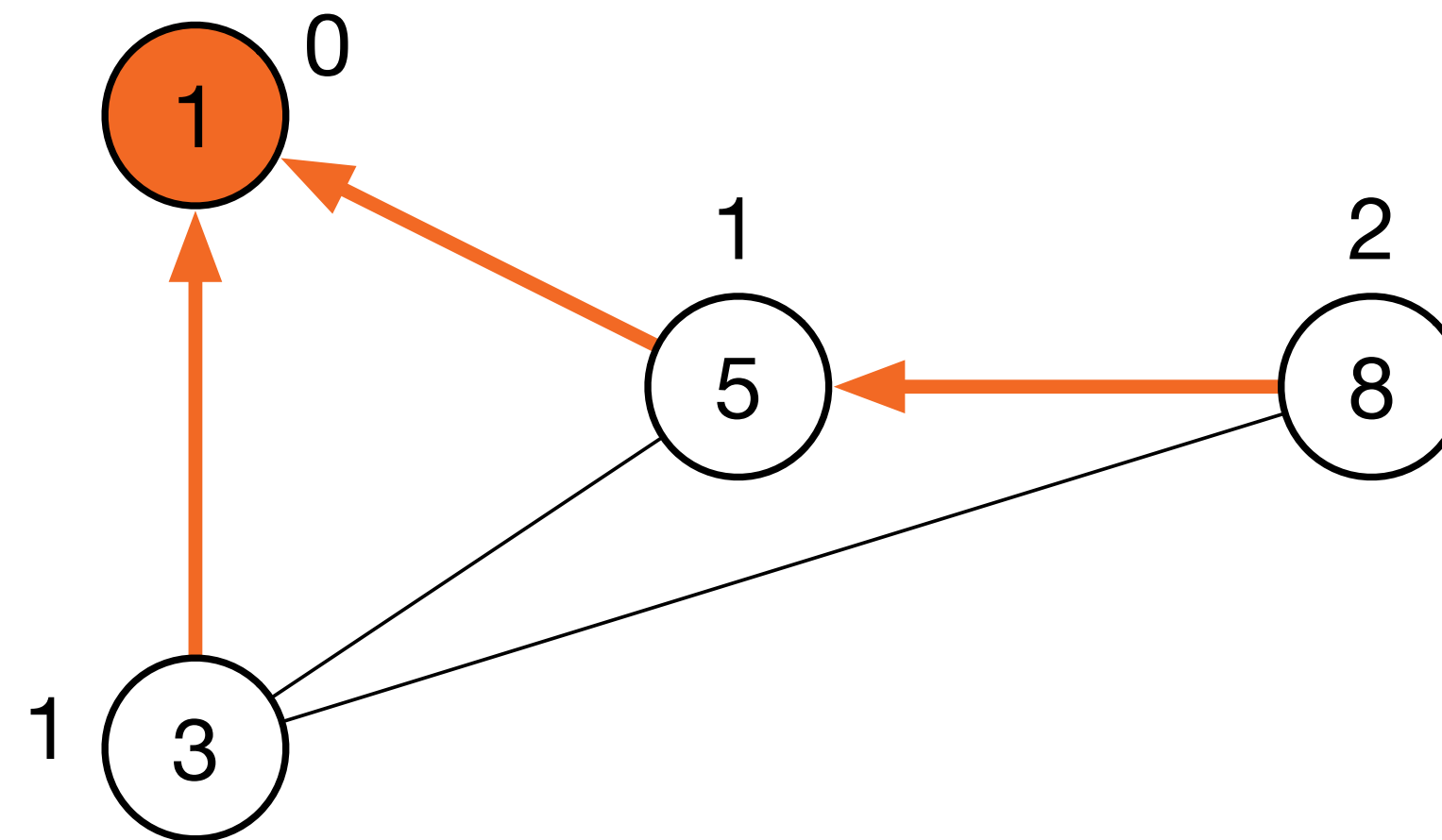
Algorithm BFS



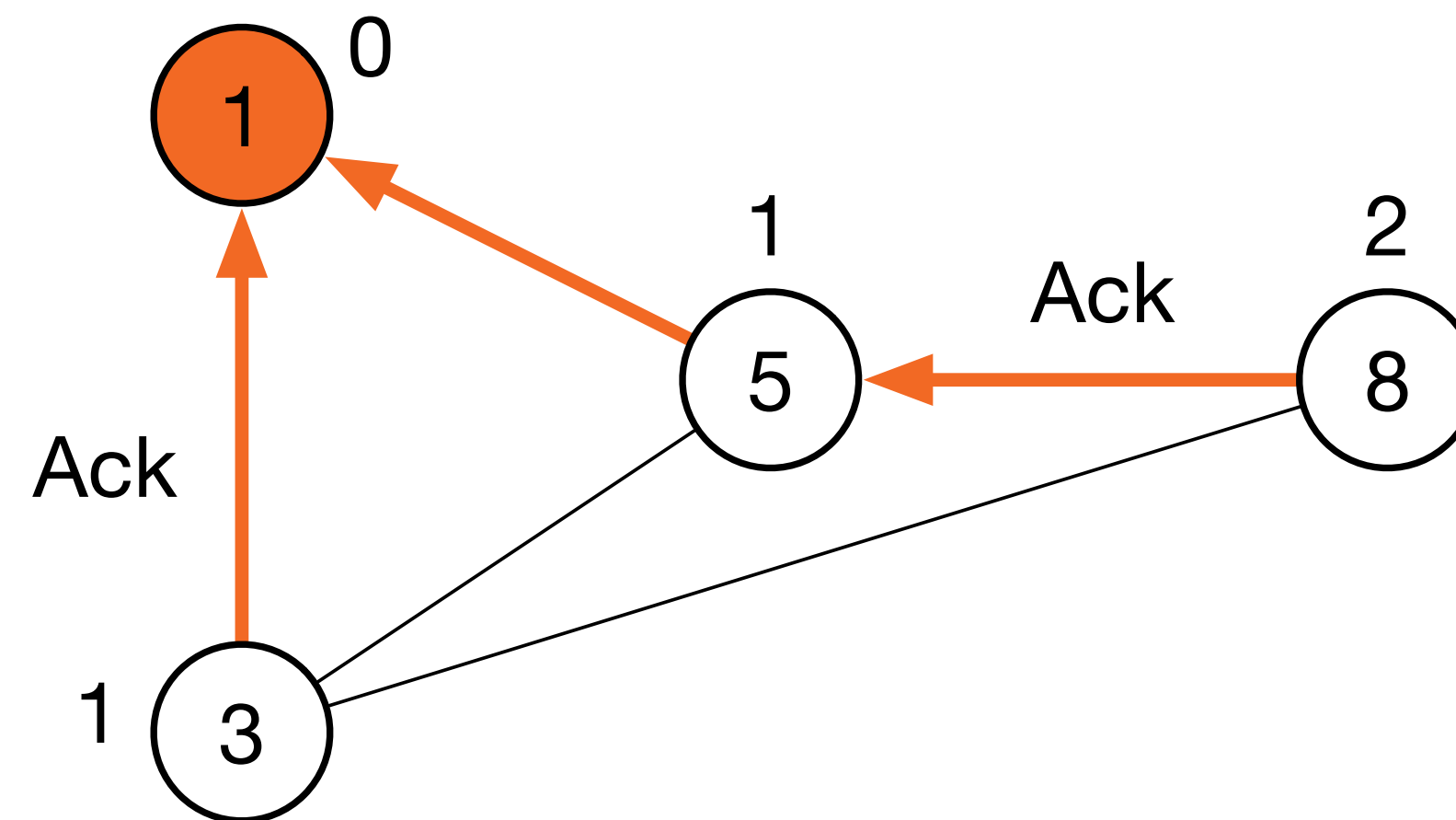
Algorithm BFS



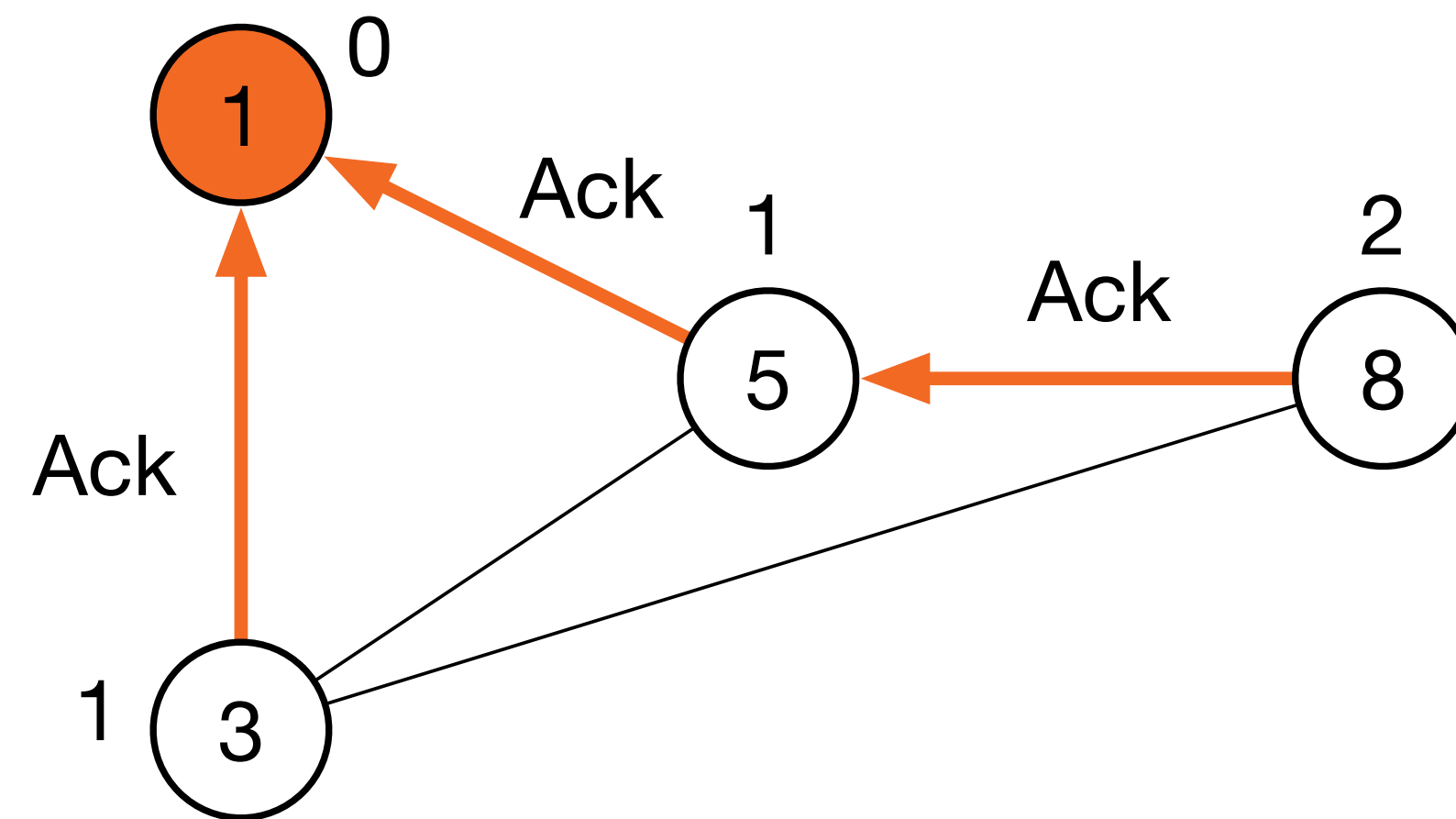
Algorithm BFS



Algorithm BFS



Algorithm BFS



Algorithm Leader

- Each node creates a separate BFS process
 - each node v pretends to be the root
 - messages of the BFS started by v contain $ID(v)$
- When two BFS processes “collide”, the one with the smaller root “wins”
 - each node only needs to send messages related to ***one BFS process***
- One tree wins everyone else → leader

Recap until now

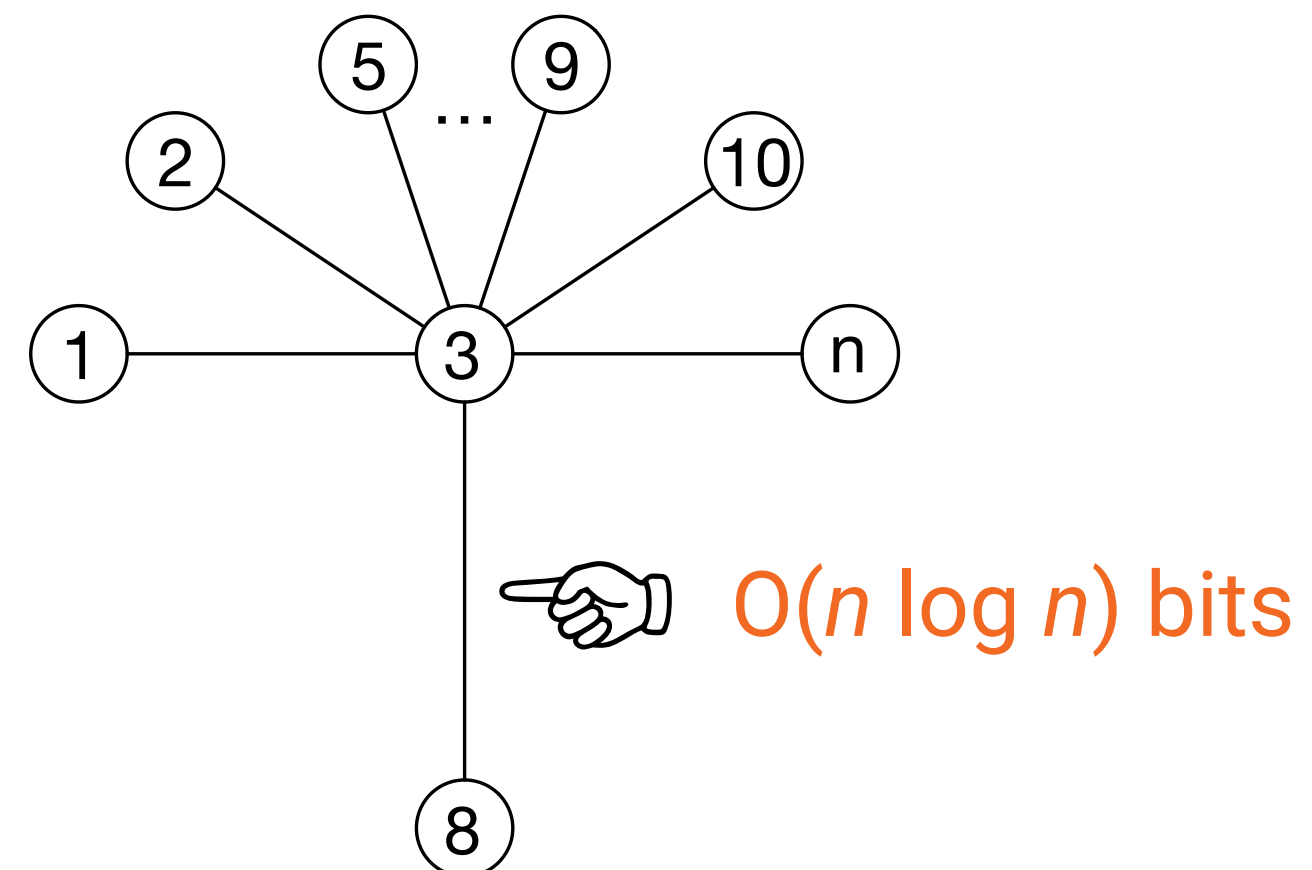
- **SSSP**: Wave algorithm
- **BFS tree**: Wave algorithm + acceptance/rejections
- **Leader election**: Many BFS in parallel
- All these problems can be solved in $O(\text{diam}(G))$ rounds in the **CONGEST** model

Algorithm APSP

- **Basic idea:** run Wave from each node
- **Challenge:** congestion

Algorithm APSP

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 - all waves parallel \rightarrow too many bits per edge



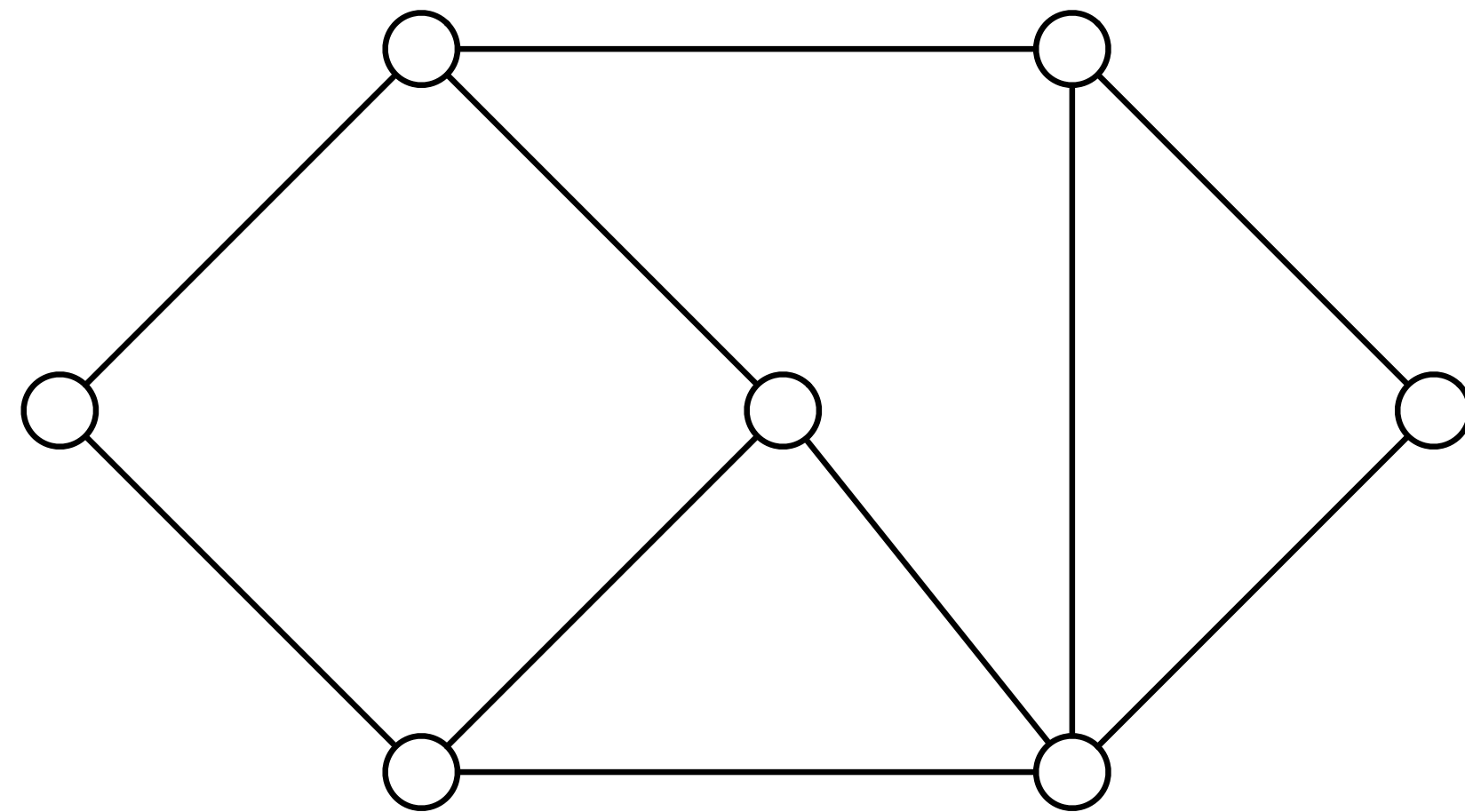
Algorithm APSP

- **Basic idea:** run Wave from each node
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 - all waves parallel → too many bits per edge
 - all waves sequentially → takes too long

Algorithm APSP

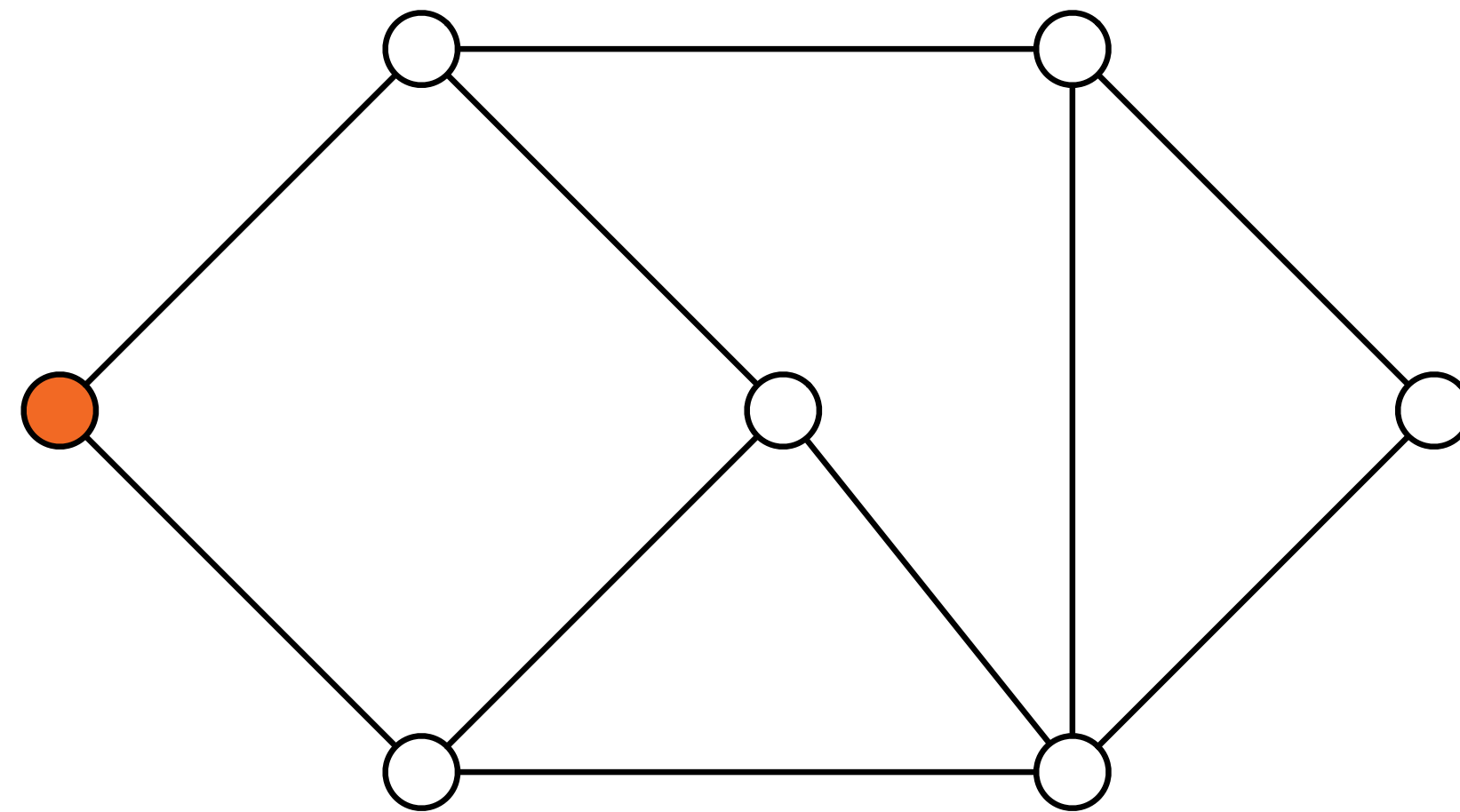
- **Basic idea:** run Wave from each node
- **Challenge:** congestion
 - all waves parallel → too many bits per edge
 - all waves sequentially → takes too long
- **Solution:** **pipelining**
 - all waves in parallel in such a way that each node propagates at most one wave per round

Algorithm APSP



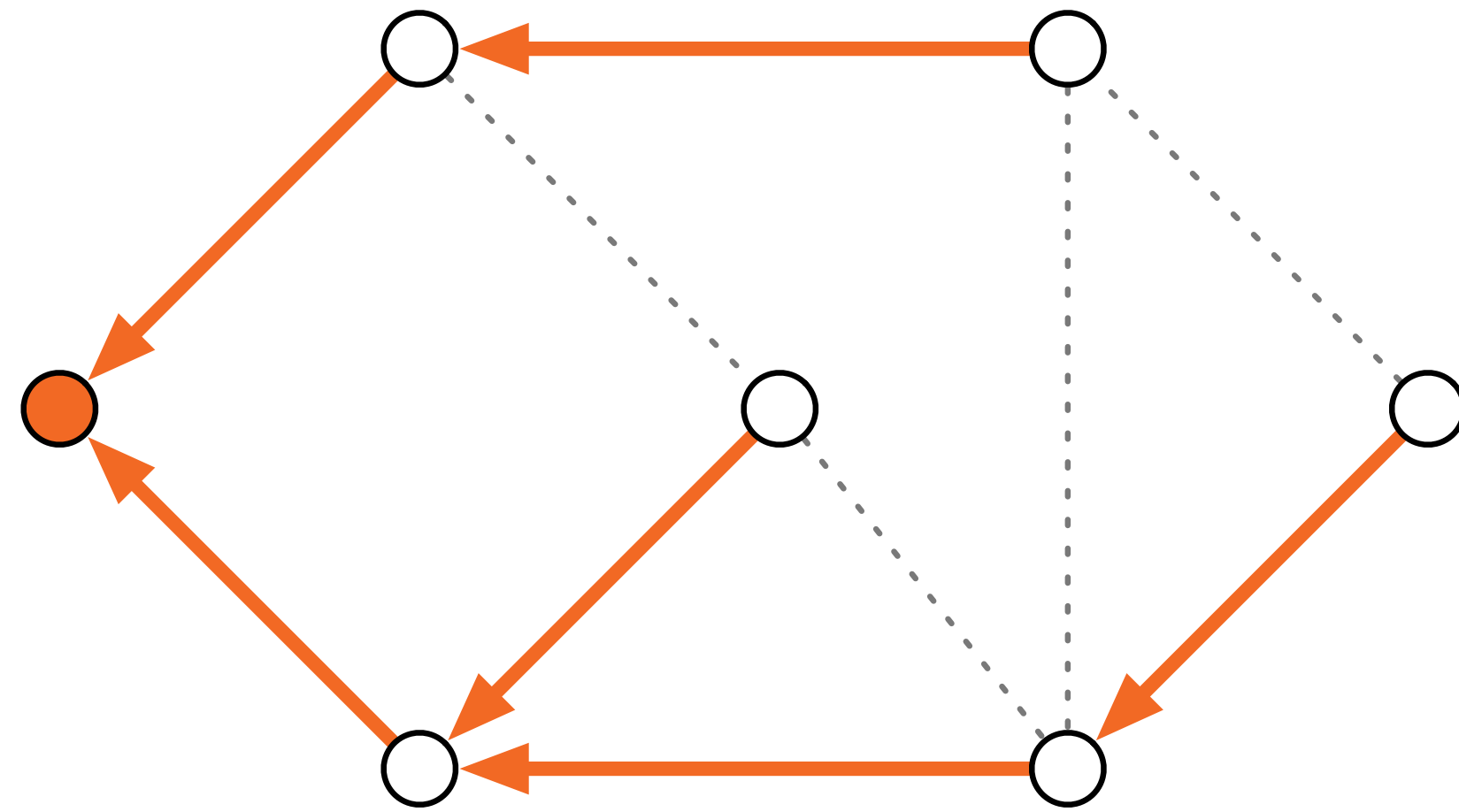
Algorithm APSP

- Elect leader



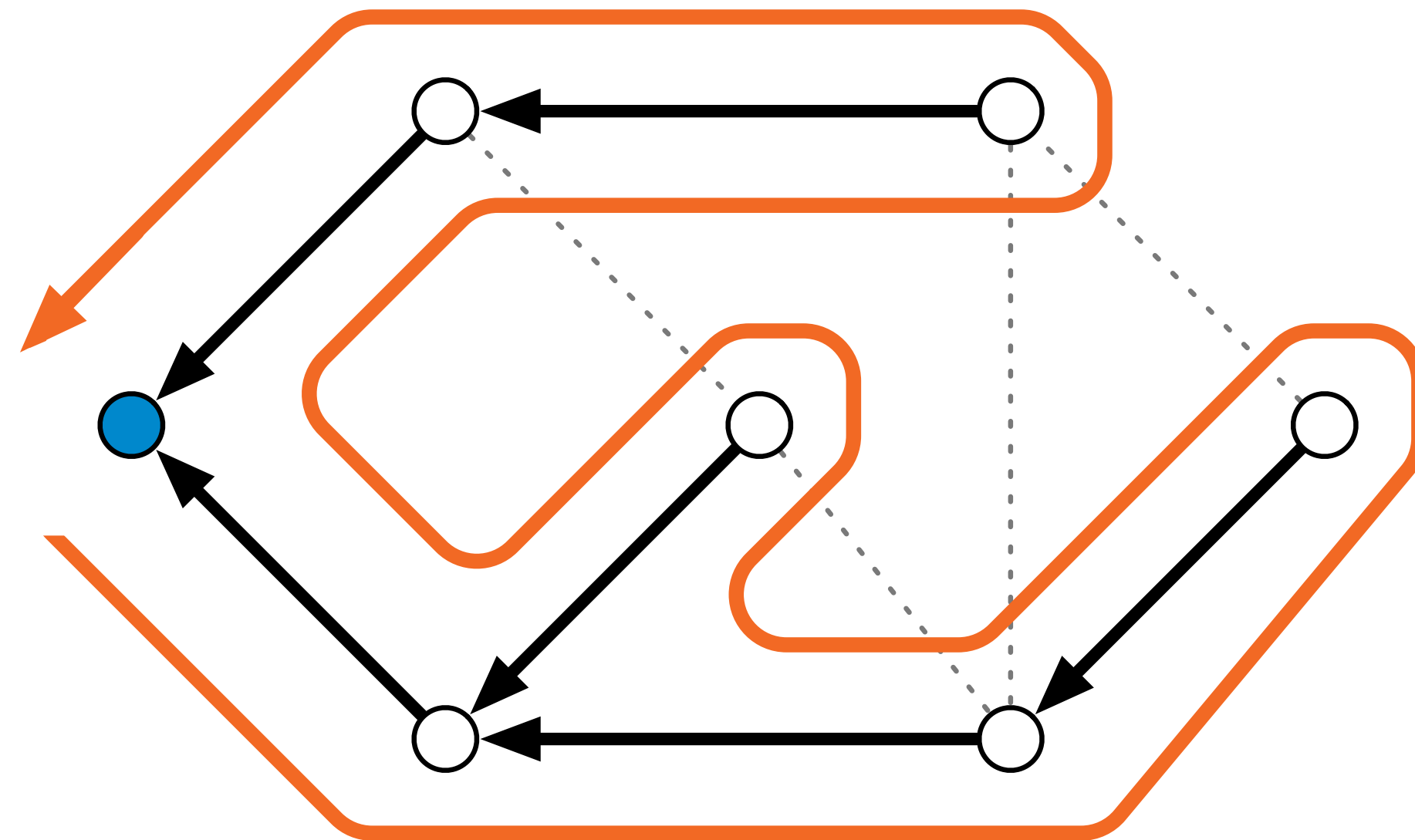
Algorithm APSP

- Elect leader, construct BFS tree



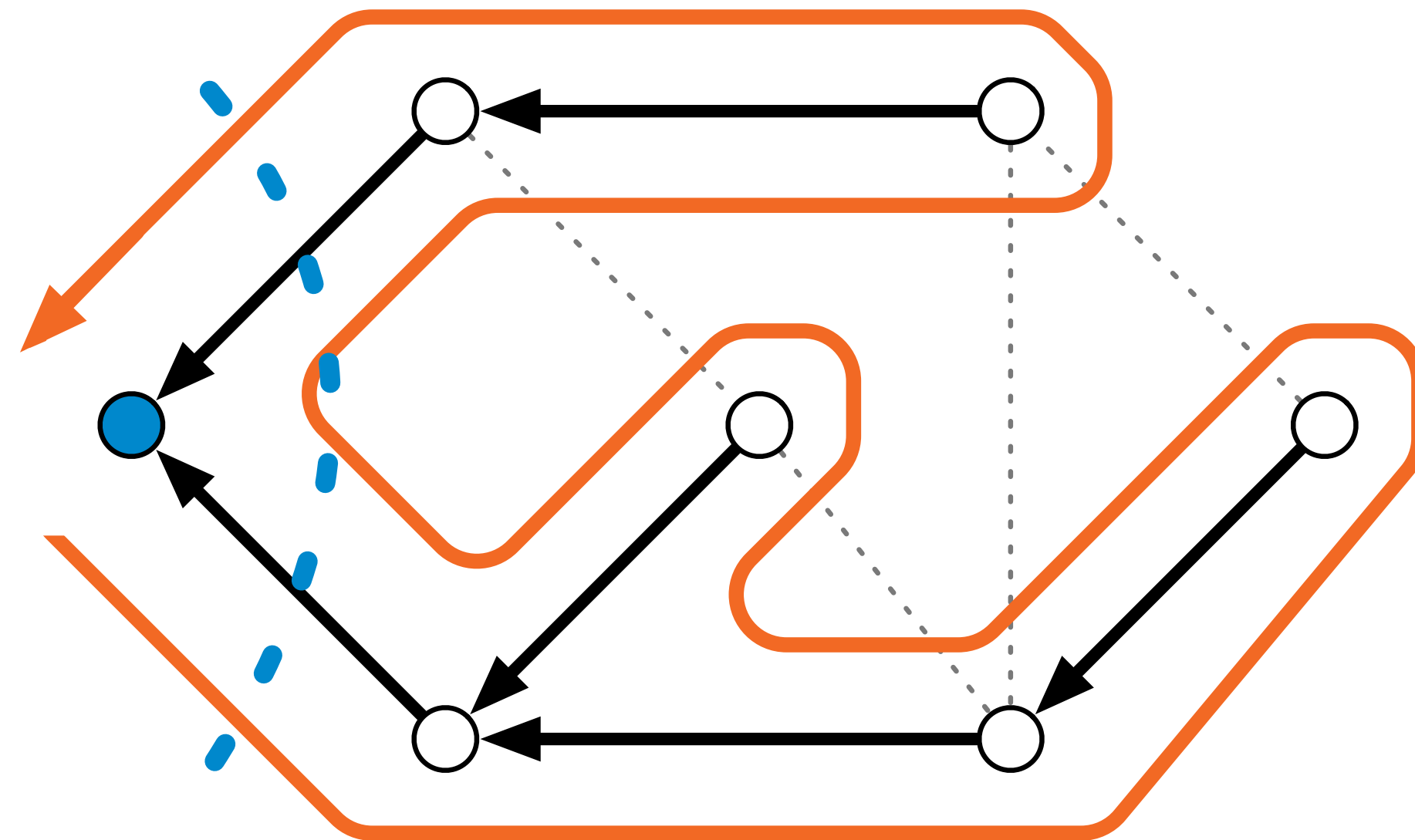
Algorithm APSP

- Move token along BFS tree slowly (every 2 rounds)



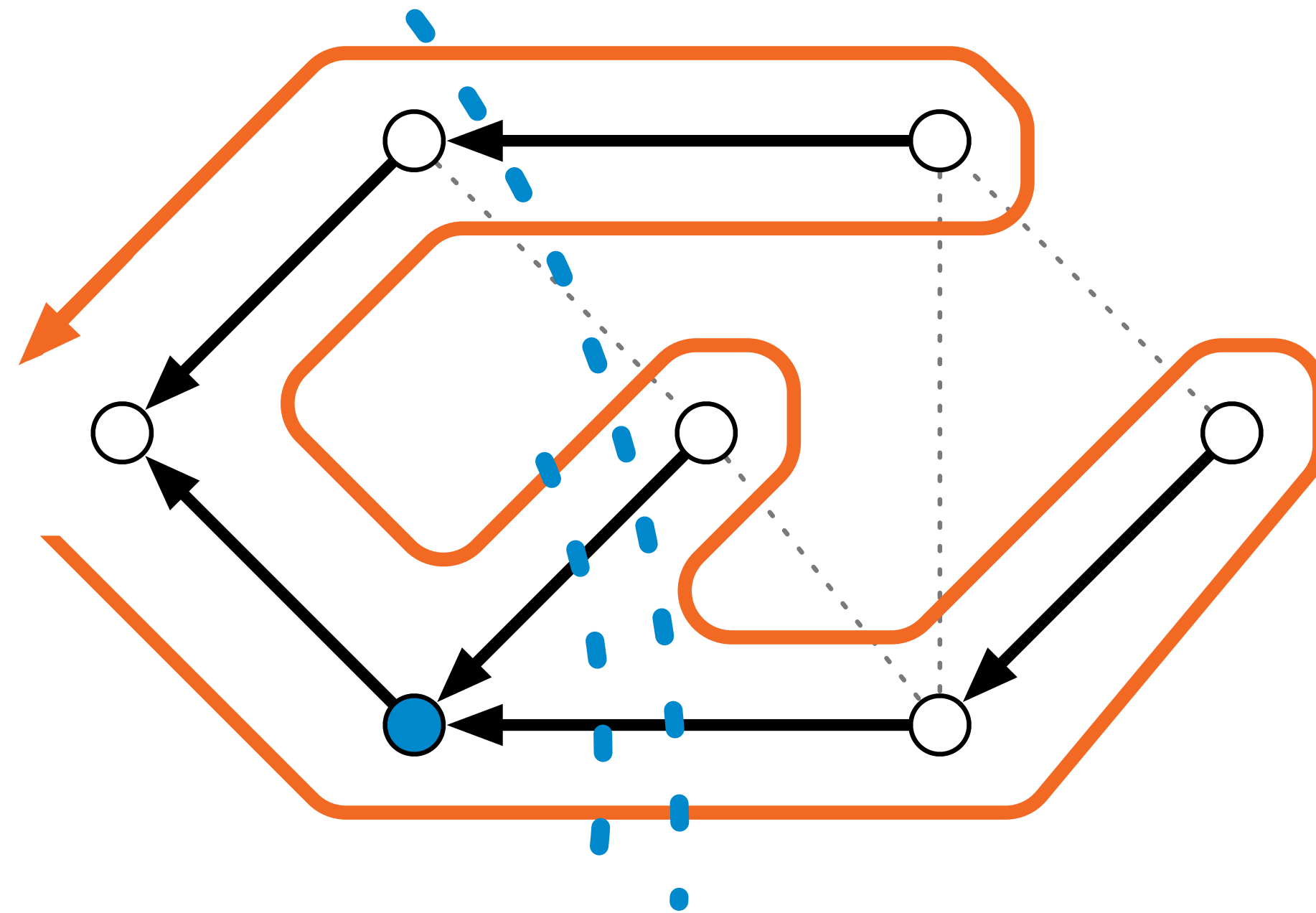
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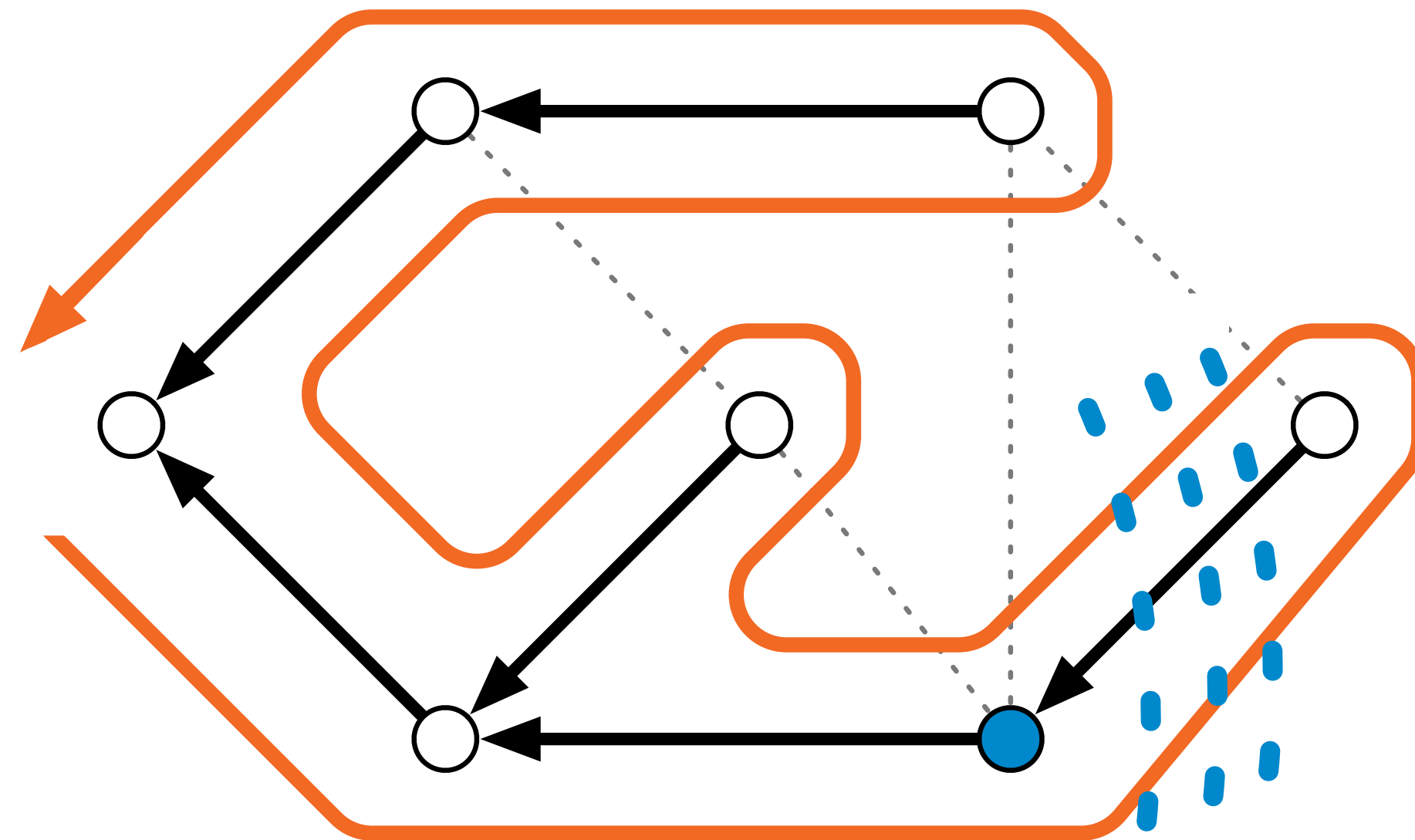
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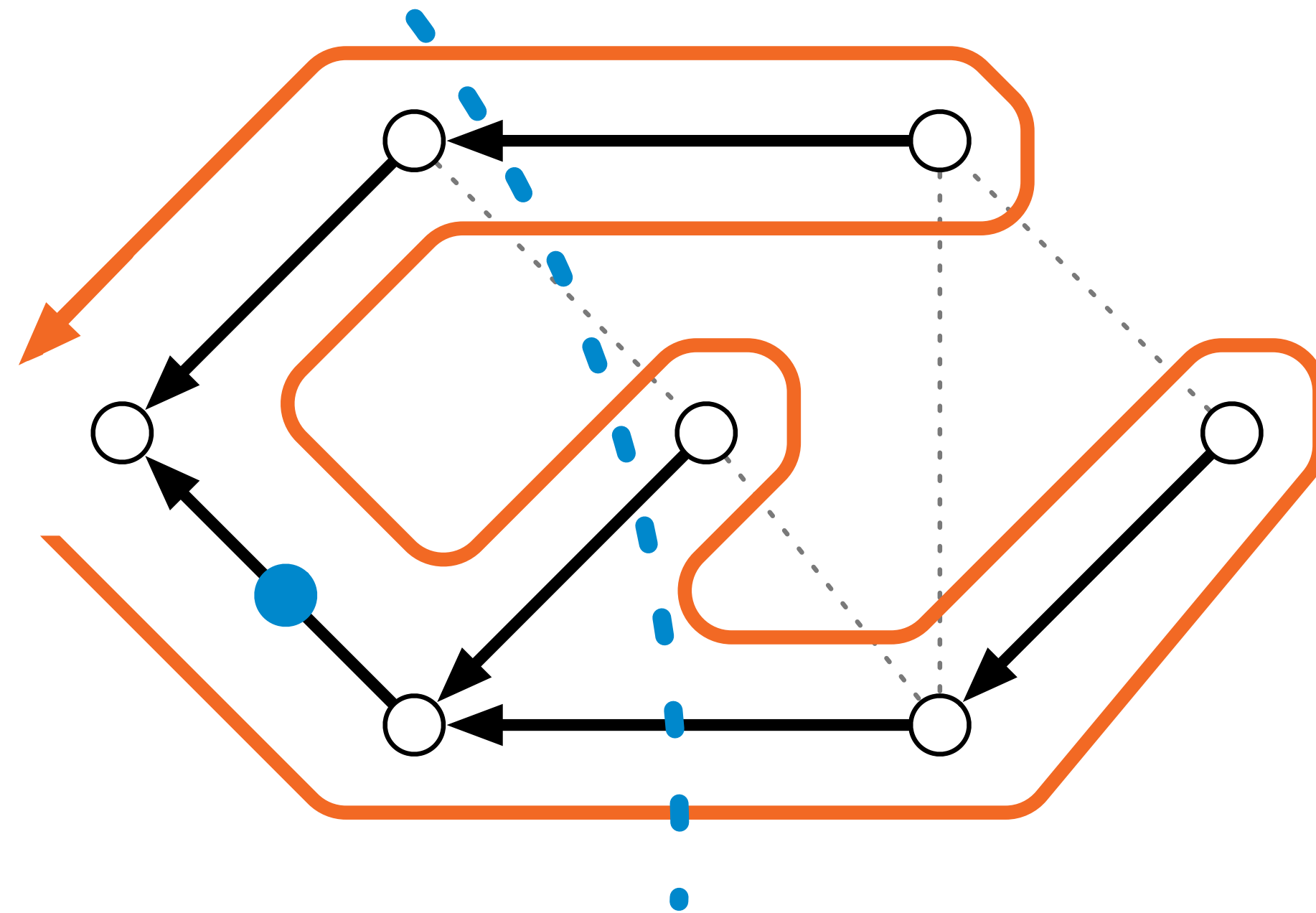
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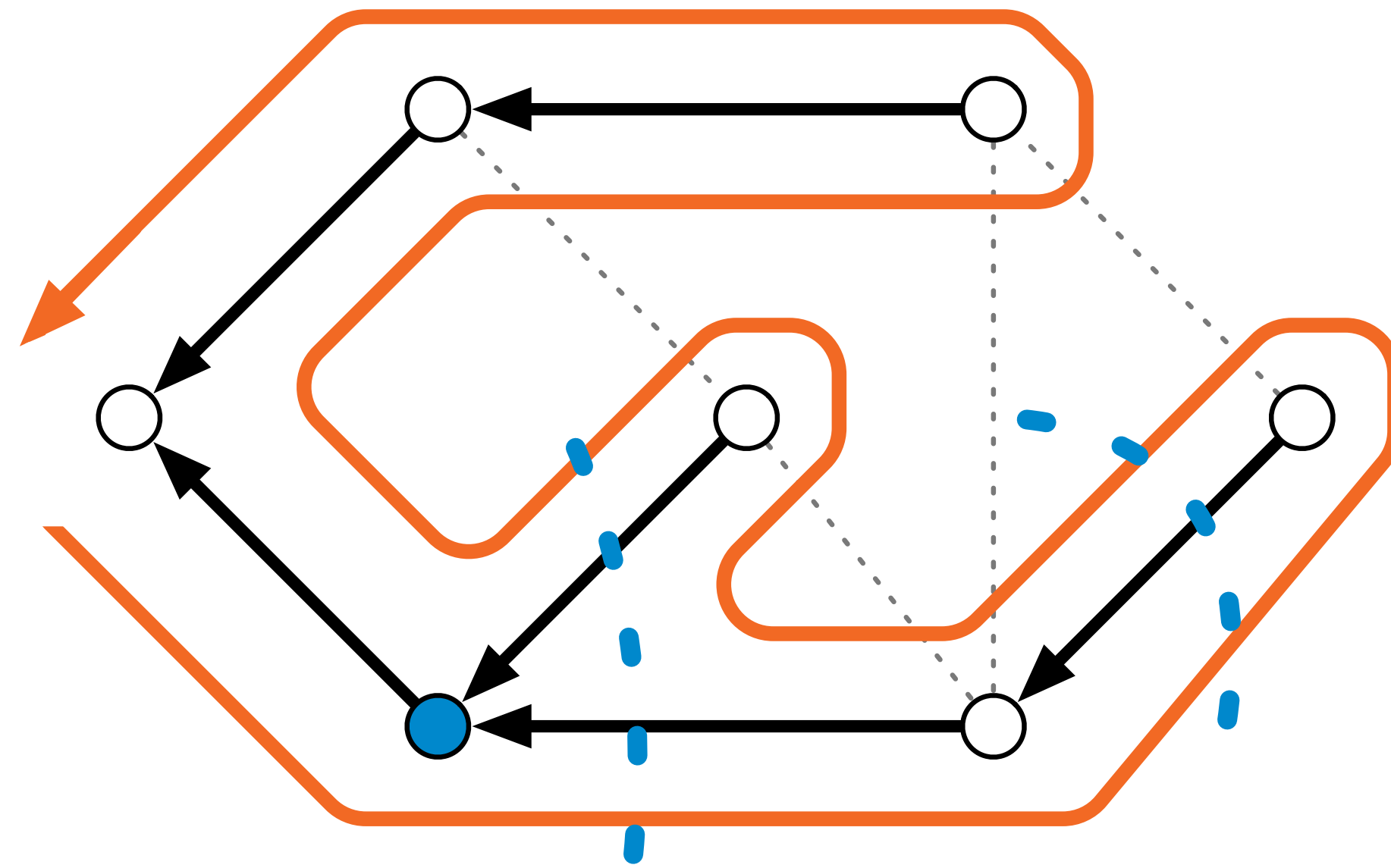
Algorithm APSP

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Algorithm APSP

- Move token along BFS tree slowly (every 2 rounds)



Algorithm APSP - Animation

- See animation here: <https://jukkasuomela.fi/apsp/>

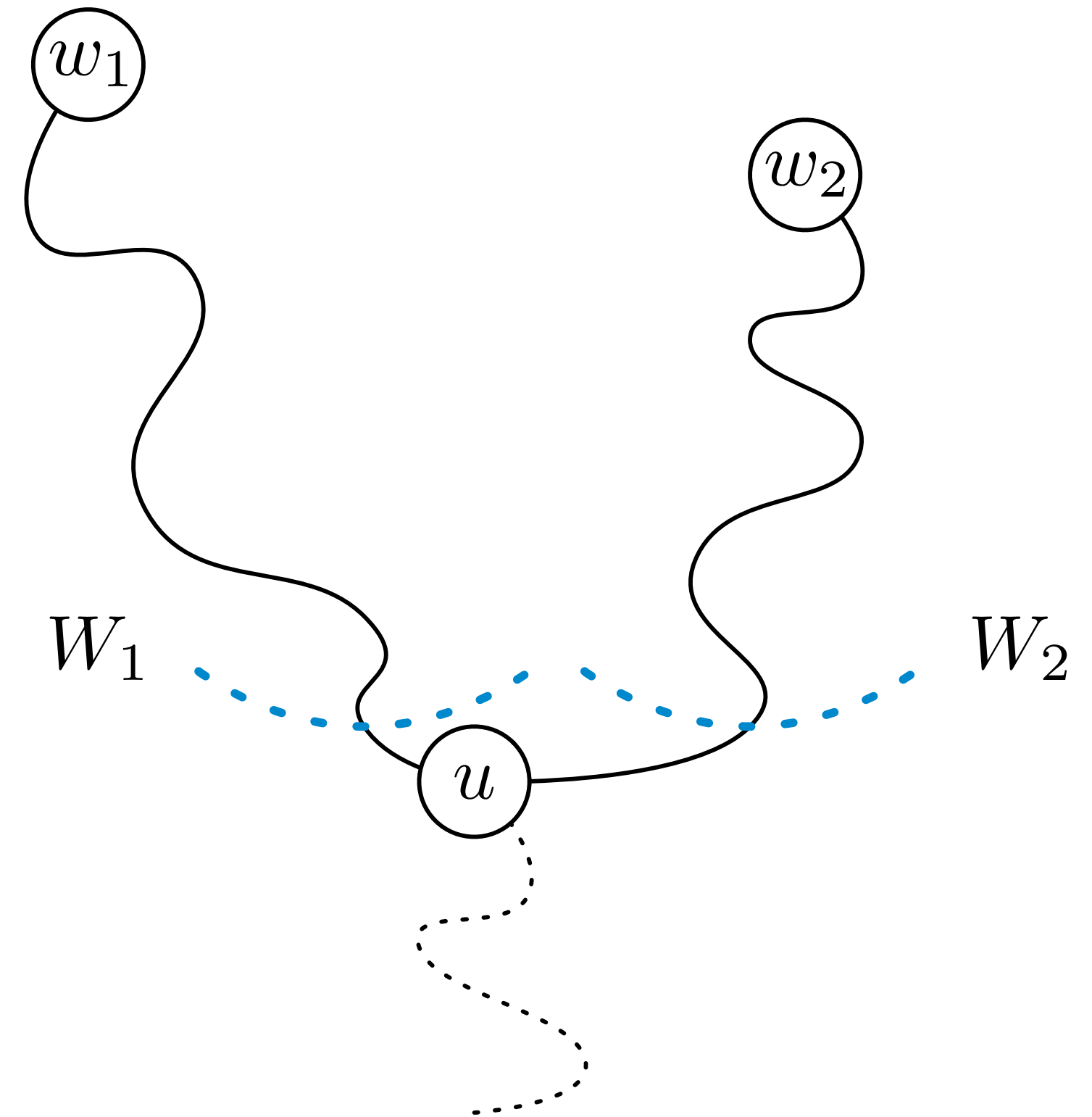
Algorithm APSP: no collisions

Claim: **at each round**, a node has to propagate **only one wave**

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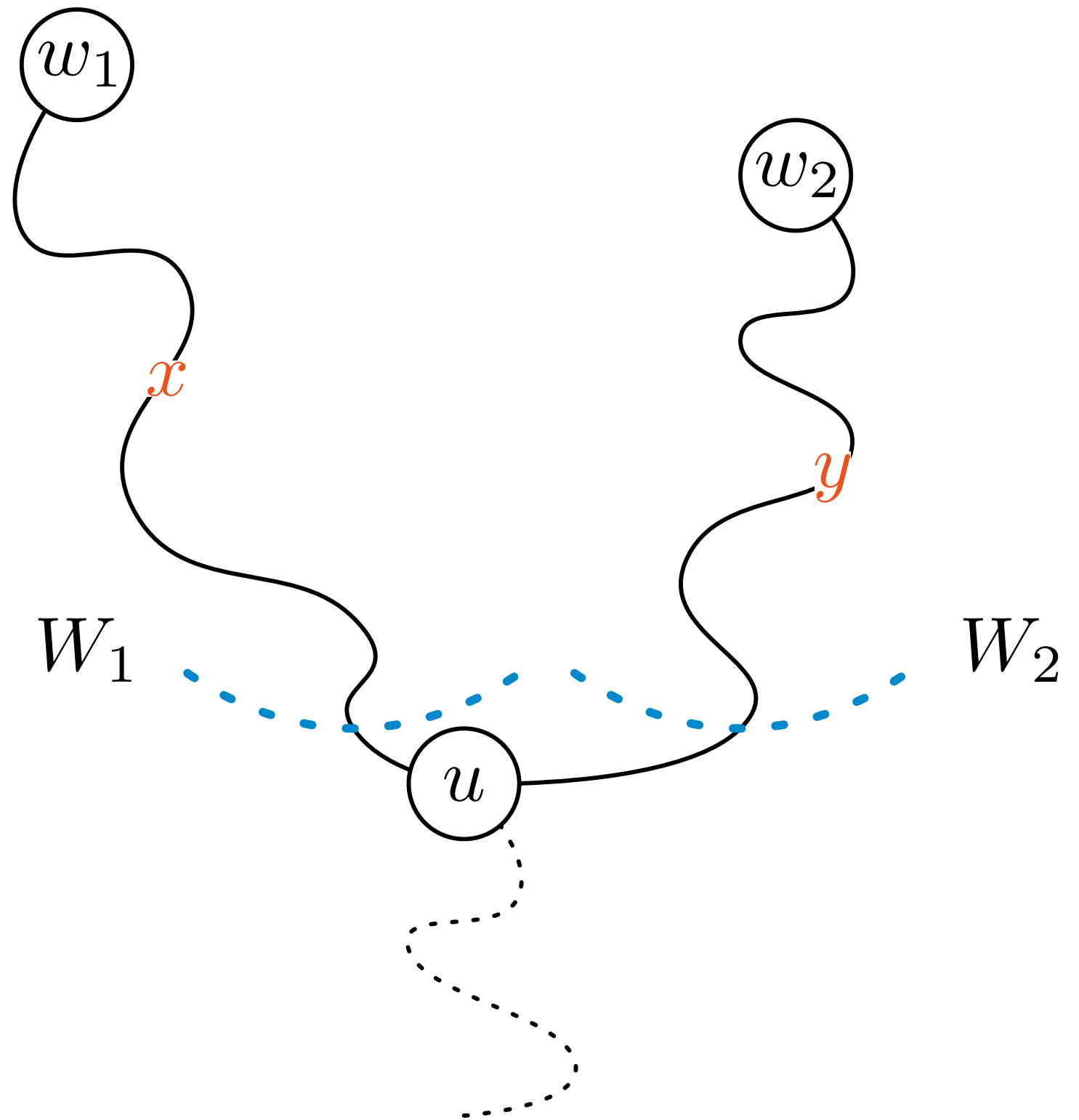
- Suppose, for a contradiction that this is not the case



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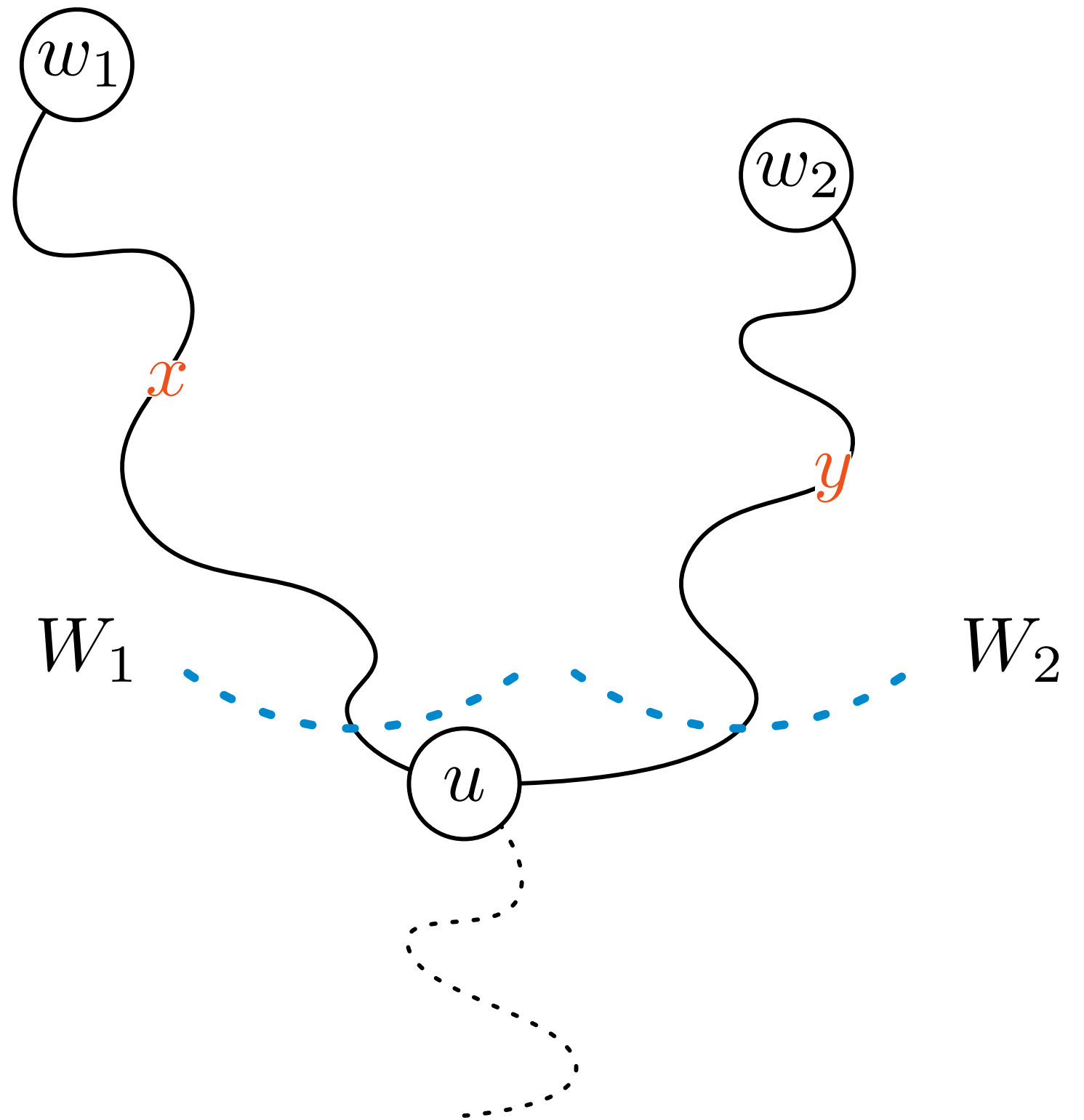
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- $d(u, w_1) = x, d(u, w_2) = y$



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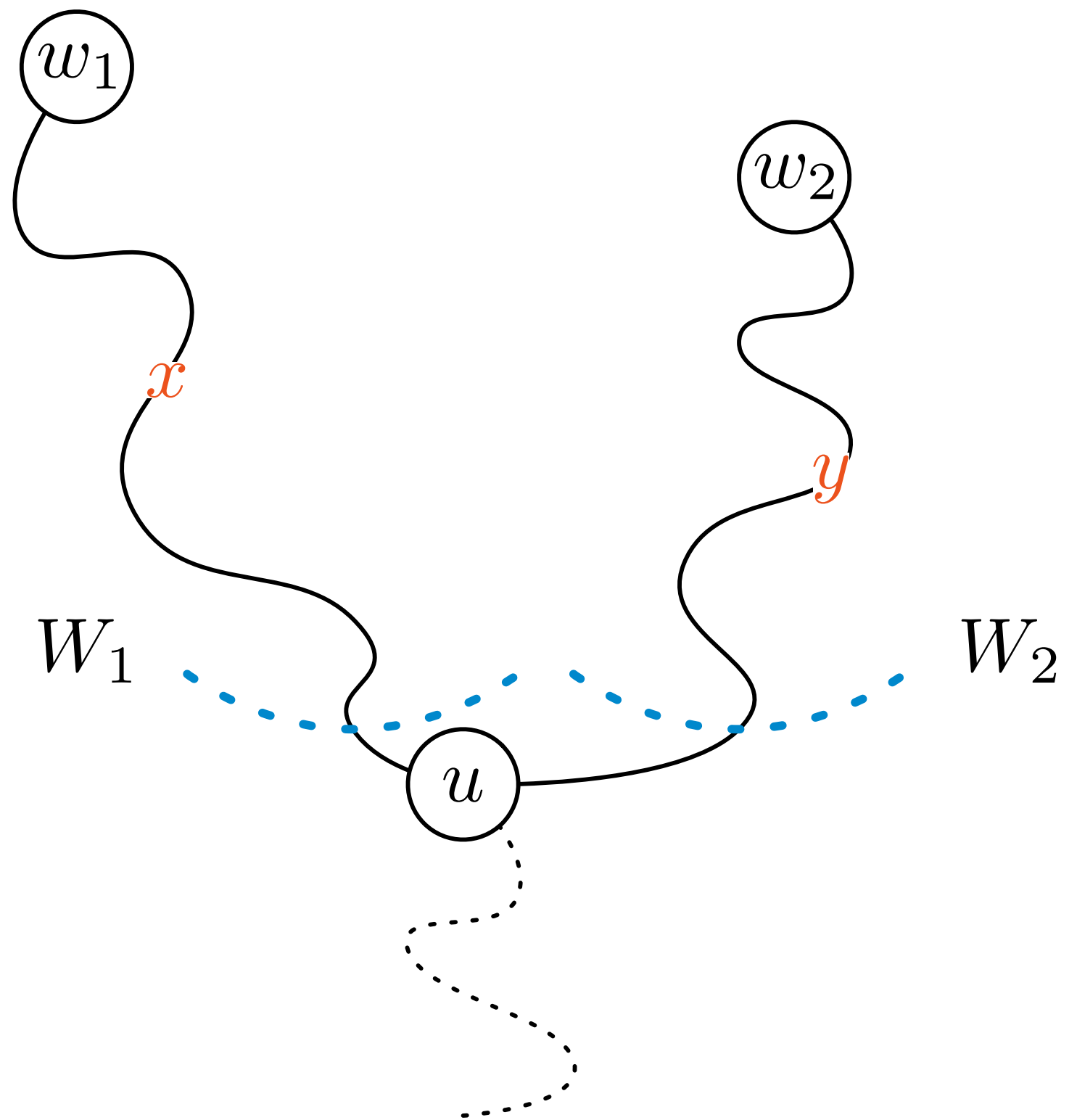
- Suppose, for a contradiction that this is not the case
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- $x \neq y$



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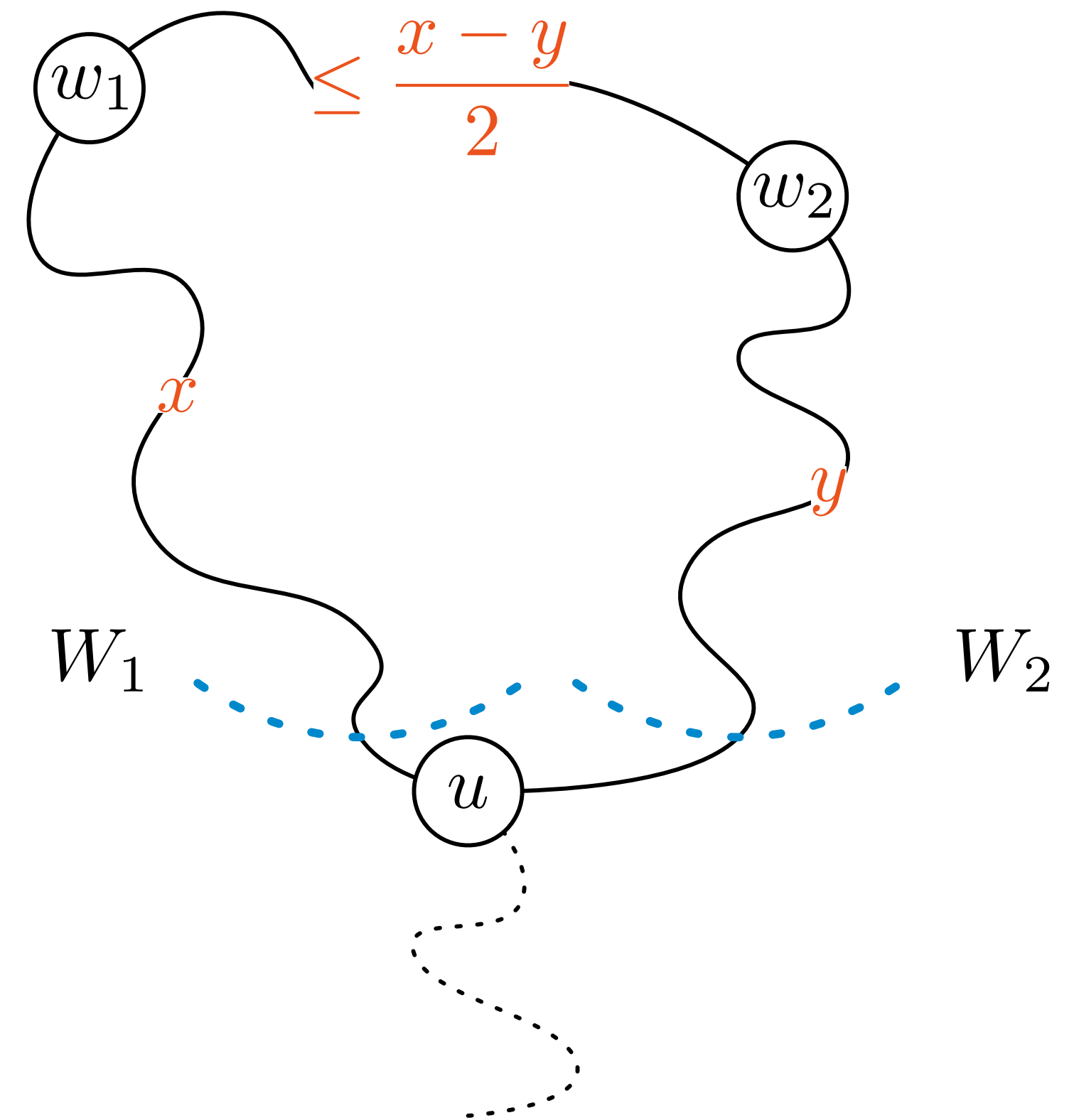
- Suppose, for a contradiction that this is not the case
- $d(u, w_1) = x, d(u, w_2) = y$
- $x \neq y$
- w.l.o.g. $x > y$



Algorithm APSP: no collisions

Claim: at each round, a node has to propagate **only one wave**

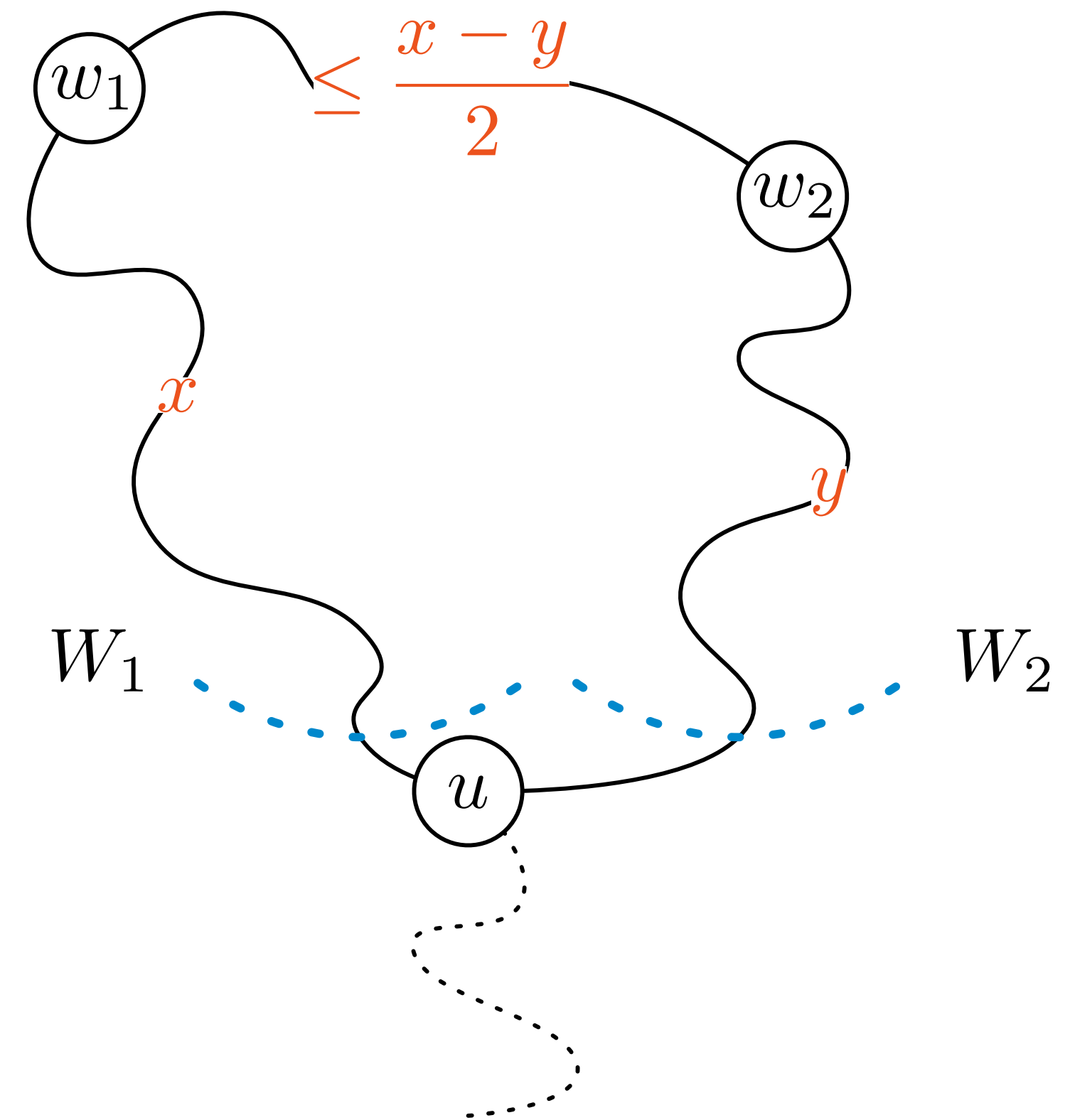
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- $d(w_1, w_2) \leq (x - y) / 2$



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- Suppose, for a contradiction that this is not the case
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- w.l.o.g. $x > y$
- $d(w_1, w_2) \leq (x - y) / 2$
- $x \leq y + (x - y) / 2$



Algorithm APSP: no collisions

Claim: at each round, a node has to propagate **only one wave**

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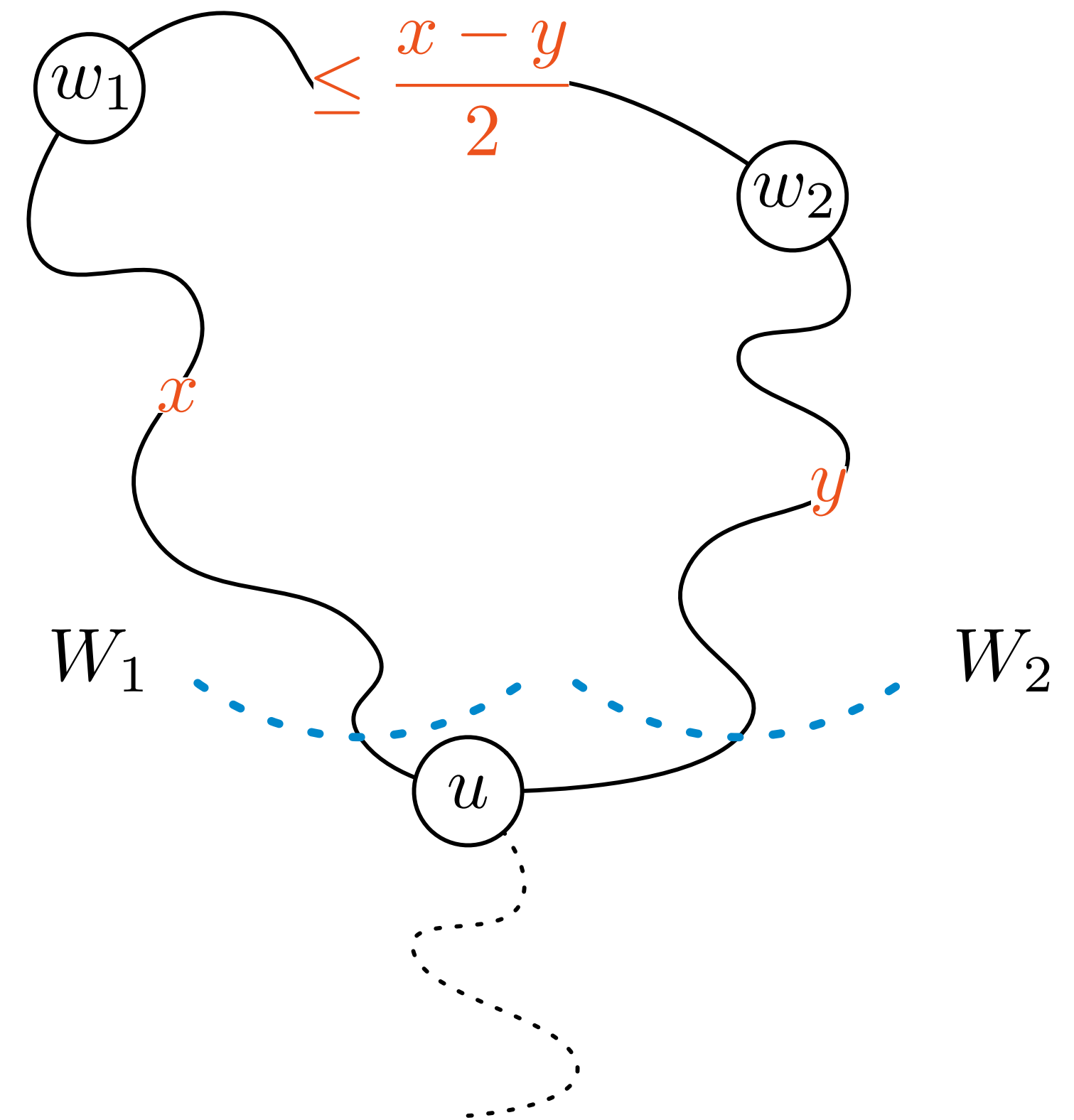
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- $x \leq y + (x - y) / 2$

- $x \leq y$



Algorithm APSP: no collisions

Claim: at each round, a node has to propagate **only one wave**

- Suppose, for a contradiction that this is not the case

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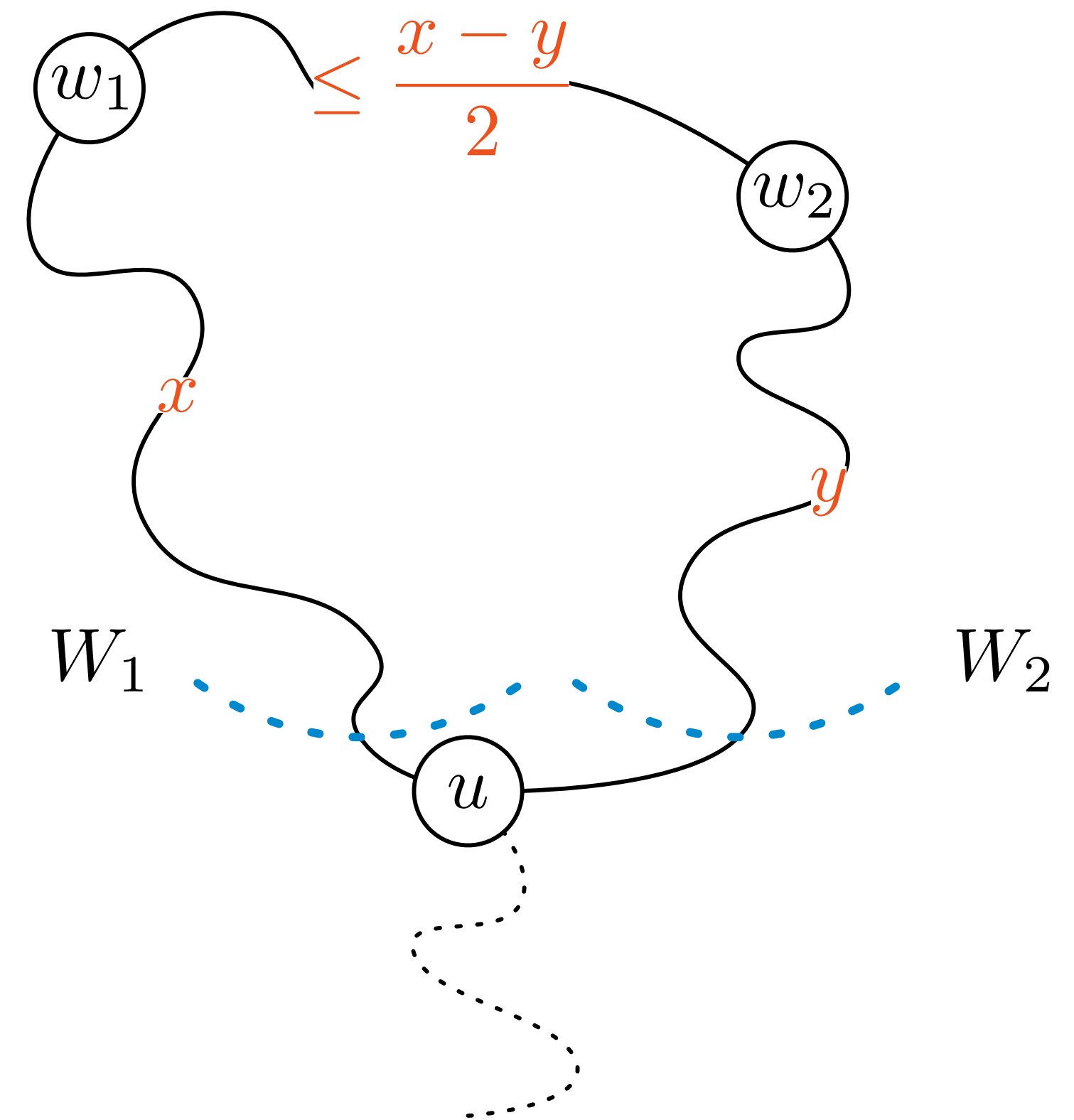
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- $d(w_1, w_2) \leq (x - y) / 2$

- $x \leq y + (x - y) / 2$

- **$x \leq y$ contradiction!**



Algorithm APSP: runtime

- Leader + BFS: $O(\text{diam}(G))$ rounds
- $|E|$ in a BFS tree: $n - 1$
- Token traverses **2** times each edge of the BFS tree
- Total number of rounds:
 - $2(2(n - 1)) + O(\text{diam}(G)) \in O(n)$ rounds

Pipelining

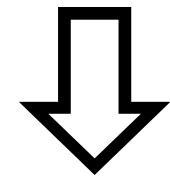
- n operations, each operation takes time t
- **Parallel**: t rounds, bad congestion
- **Sequential**: nt rounds, no congestion
- **Pipelining**: $n + t$ rounds, no congestion

Lower bound for APSP

APSP requires $\Omega(n / \log n)$ rounds

Lower bound for APSP

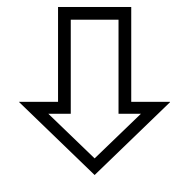
Lower bound of $\Omega(n / \log n)$ for **computing the diameter**



Lower bound of $\Omega(n / \log n)$ for **APSP**

Lower bound for APSP

Lower bound of $\Omega(n / \log n)$ for **computing the diameter**



Lower bound of $\Omega(n / \log n)$ for **APSP**

1. Show that a **lower bound** for **computing the diameter** implies the same **lower bound** for **APSP**
2. Focus only on the **lower bound** for **computing the diameter**

From APSP to Diameter

- Given a **solution for APSP**, we can compute the **diameter** in $O(\text{diam}(G))$ rounds

From APSP to Diameter

- Given a **solution for APSP**, we can compute the **diameter** in **$O(\text{diam}(G))$** rounds
 - Each node stores the maximum length of all its shortest paths in G

From APSP to Diameter

- Given a **solution for APSP**, we can compute the **diameter** in $O(\text{diam}(G))$ rounds
 - Each node stores the maximum length of all its shortest paths in G
 - Construct a BFS tree in $O(\text{diam}(G))$

From APSP to Diameter

- Given a **solution for APSP**, we can compute the **diameter** in $O(\text{diam}(G))$ rounds
 - Each node stores the maximum length of all its shortest paths in G
 - Construct a BFS tree in $O(\text{diam}(G))$
 - Leaves send their maximum distance to parent

From APSP to Diameter

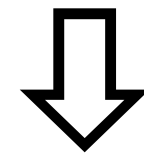
- Given a **solution for APSP**, we can compute the **diameter** in $O(\text{diam}(G))$ rounds
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 - Non-leaves compute the maximum distance among their own and the ones of its children, send to parent

From APSP to Diameter

- Given a **solution for APSP**, we can compute the **diameter** in $O(\text{diam}(G))$ rounds
 - Each node stores the maximum length of all its shortest paths in G
 - Construct a BFS tree in $O(\text{diam}(G))$
 - Leaves send their maximum distance to parent
 - Non-leaves compute the maximum distance among their own and the ones of its children, send to parent
 - Broadcast the value of the diameter

Diameter lower bound \implies APSP lower bound

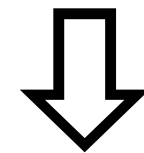
- Compute diameter in: $T(\text{APSP}) + O(\text{diam}(G))$ rounds
- If computing the **diameter requires $\Omega(n / \log n)$** rounds




- **APSP must require $\Omega(n / \log n)$** rounds in all graphs with diameter $o(n / \log n)$
 - $T(\text{APSP}) + o(n / \log n) \in \Omega(n / \log n) \implies \mathbf{T(\text{APSP}) \in \Omega(n / \log n)}$

Diameter lower bound \implies APSP lower bound

- Compute diameter in: $T(\text{APSP}) + O(\text{diam}(G))$ rounds
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Computing the diameter

- Computing the diameter requires $\Omega(n/\log n)$ [Frischknecht, Holzer, Wattenhofer]
- The proof uses known results from 2-party **communication complexity**
 - Studies the **minimum amount of communication** (number of bits) needed in order to compute functions whose **arguments** are **distributed** among several parties
 - **Set disjointness** between **2 communication parties**

Set disjointness

Set disjointness



Set disjointness



Set disjointness



A



Set disjointness



A



B

Set disjointness

- $A, B \subseteq \{1, 2, \dots, k\}$



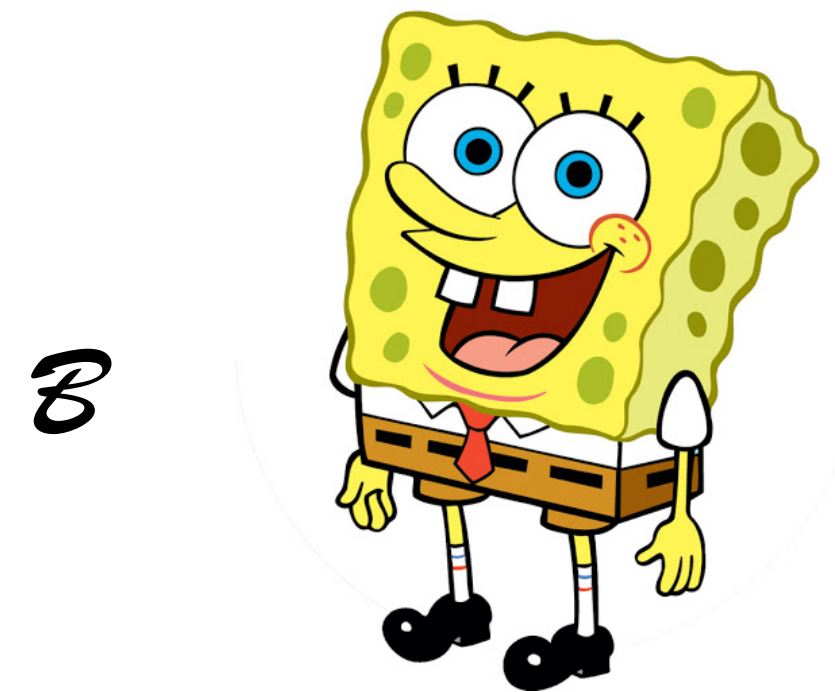
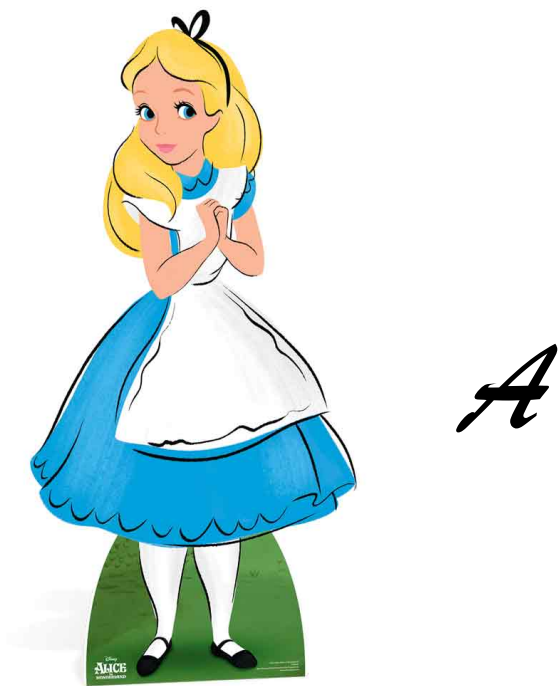
A



B

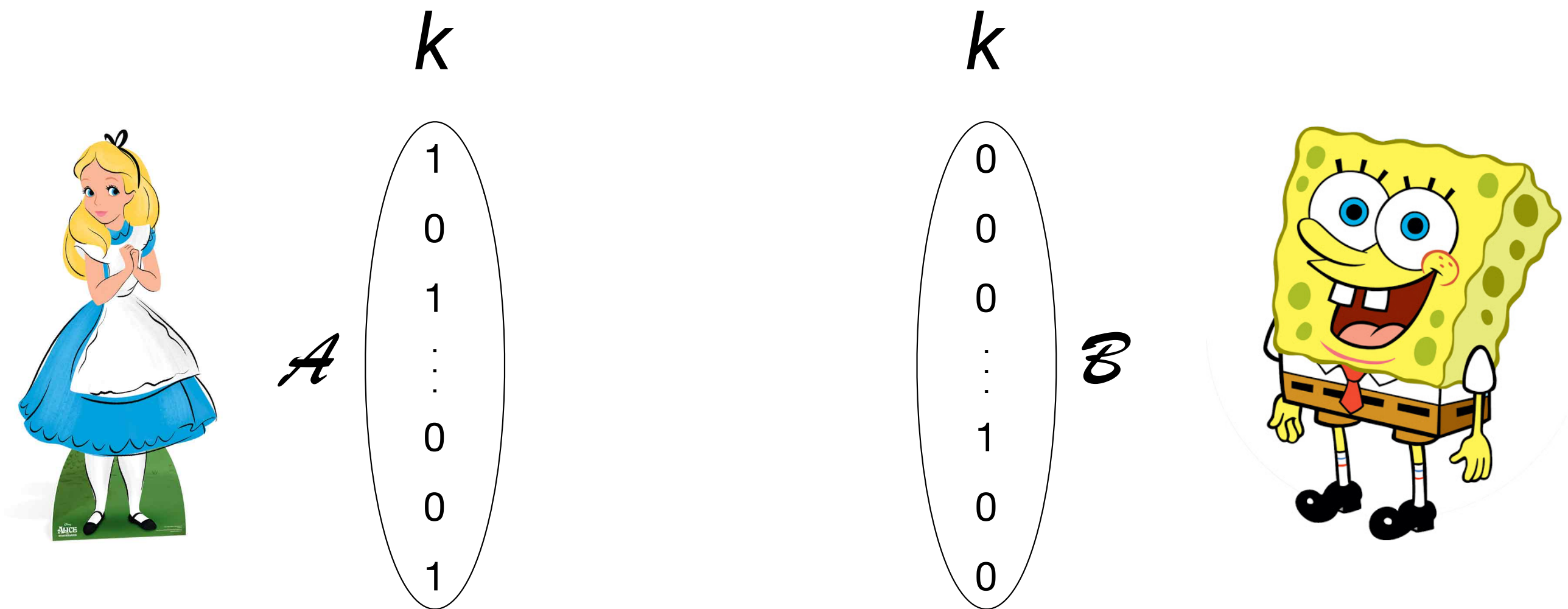
Set disjointness

- $A, B \subseteq \{1, 2, \dots, k\}$
- Output: 1 if $A \cap B = \emptyset$; 0 otherwise



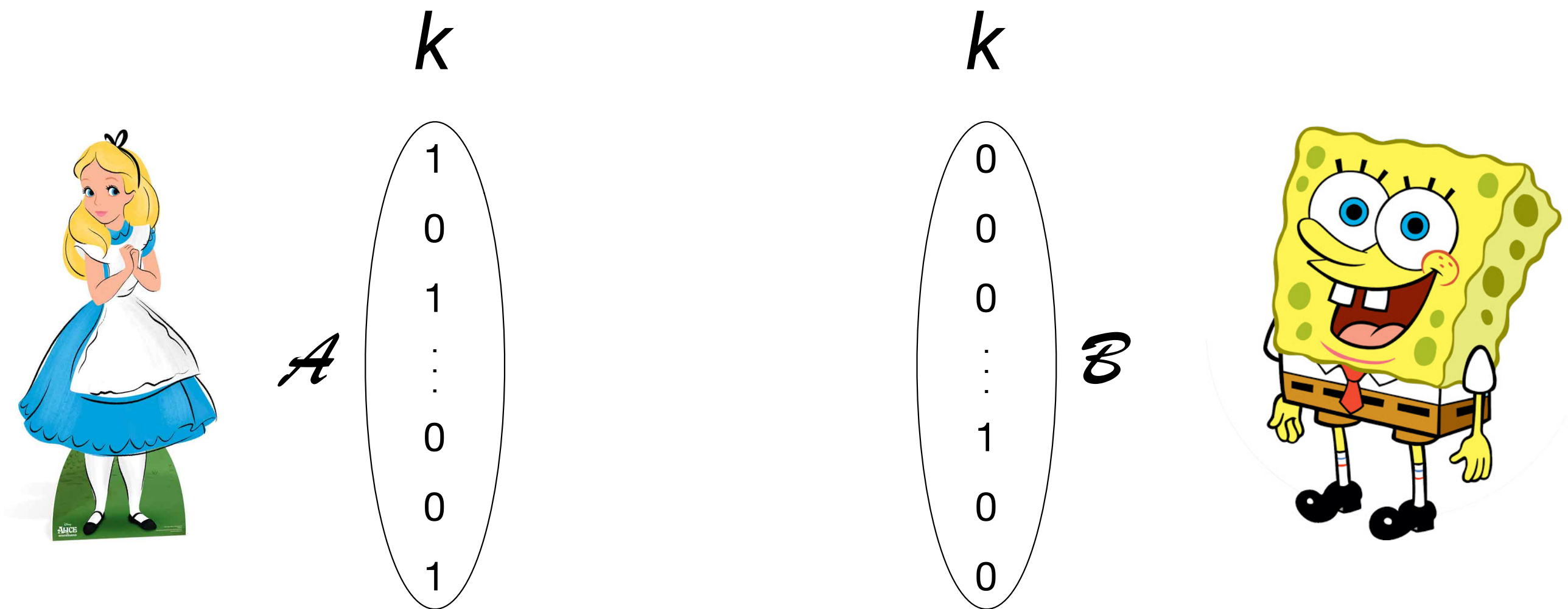
Set disjointness

- $A, B \subseteq \{1, 2, \dots, k\}$
- Output: 1 if $A \cap B = \emptyset$; 0 otherwise
- String of k bits: 1 in position i if the i -th element is present, 0 otherwise

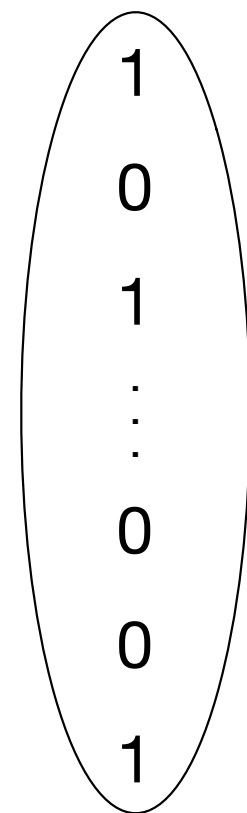


Set disjointness

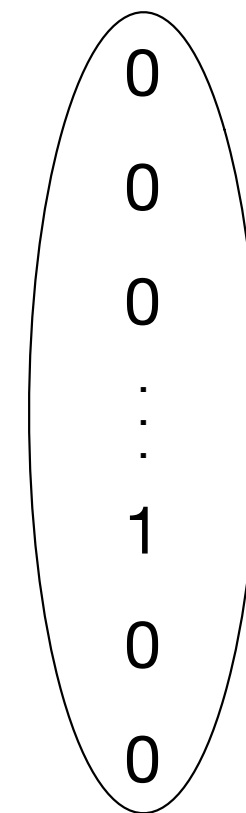
Alice and Bob need to **exchange $\Omega(k)$ bits** in order to solve set disjointness



Computing the diameter: Lower bound idea



k



k

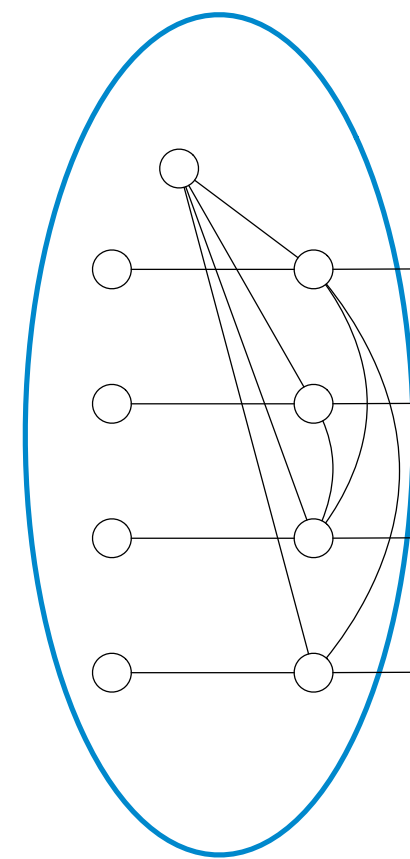


Computing the diameter: Lower bound idea

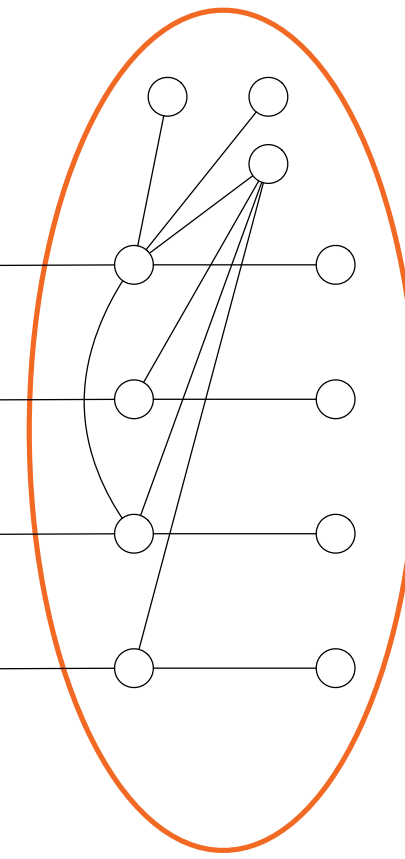
Algorithm that computes the diameter \implies Solution to the set disjointness problem



k



$\Theta(\sqrt{k})$



$\Theta(\sqrt{k})$



k

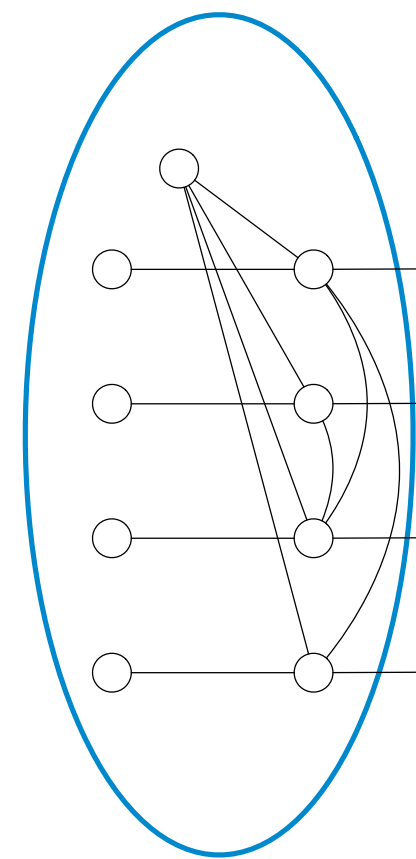
Computing the diameter: Lower bound idea

Diameter = 4 \Rightarrow the sets are **disjoint**

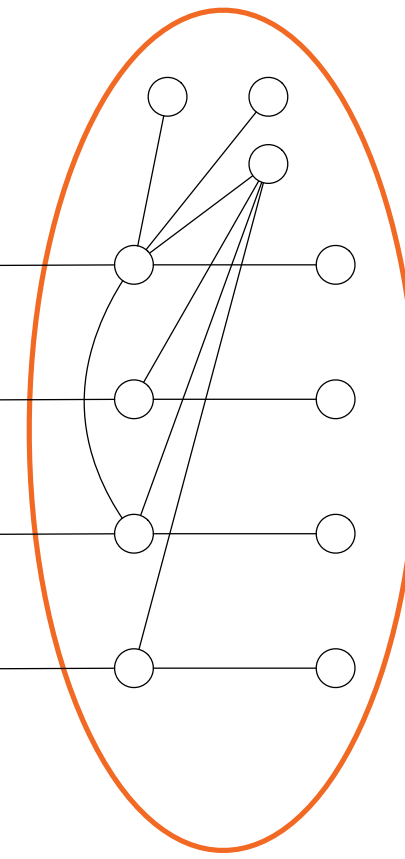
Diameter \geq 5 \Rightarrow the sets are **not** disjoint



k



$\Theta(\sqrt{k})$



$\Theta(\sqrt{k})$



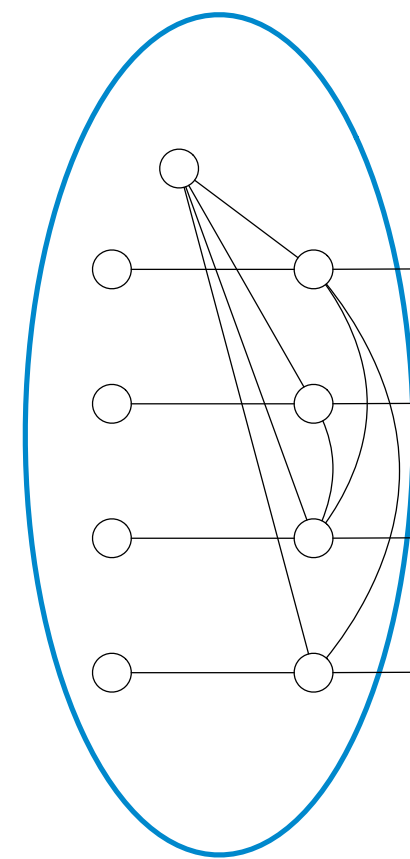
k

Computing the diameter: Lower bound idea

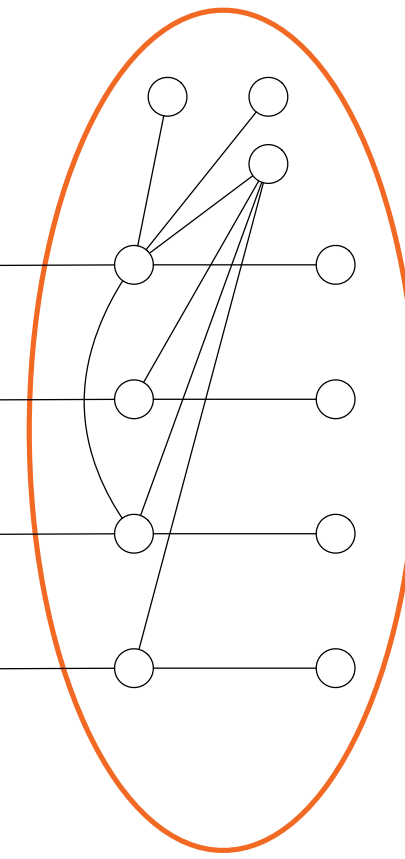
Diameter in $o(n / \log n)$ rounds \Rightarrow Diameter exchanging $o(k)$ bits \Rightarrow
Set disjointness exchanging $o(k)$ bits \Rightarrow Contradict the lower bound



k



$\Theta(\sqrt{k})$



$\Theta(\sqrt{k})$



k

Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

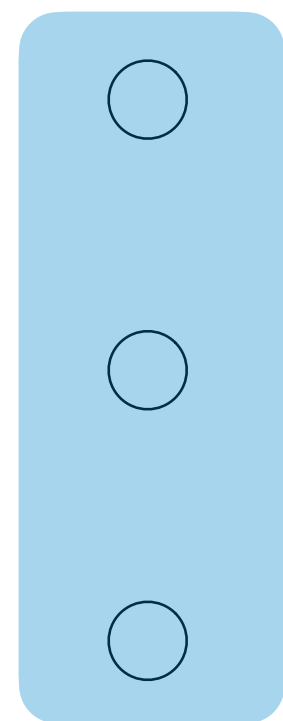
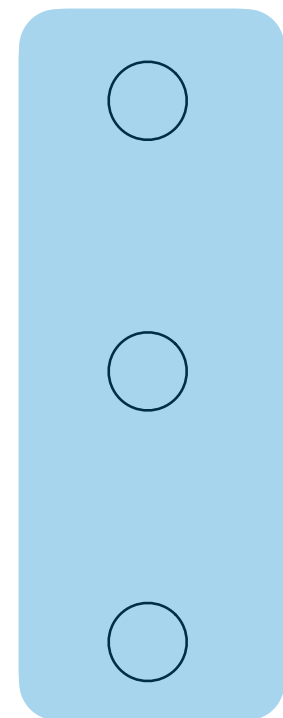


Lower bound for computing the diameter

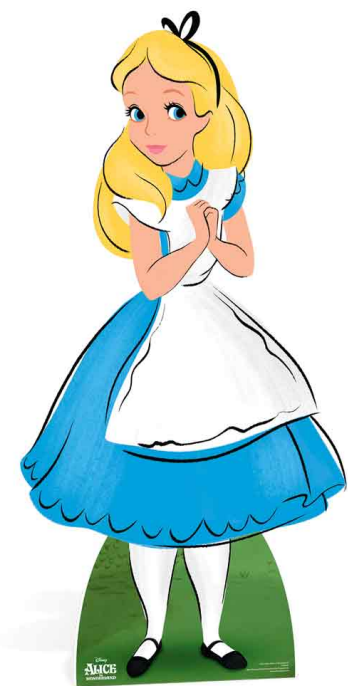


1
0
0
0
1
1
0
1
0

$k = 9$



Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

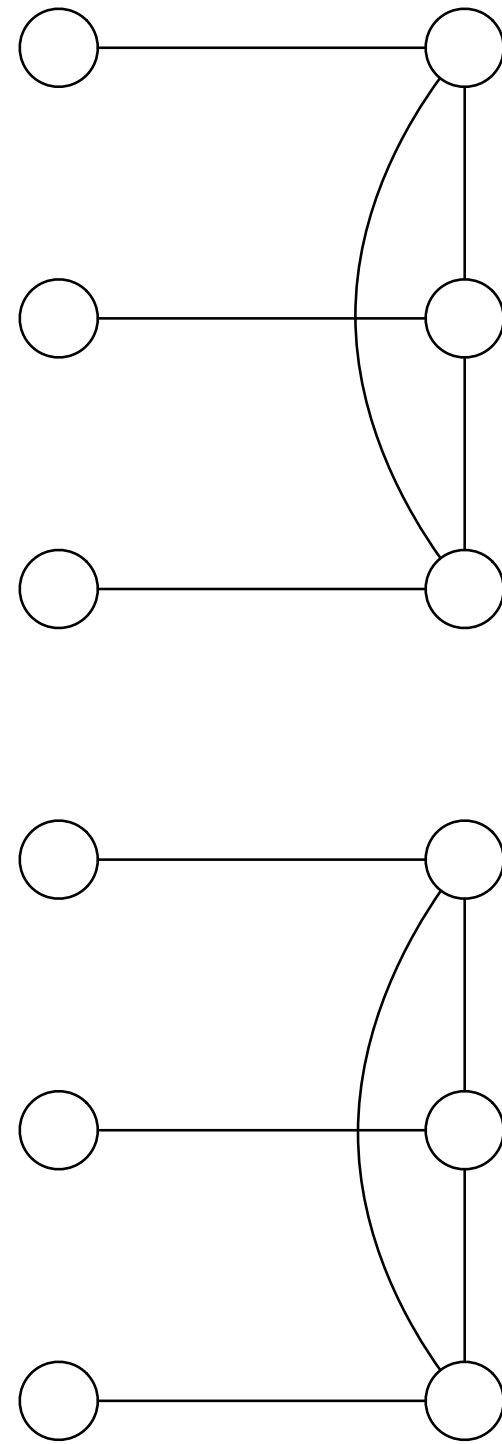


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

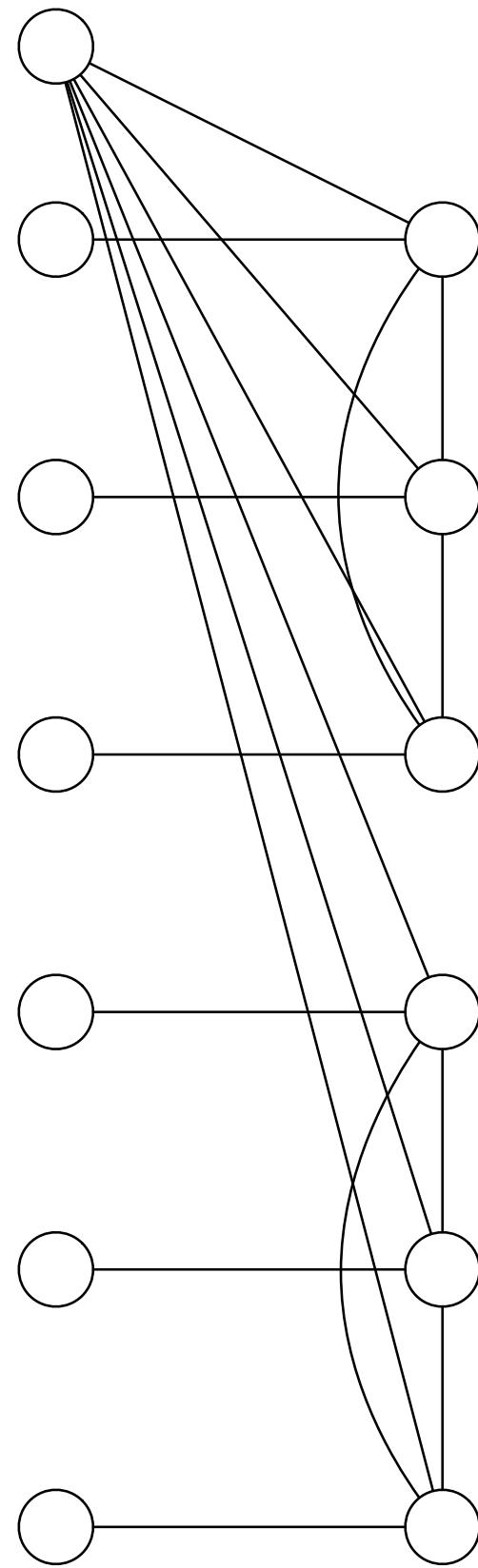


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

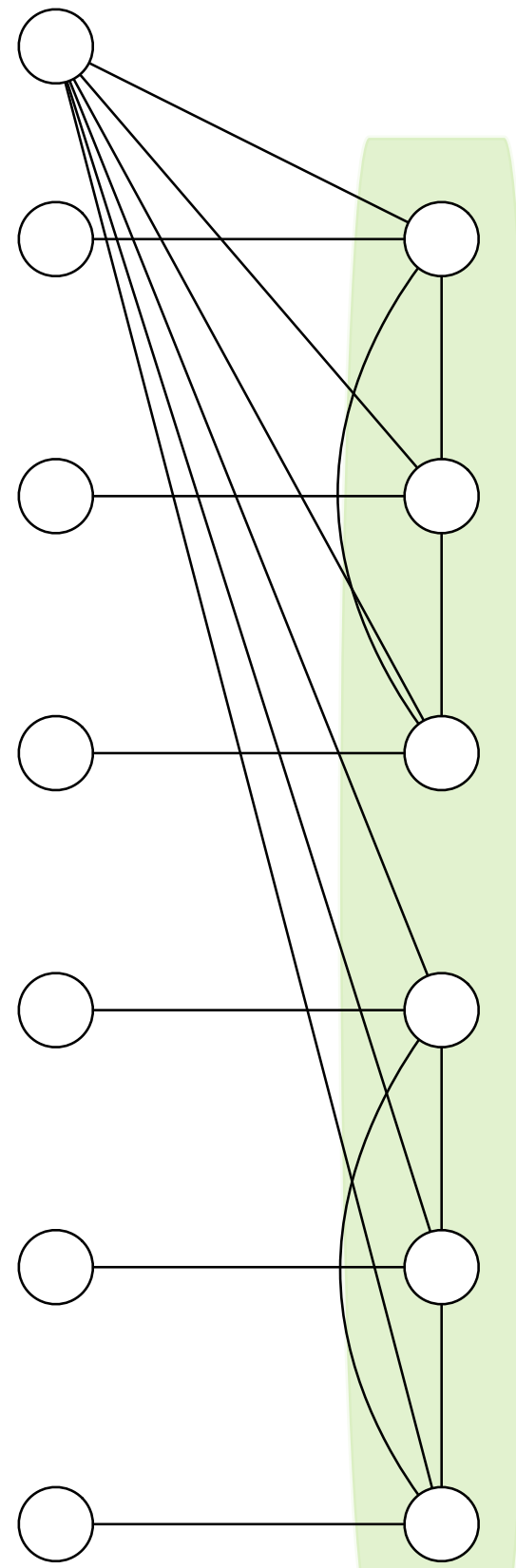


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

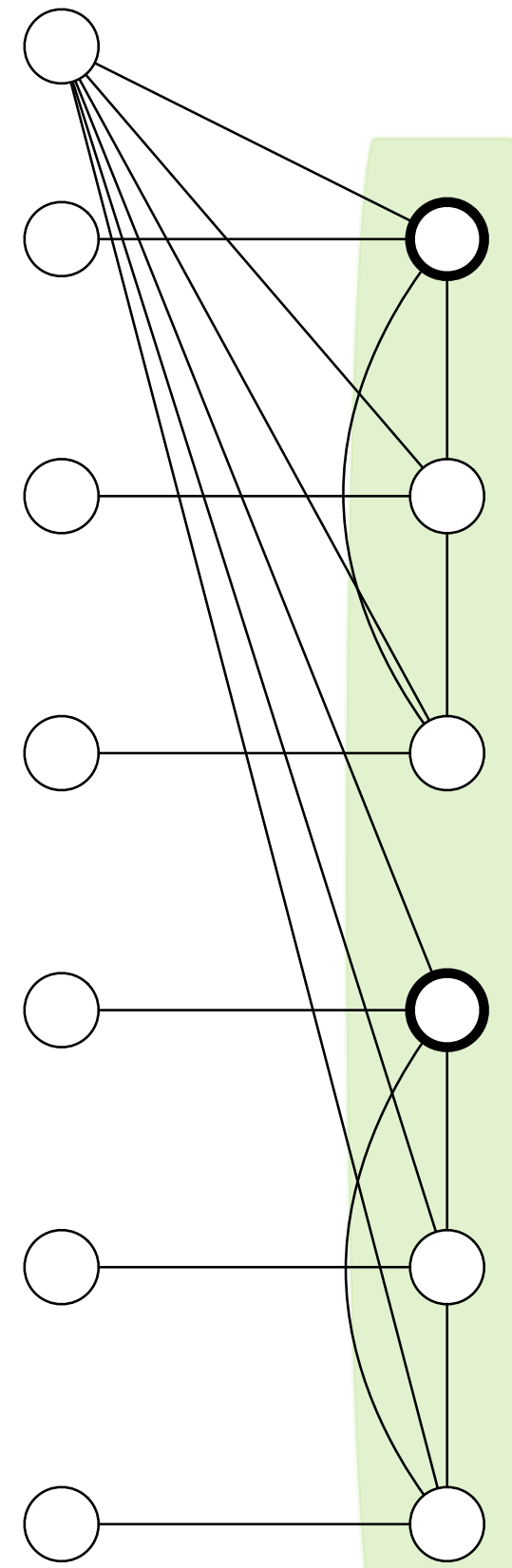


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

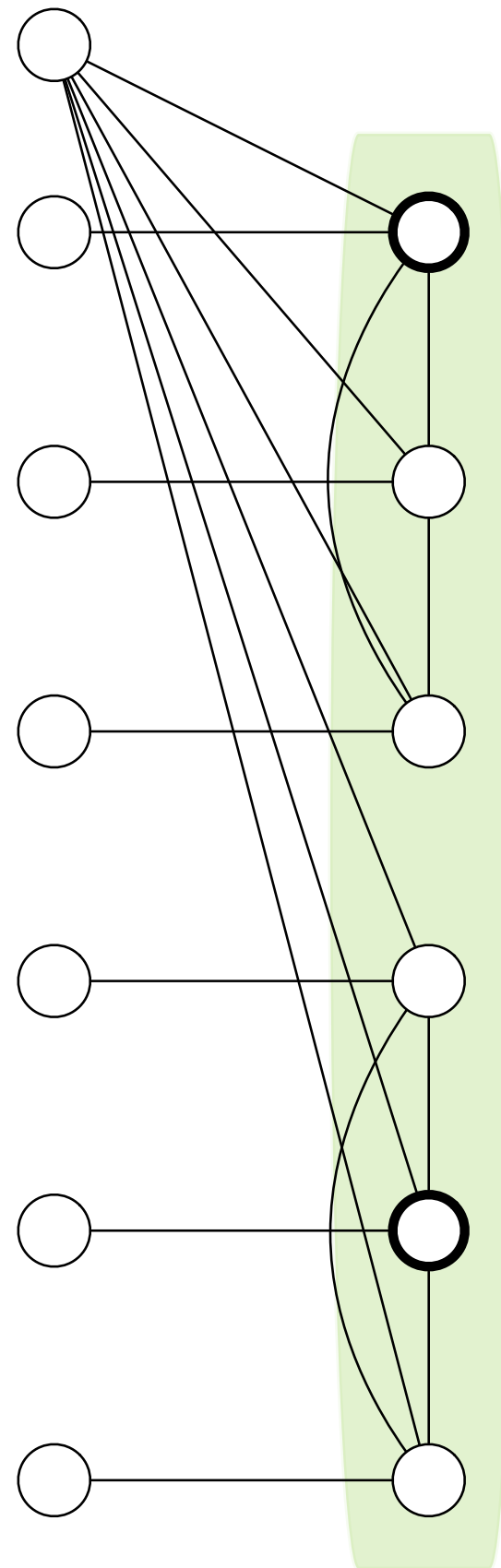


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

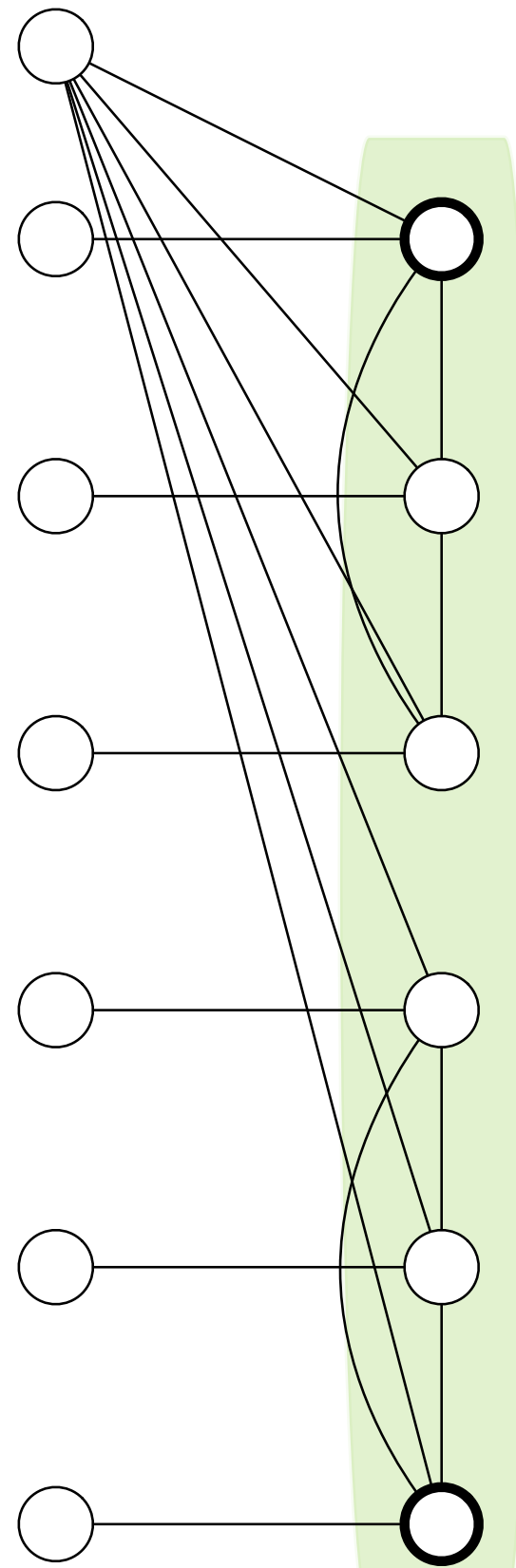


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

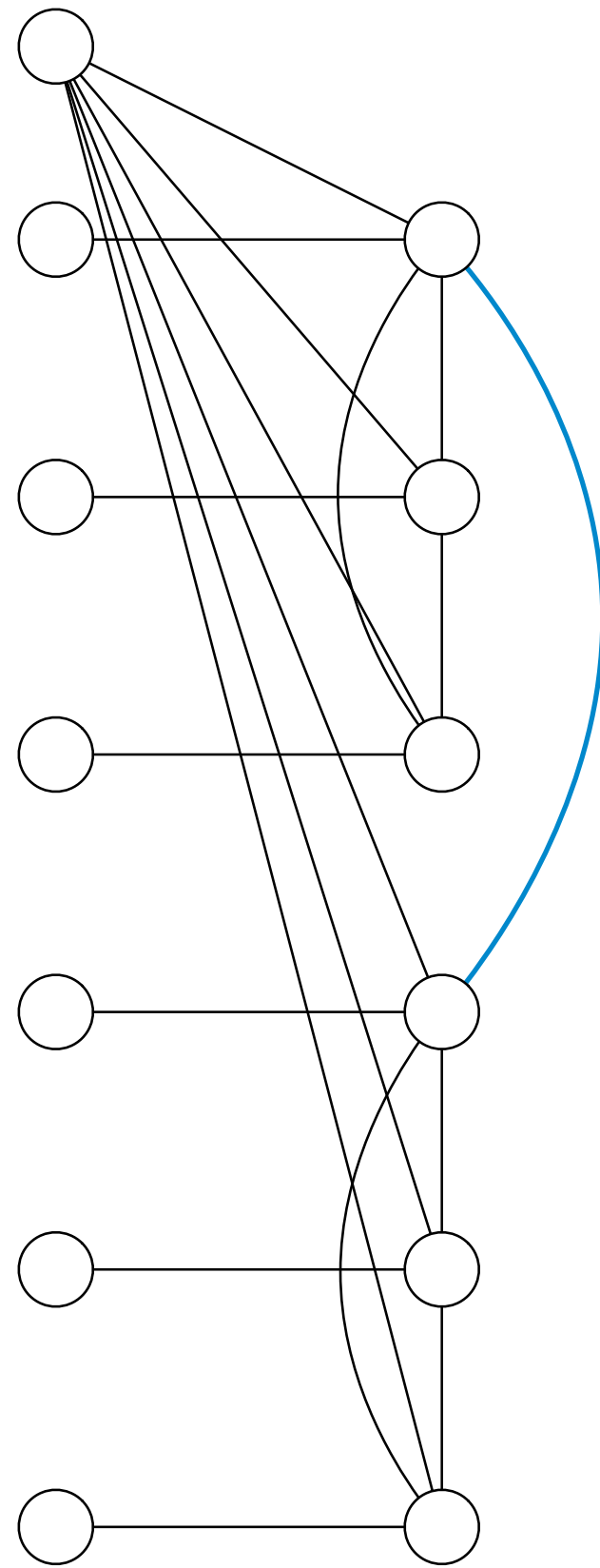


Lower bound for computing the diameter

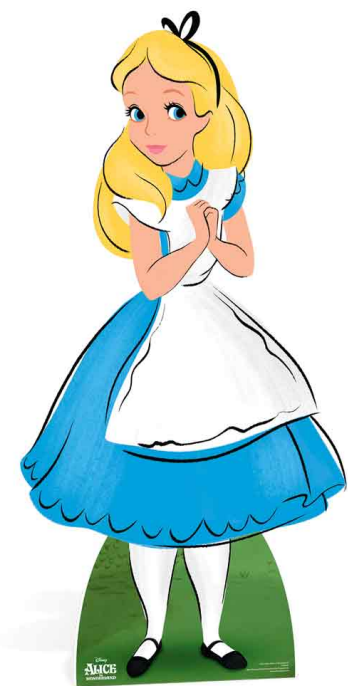


1
0
0
0
1
1
0
1
0

$k = 9$

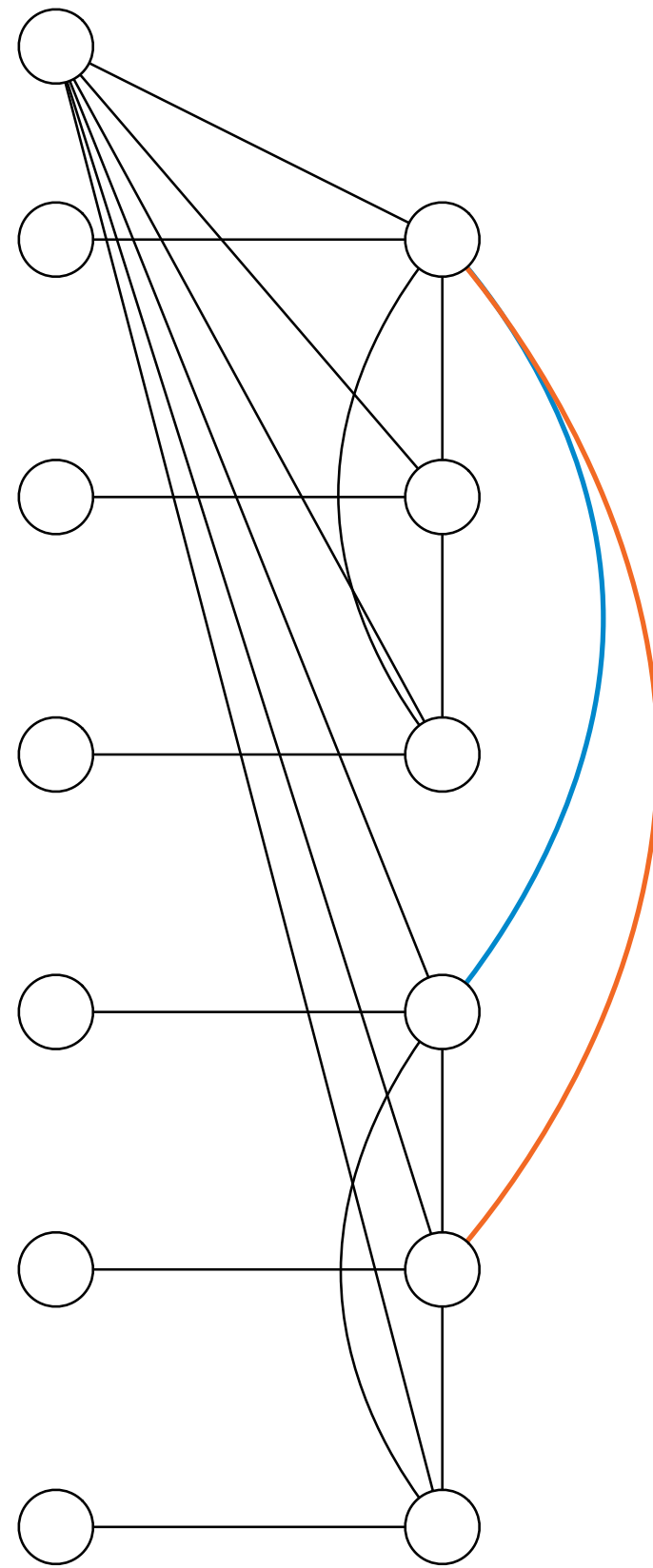


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

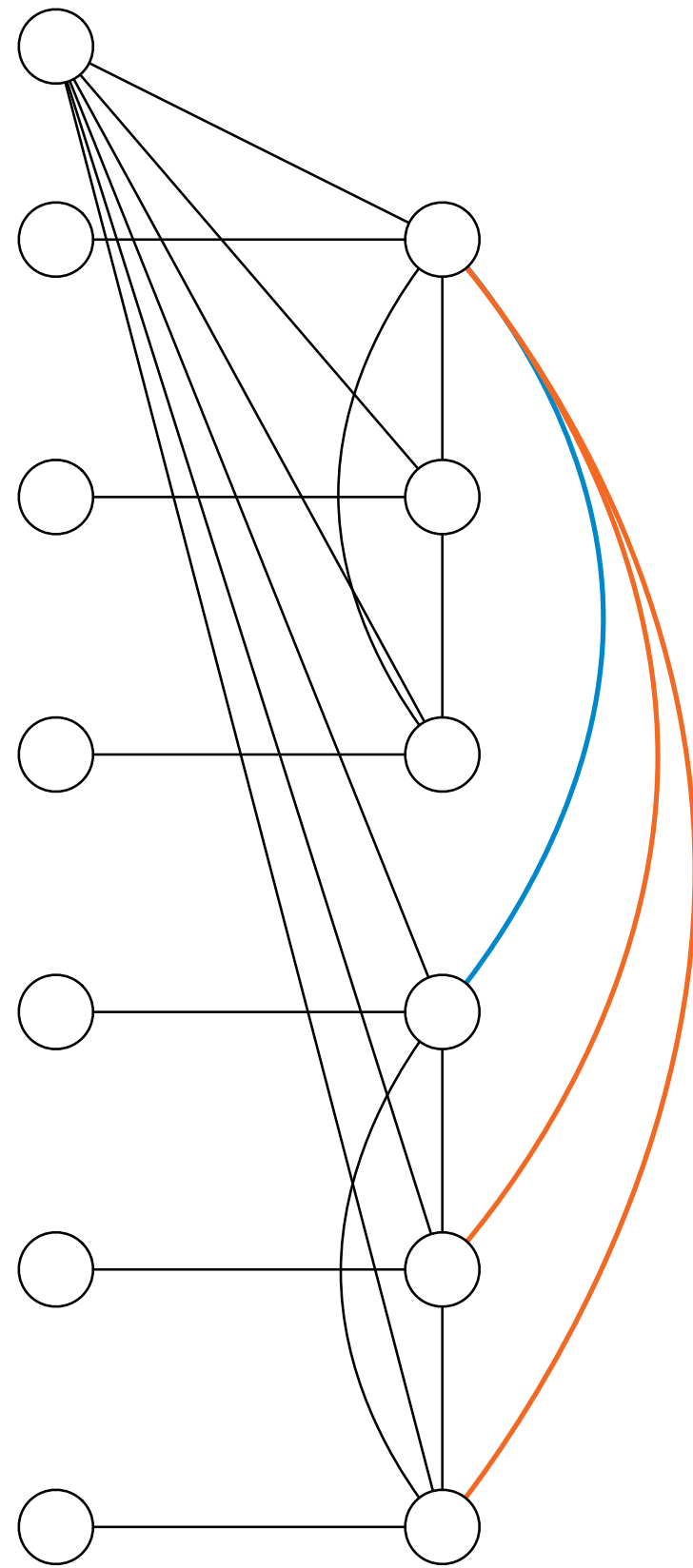


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

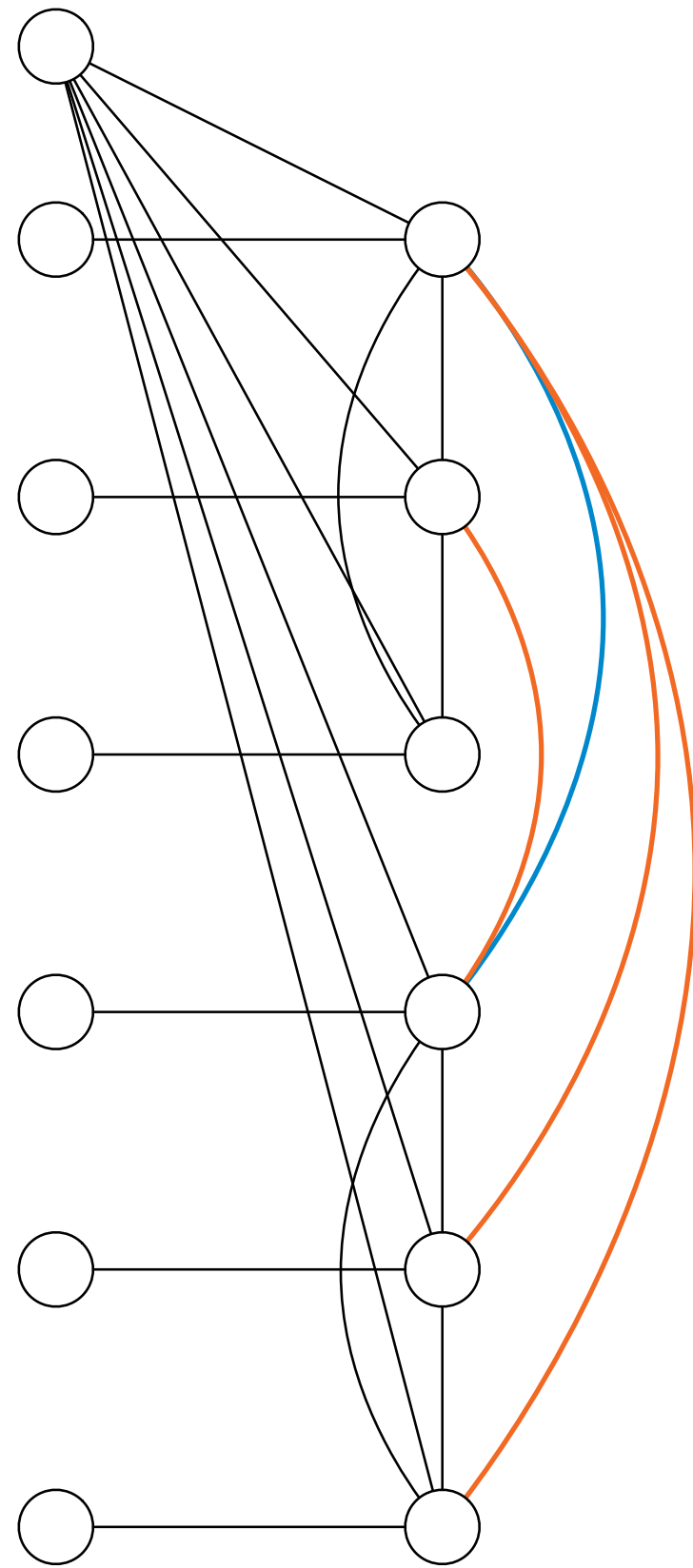


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

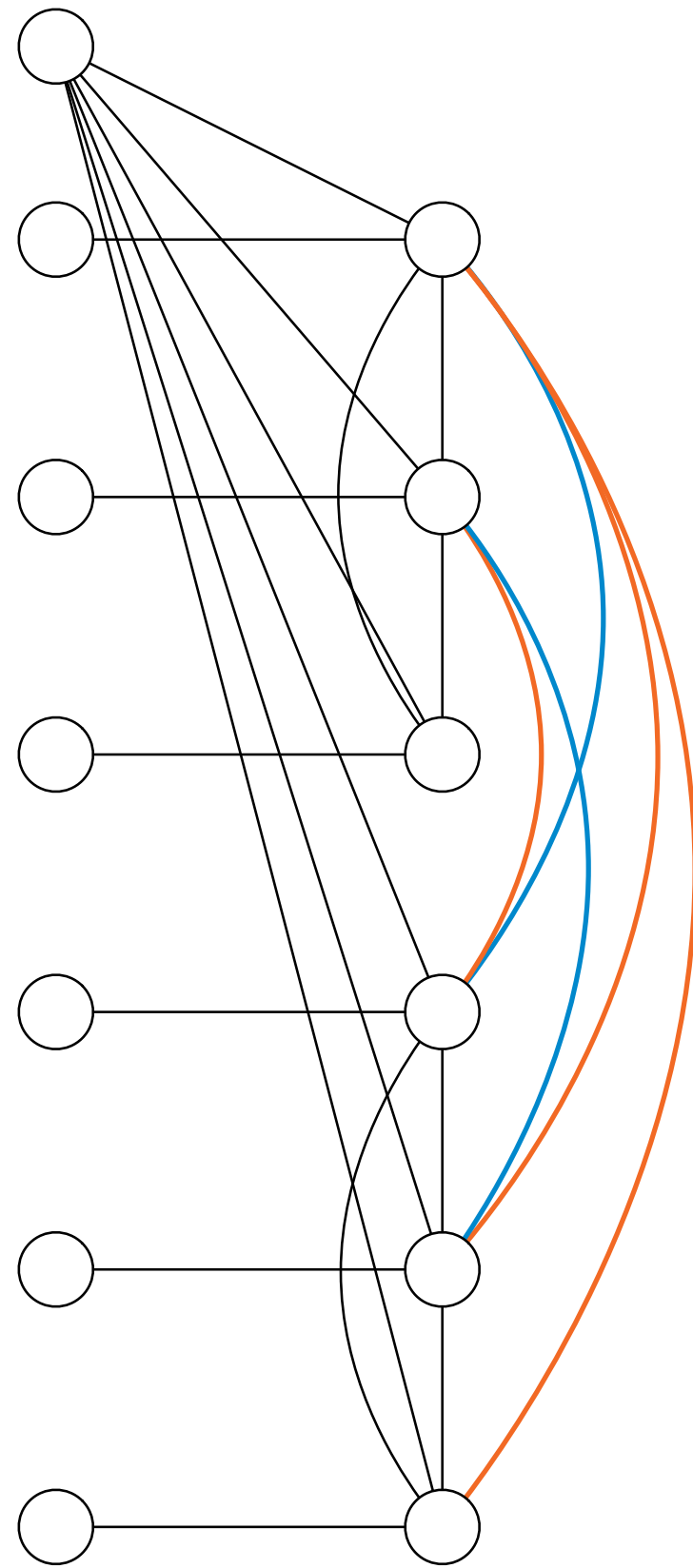


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

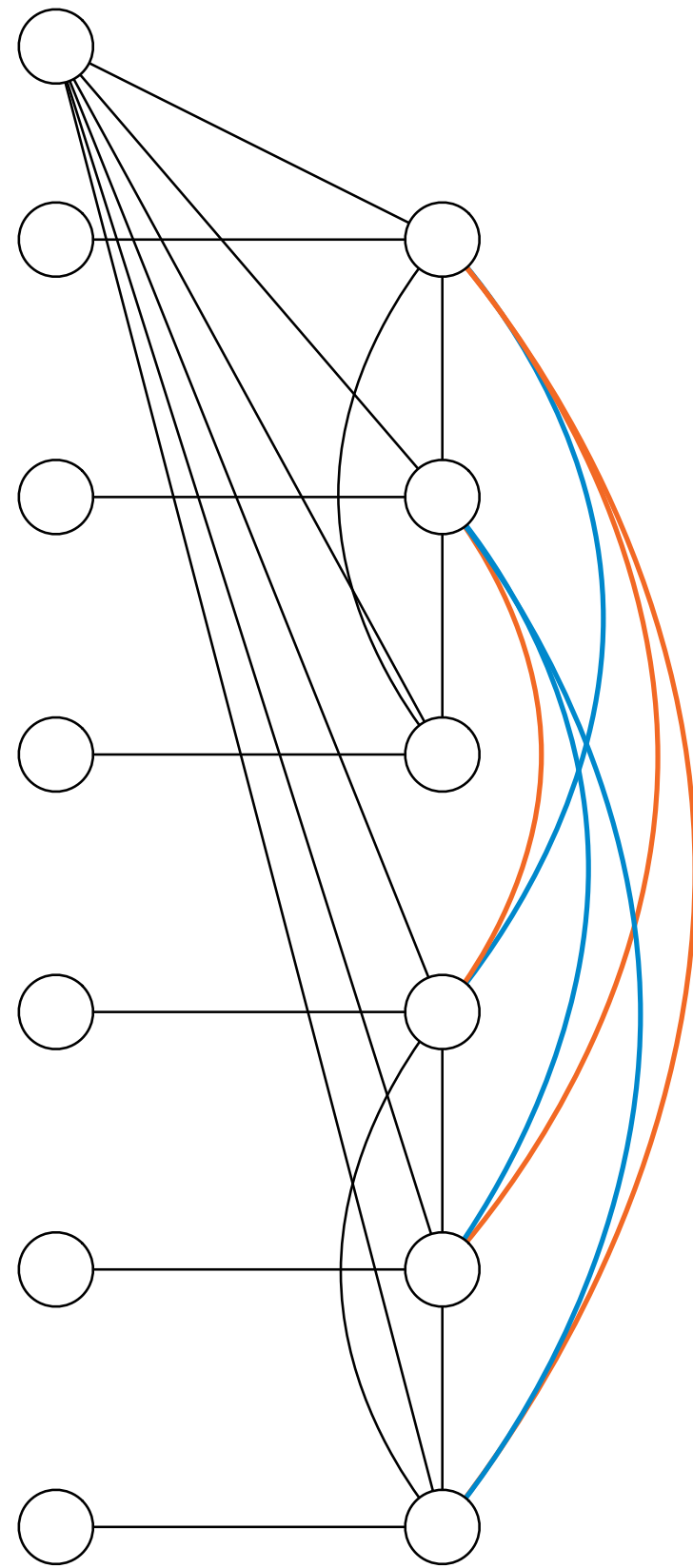


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$

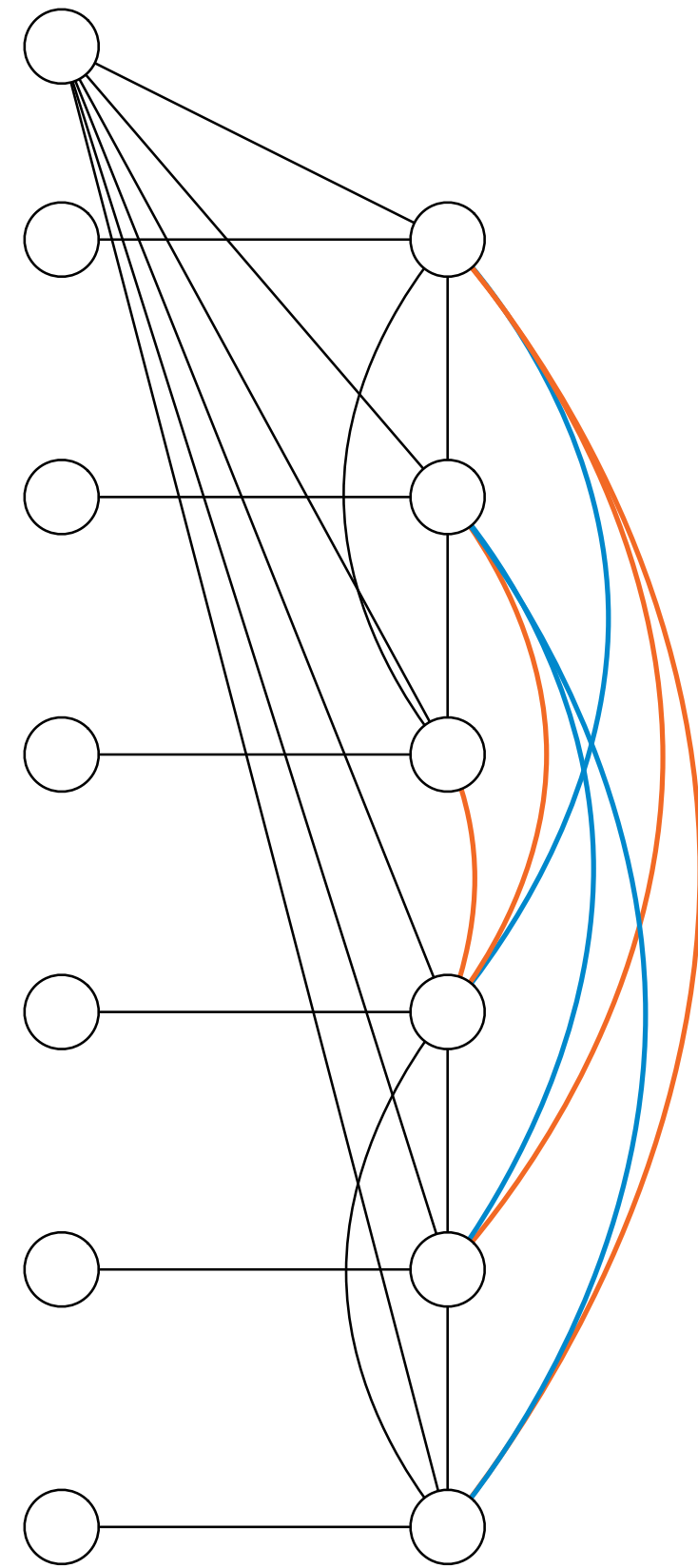


Lower bound for computing the diameter



1
0
0
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0
1
1
0
1
0

$k = 9$

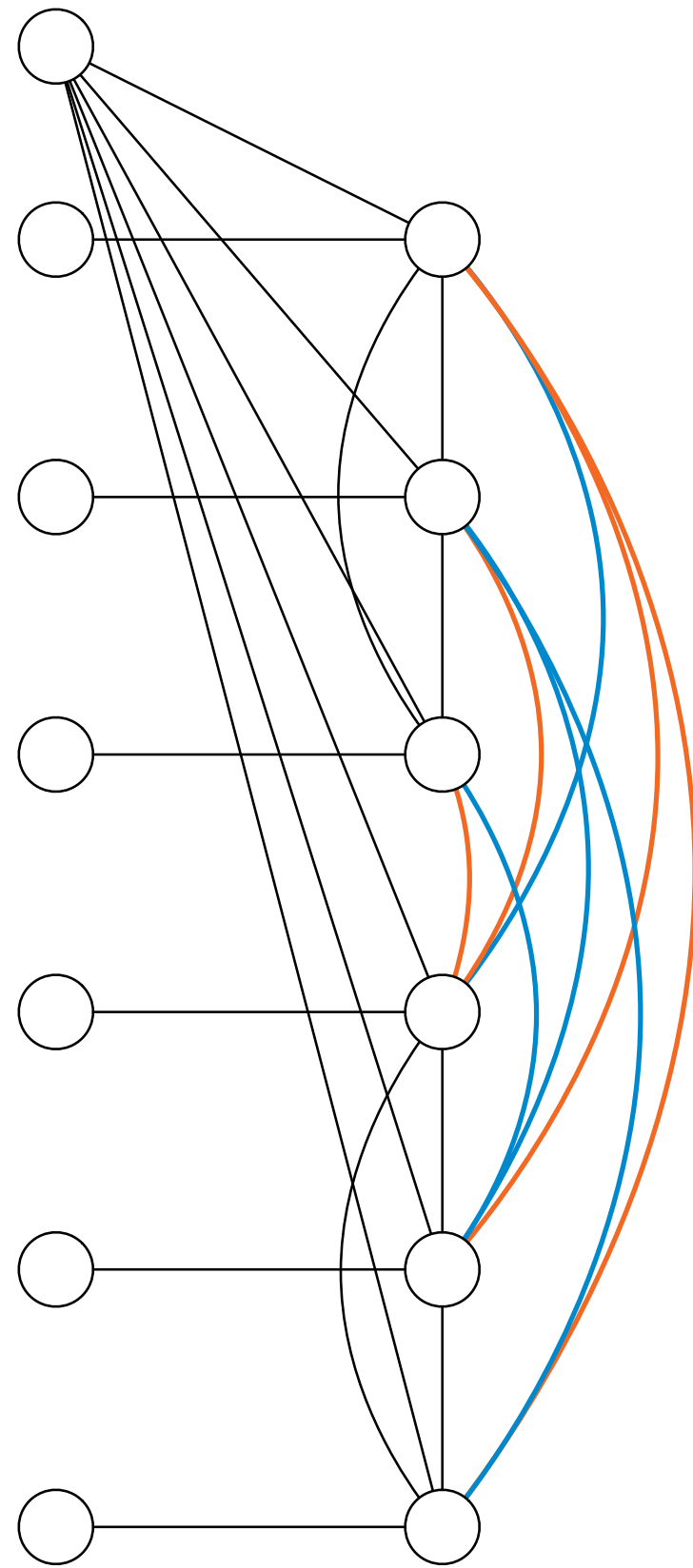


Lower bound for computing the diameter



1
0
0
0
0
1
1
0
1
0

$k = 9$

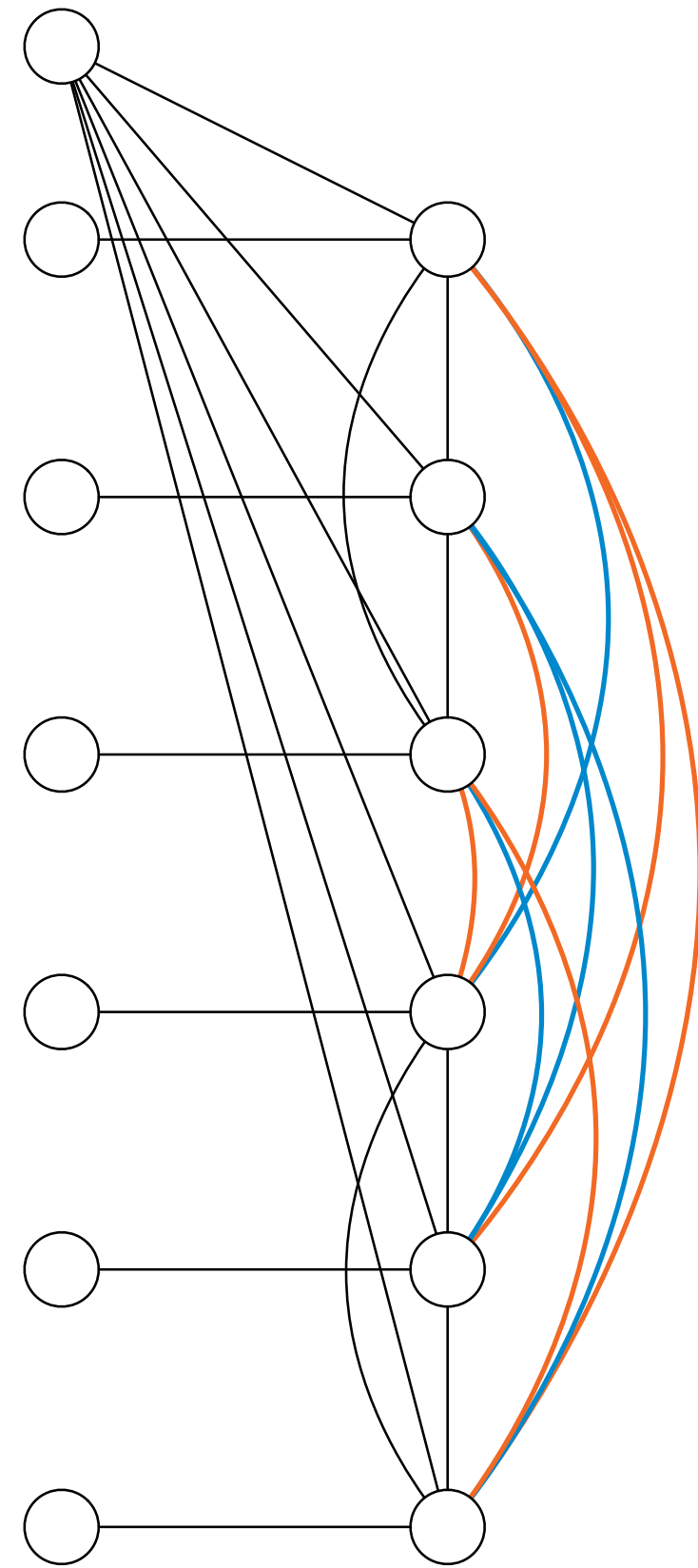


Lower bound for computing the diameter

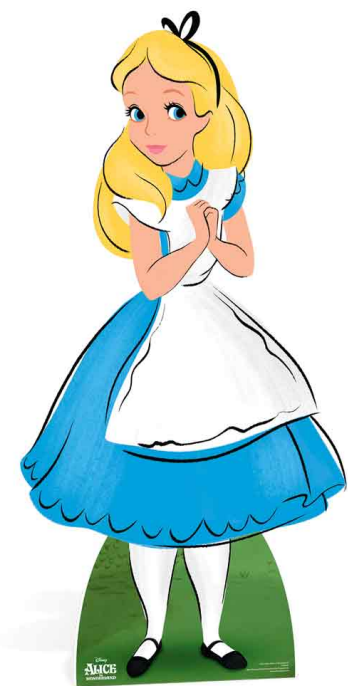


1
0
0
0
0
1
1
0
1
0

$k = 9$

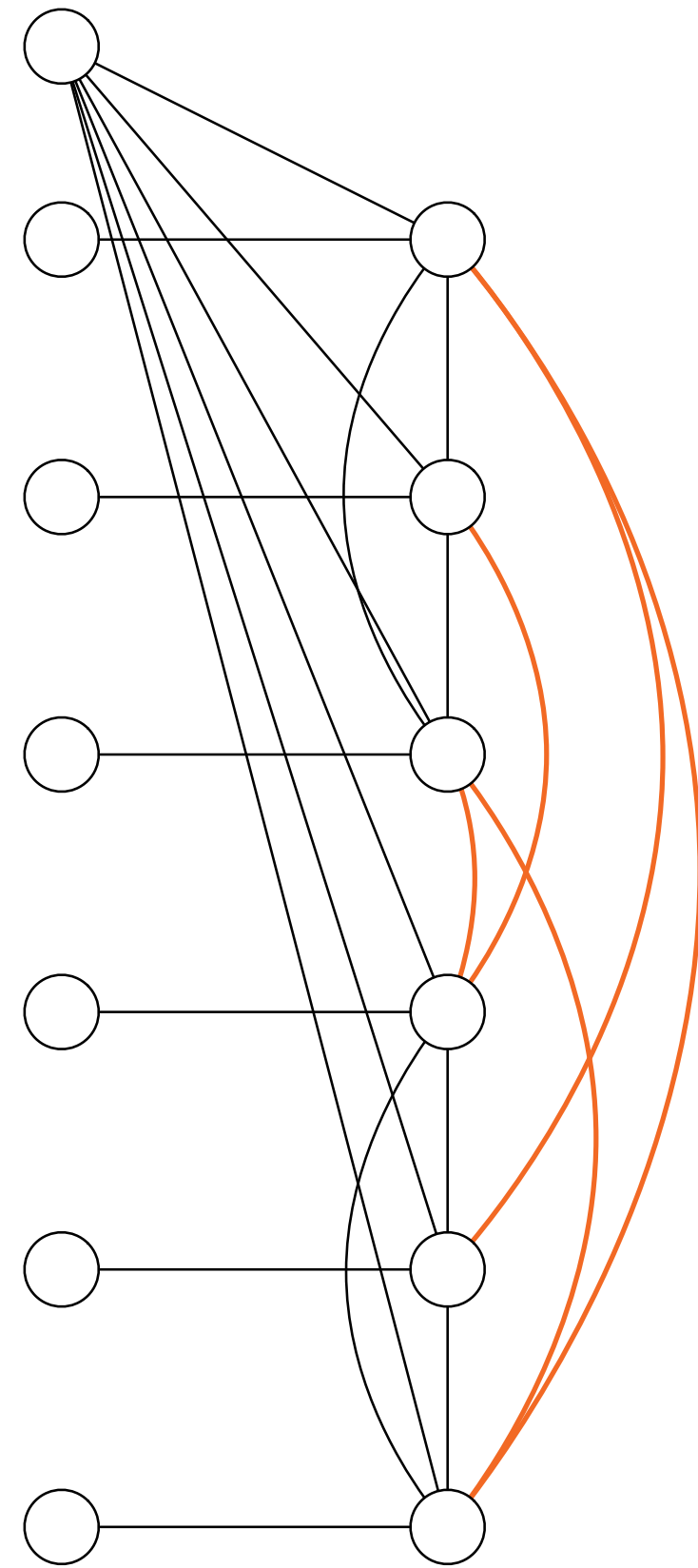


Lower bound for computing the diameter



1
0
0
0
0
1
1
0
1
0

$k = 9$

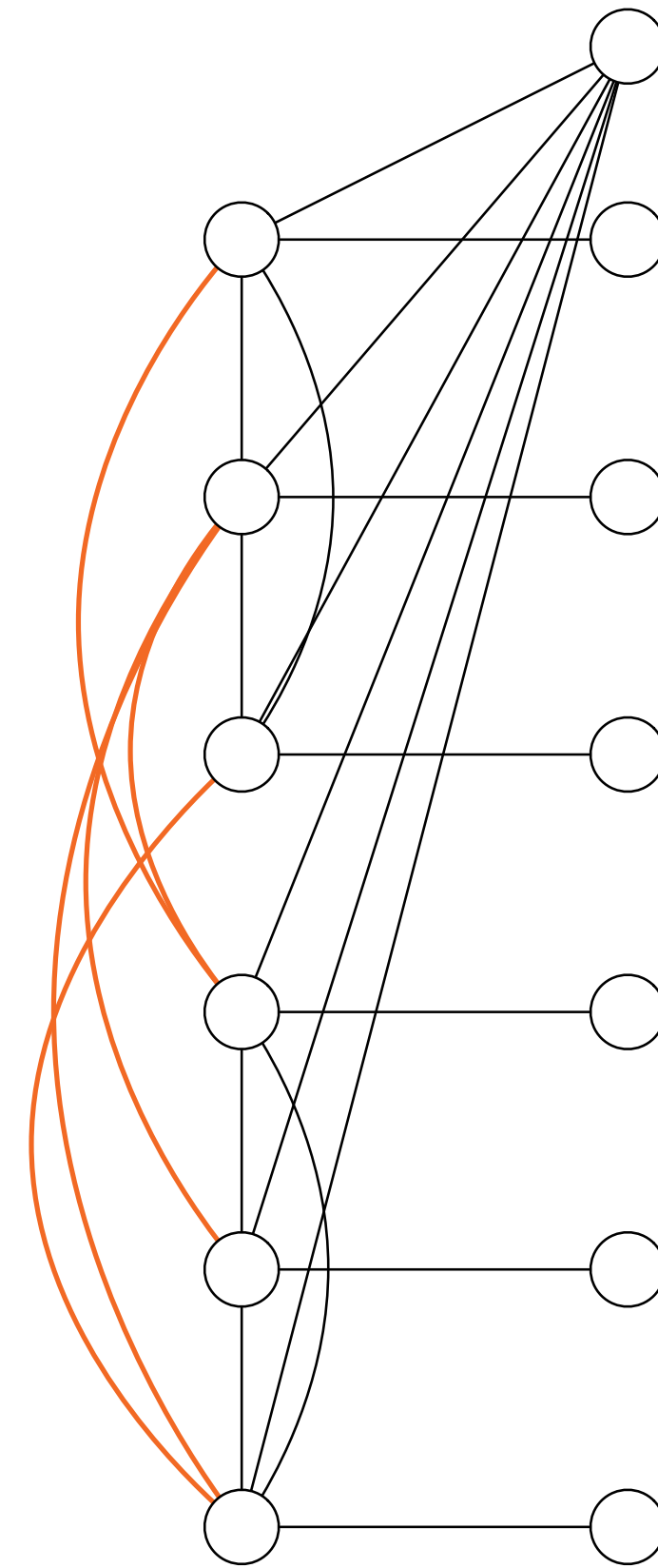


Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$



0
1
1
0
0
0
1
1
0



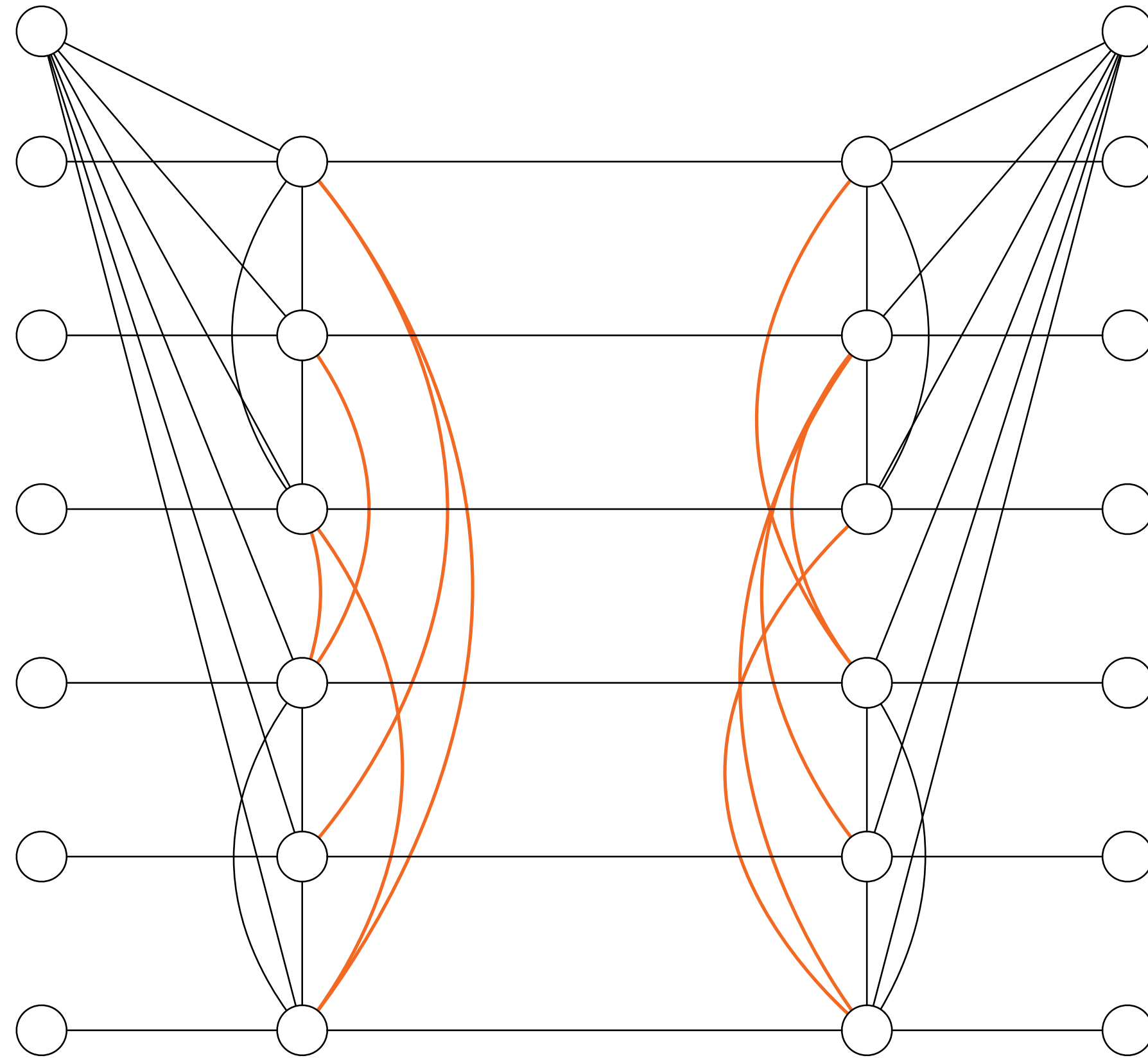
$k = 9$

Lower bound for computing the diameter



1
0
0
0
1
1
0
1
0

$k = 9$



0
1
1
0
0
0
1
1
0



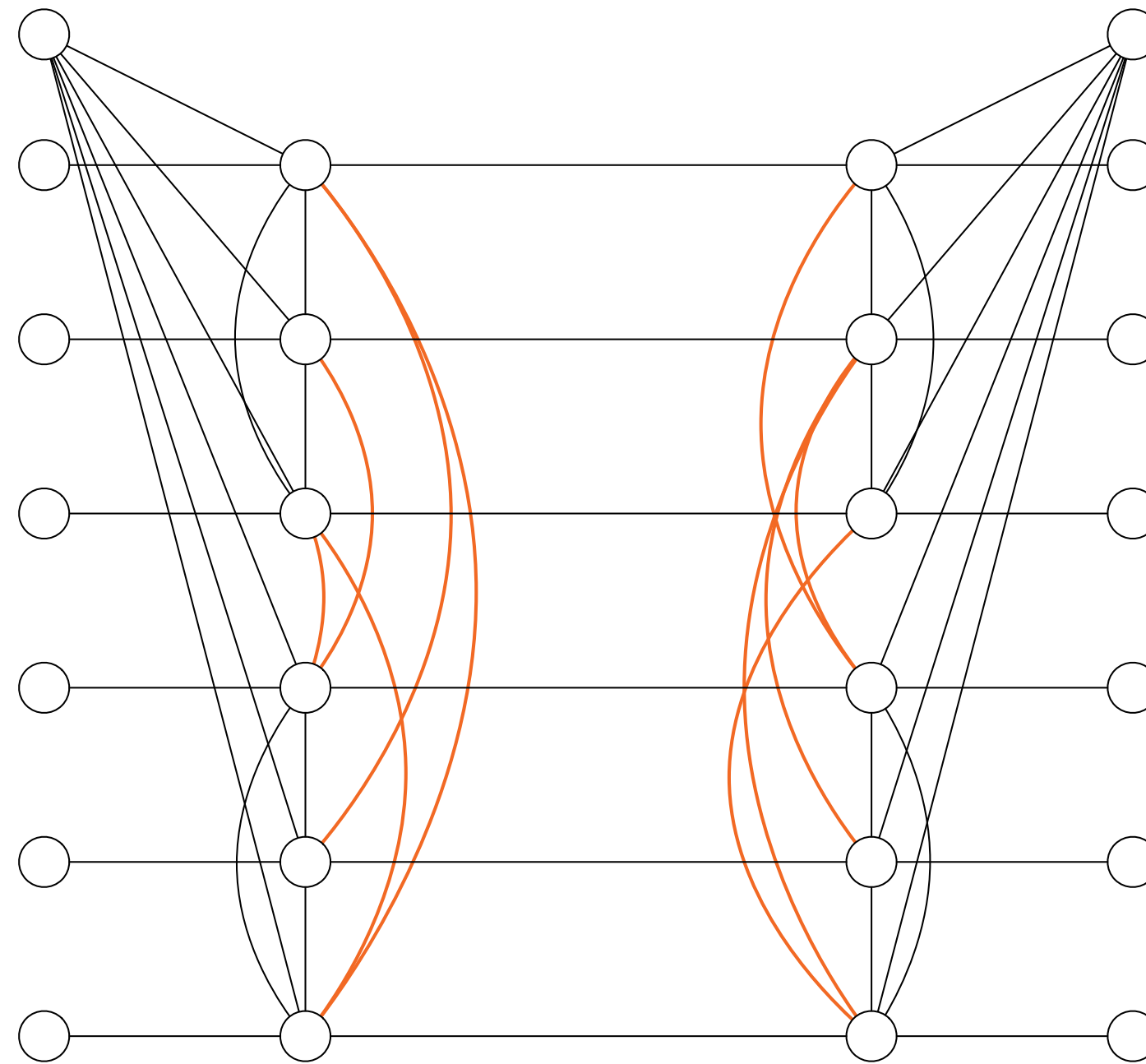
$k = 9$

Lower bound for computing the diameter

Diameter = 4 if the **sets are disjoint**, otherwise **diameter ≥ 5**



1
0
0
0
1
1
0
1
0

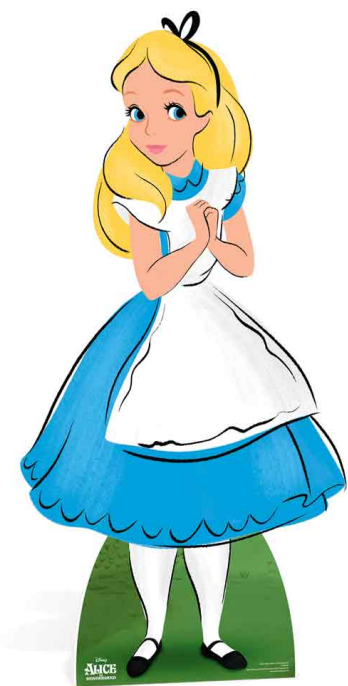


0
1
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0
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0

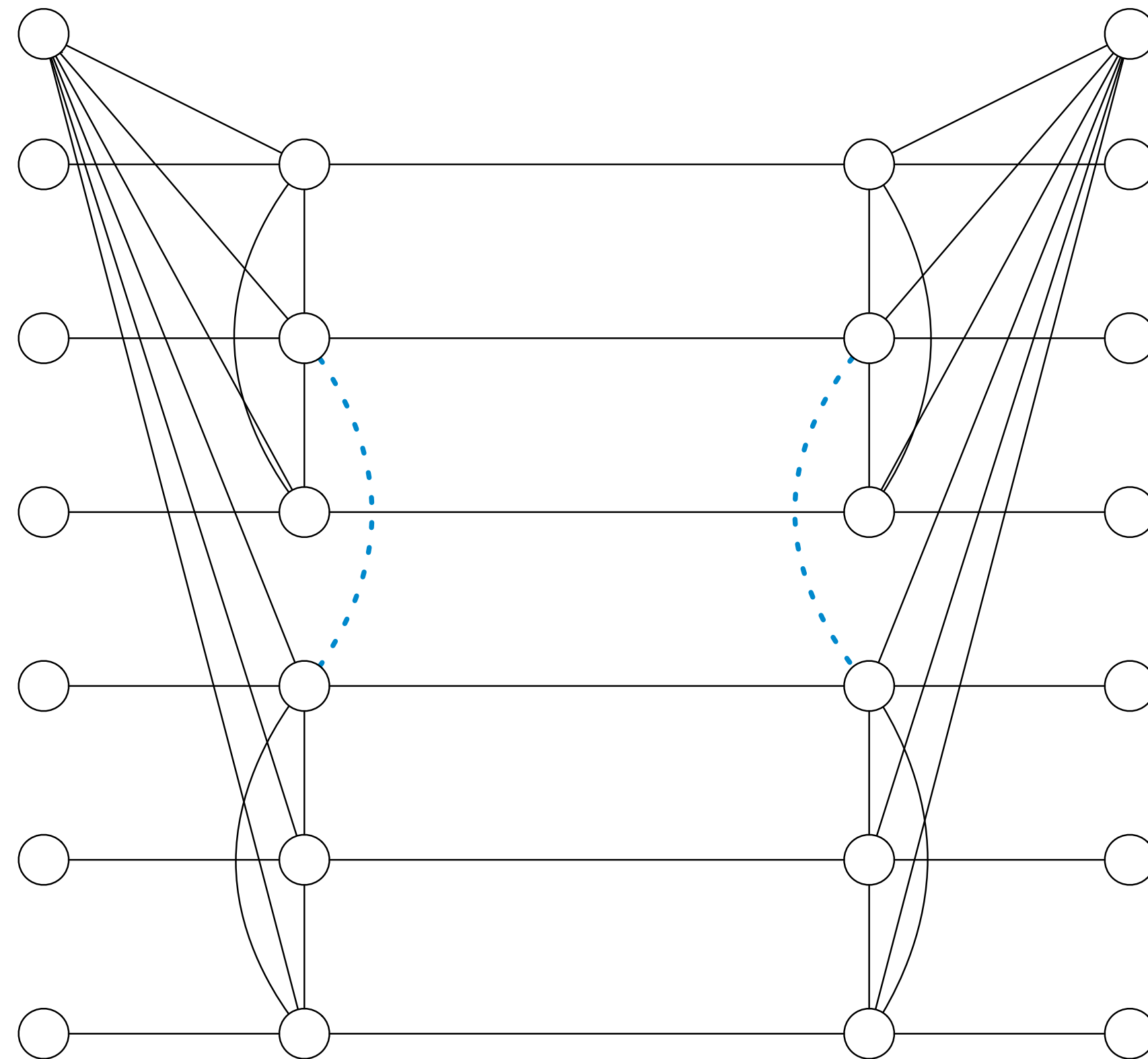


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
1
1
0
1
0



0
1
1
1
0
0
1
0
0

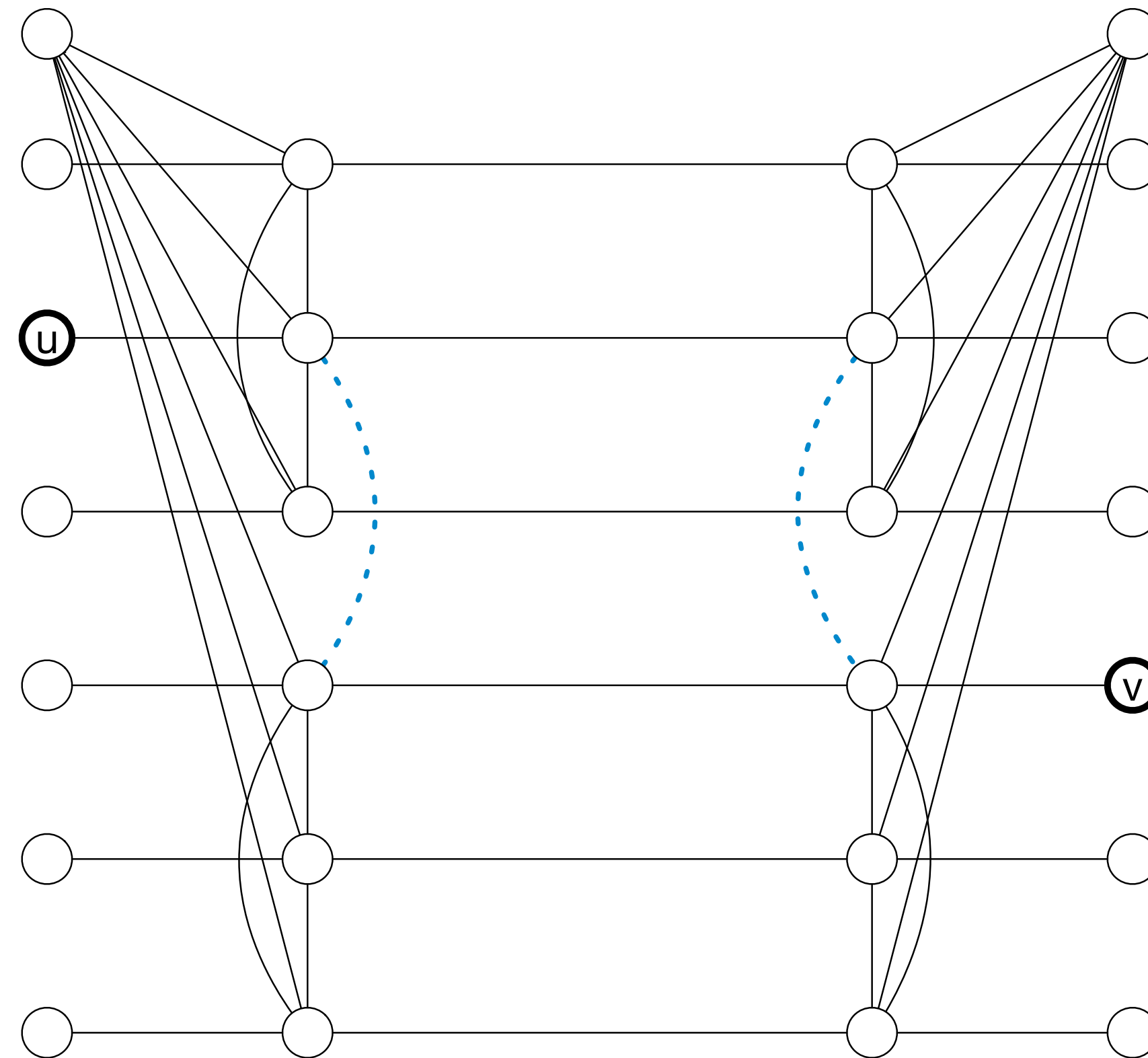


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
1
1
0
1
0



0
1
1
1
0
0
1
0
0

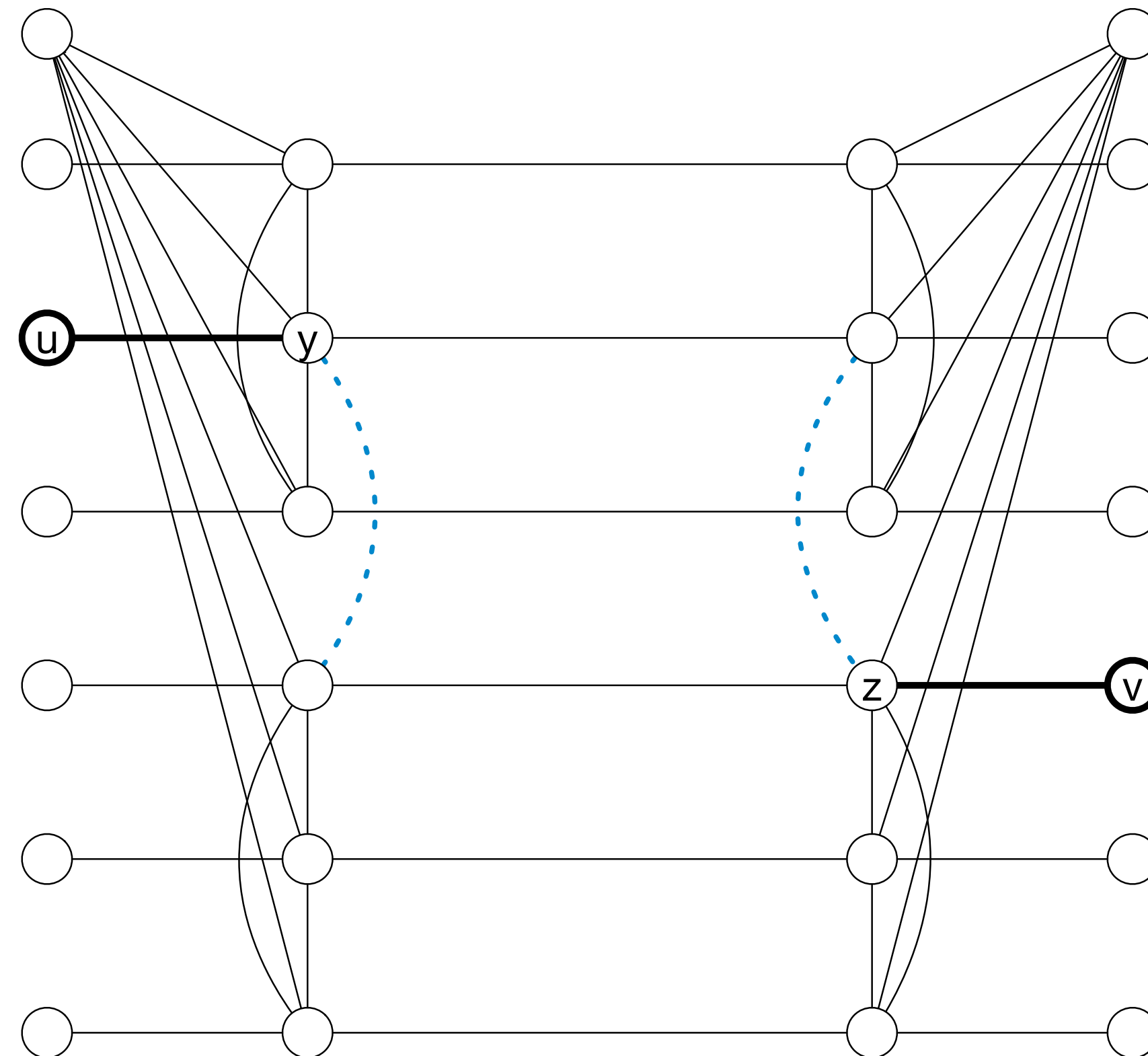


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
1
1
0
1
0



0
1
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1
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0

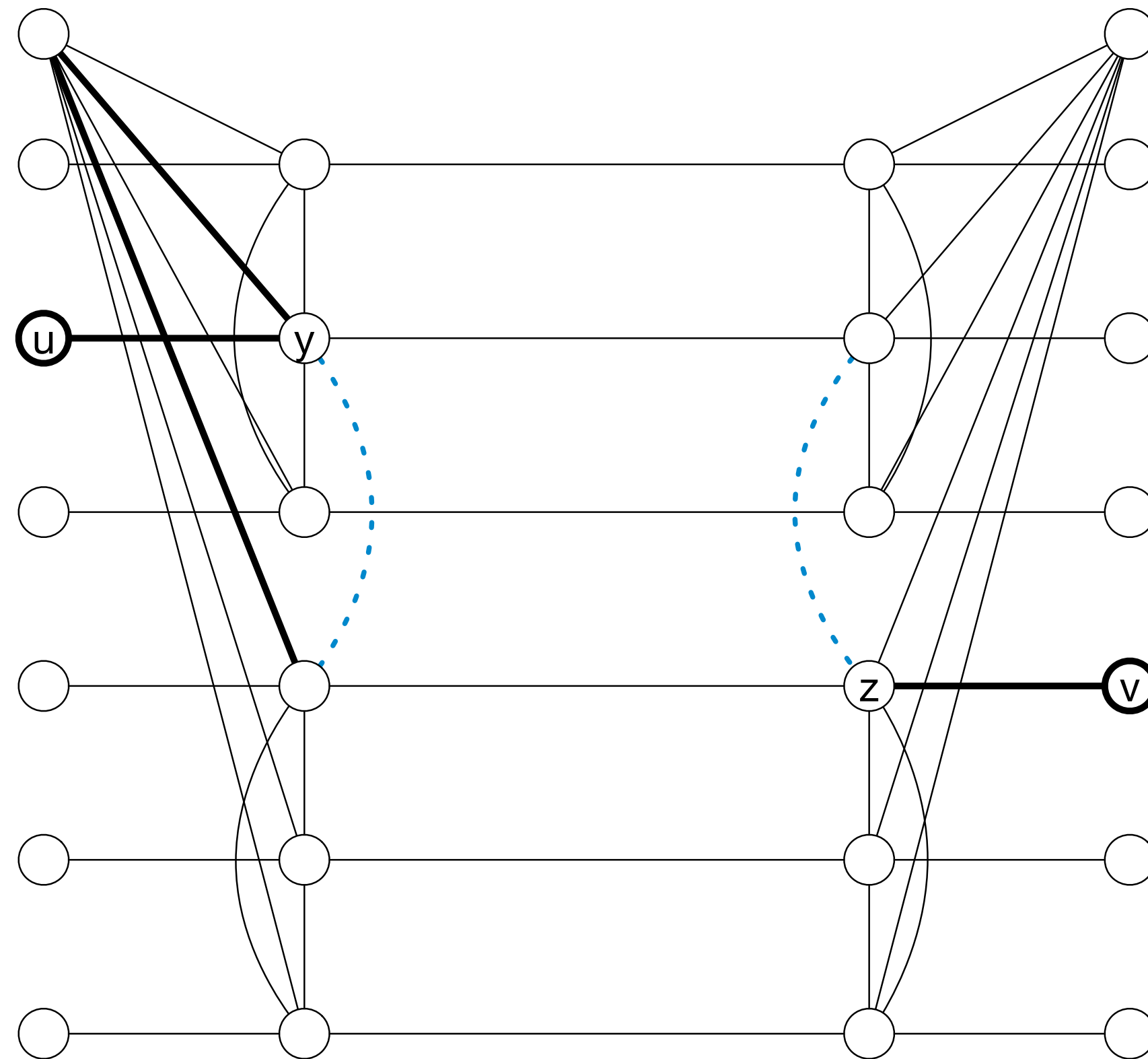


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
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1
0
1
0



0
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1
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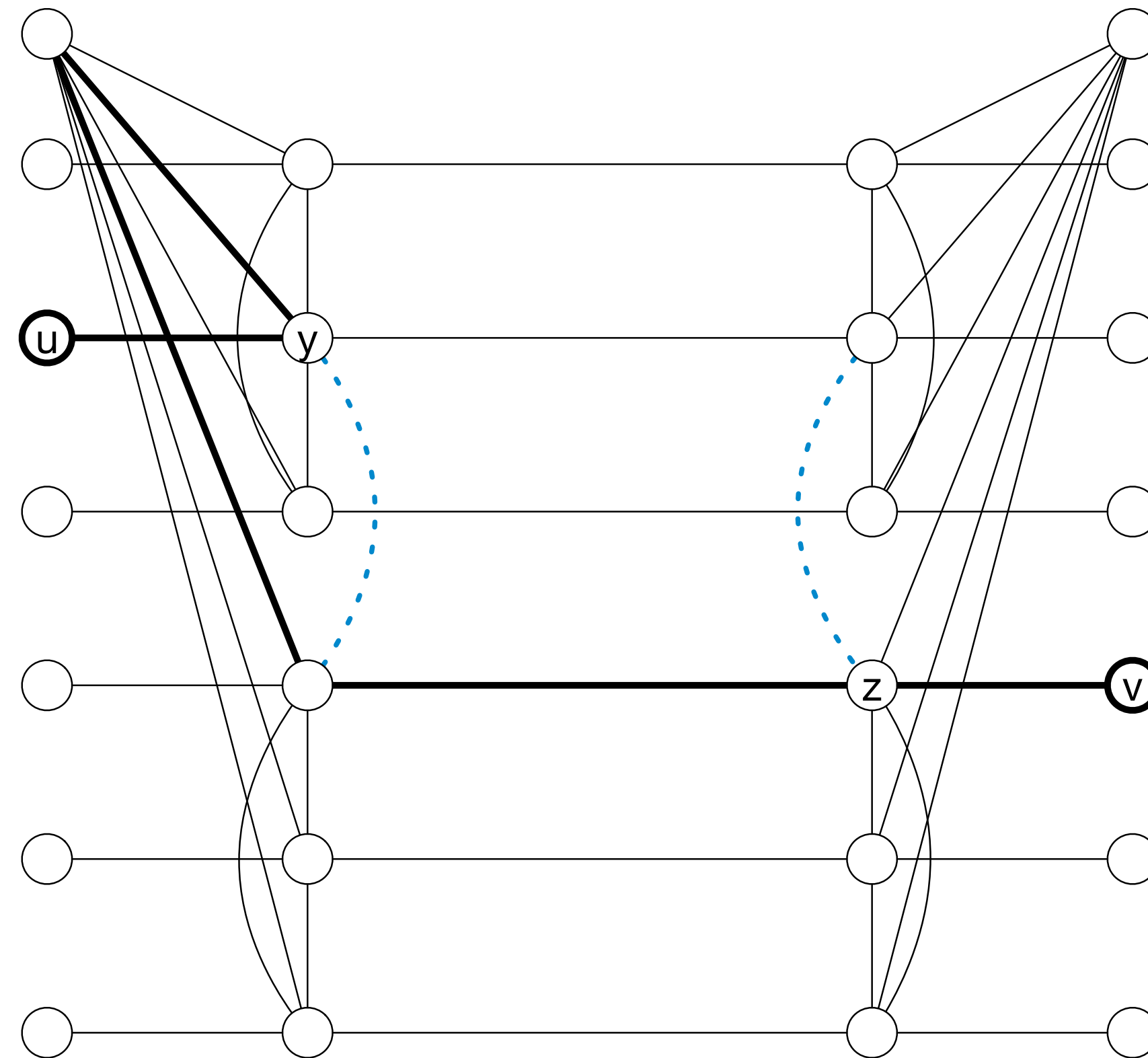


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
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0
1
0



0
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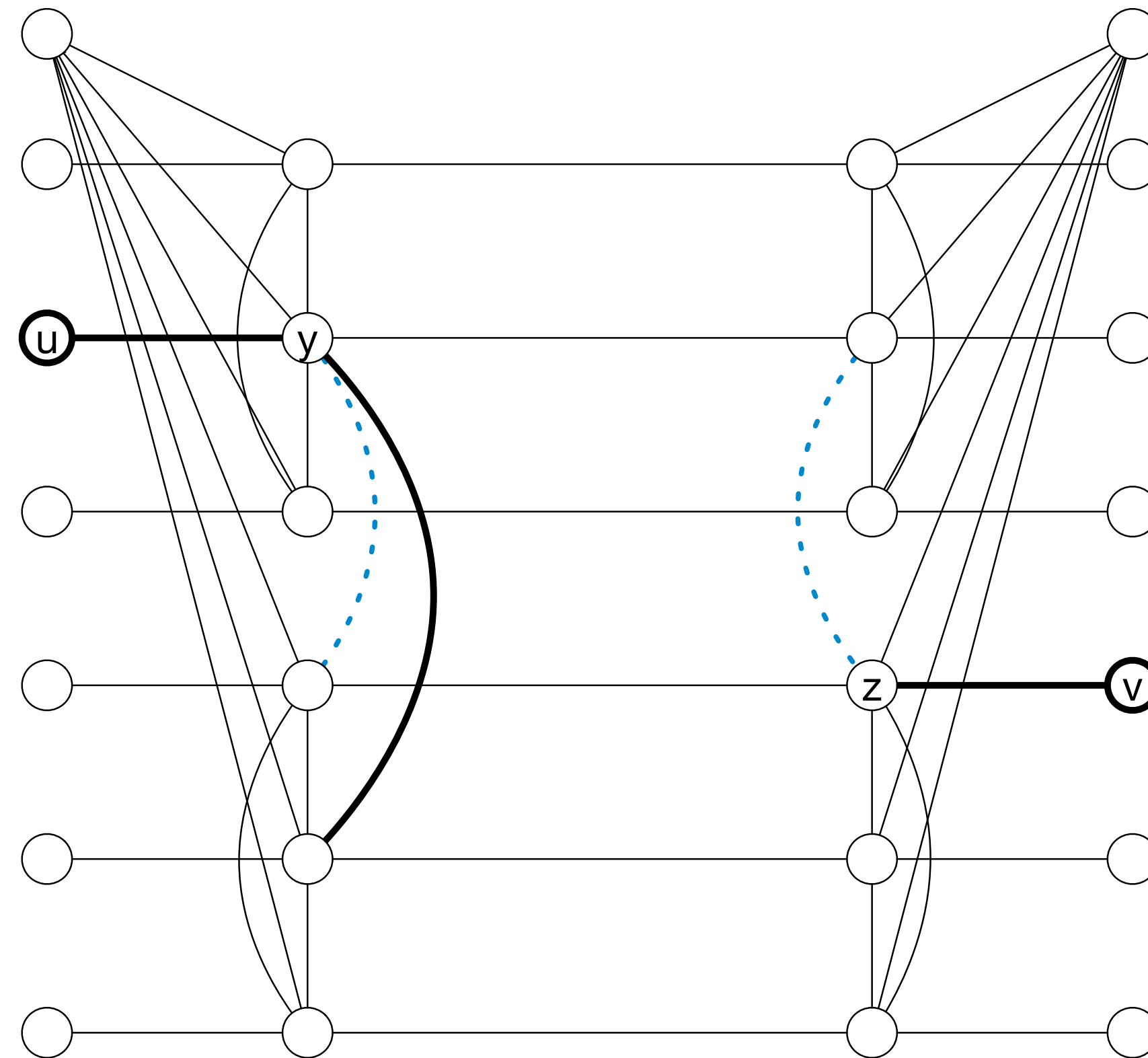


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
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0



0
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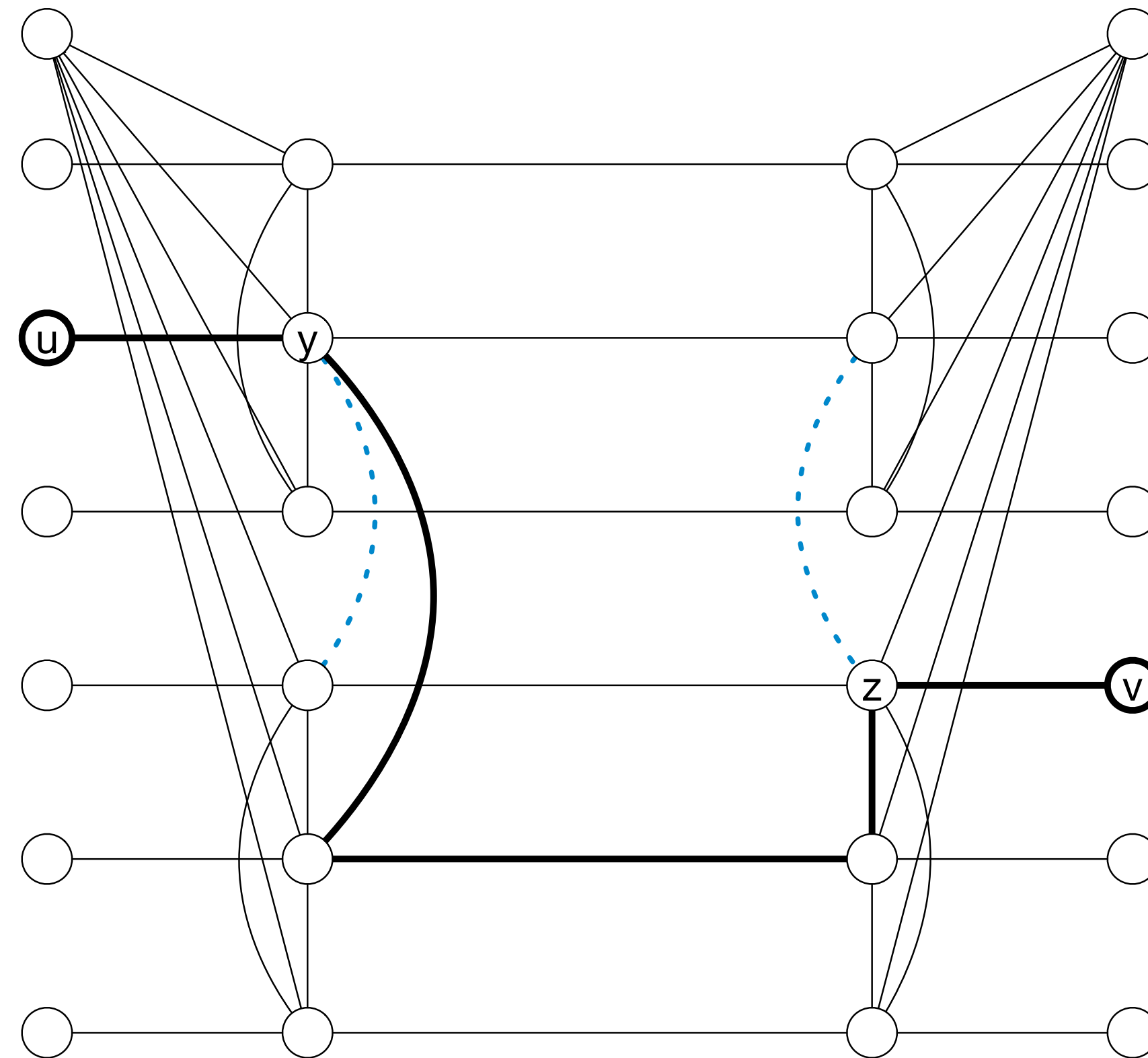


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
1
1
0
1
0

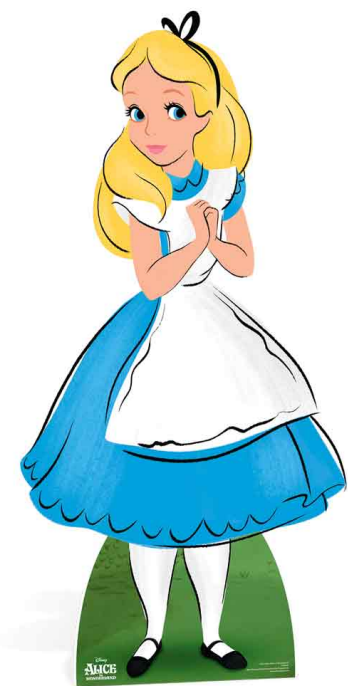


0
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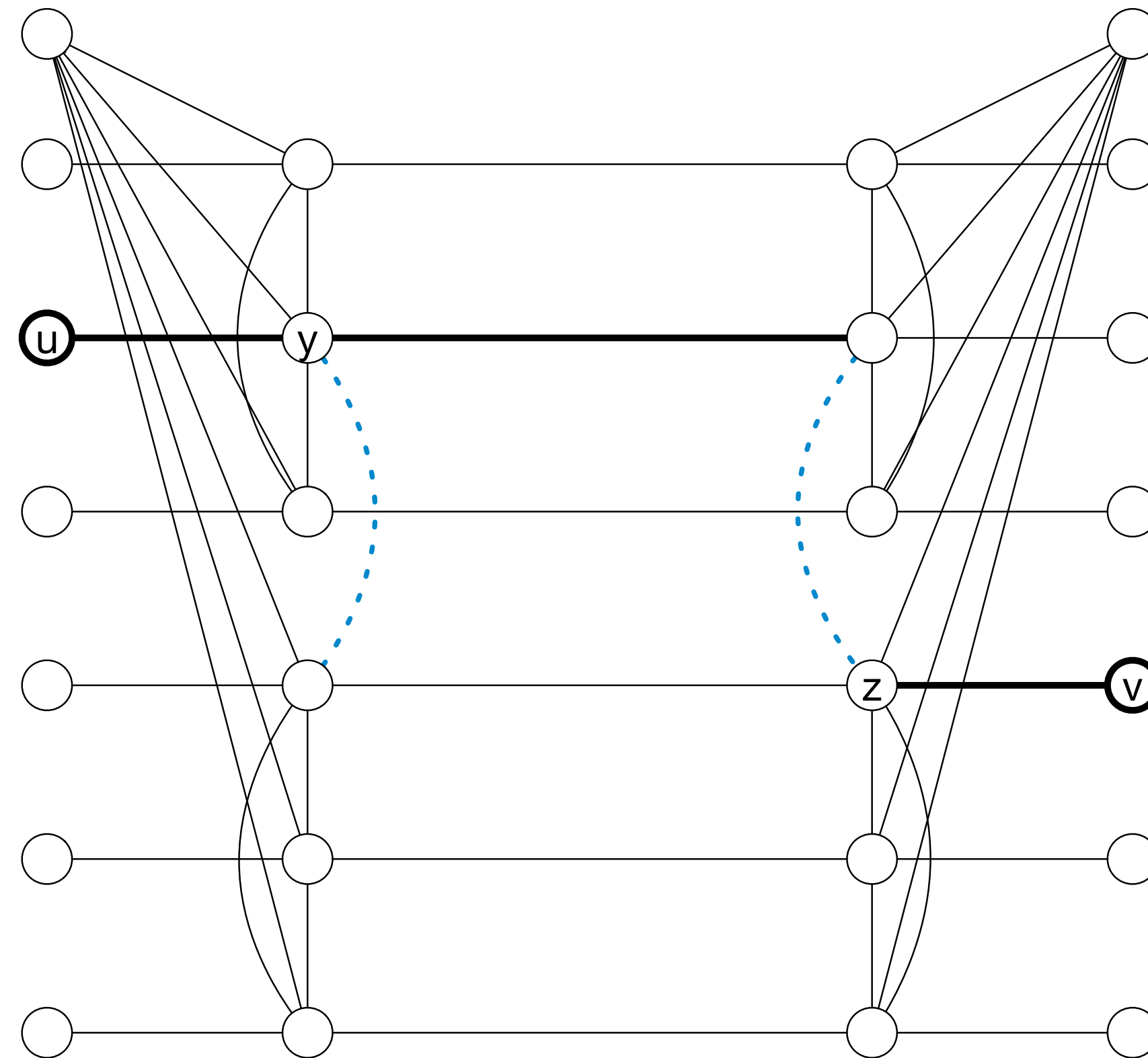


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
0
0
1
1
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0
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0

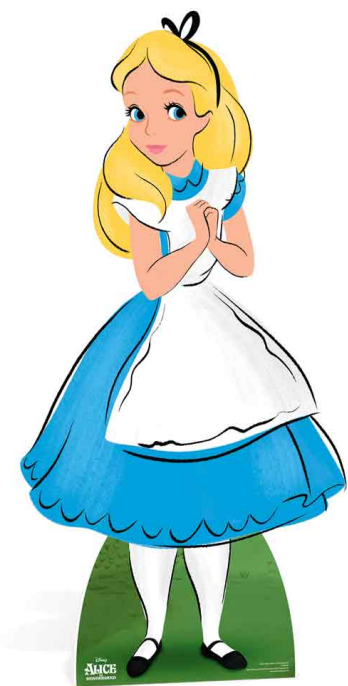


0
1
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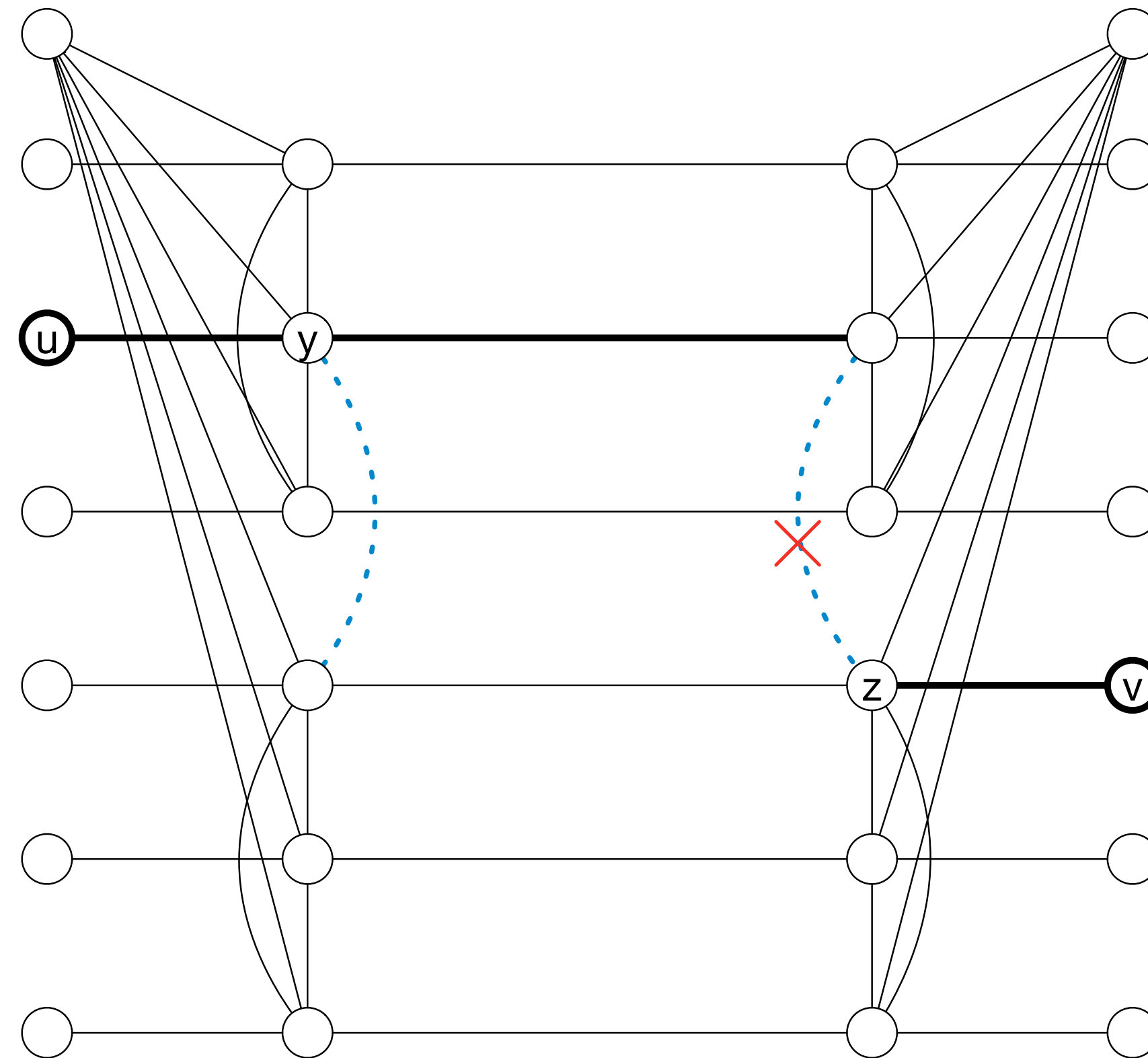


Lower bound for computing the diameter

Sets are **not** disjoint, **diameter ≥ 5**



1
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0
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0
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1
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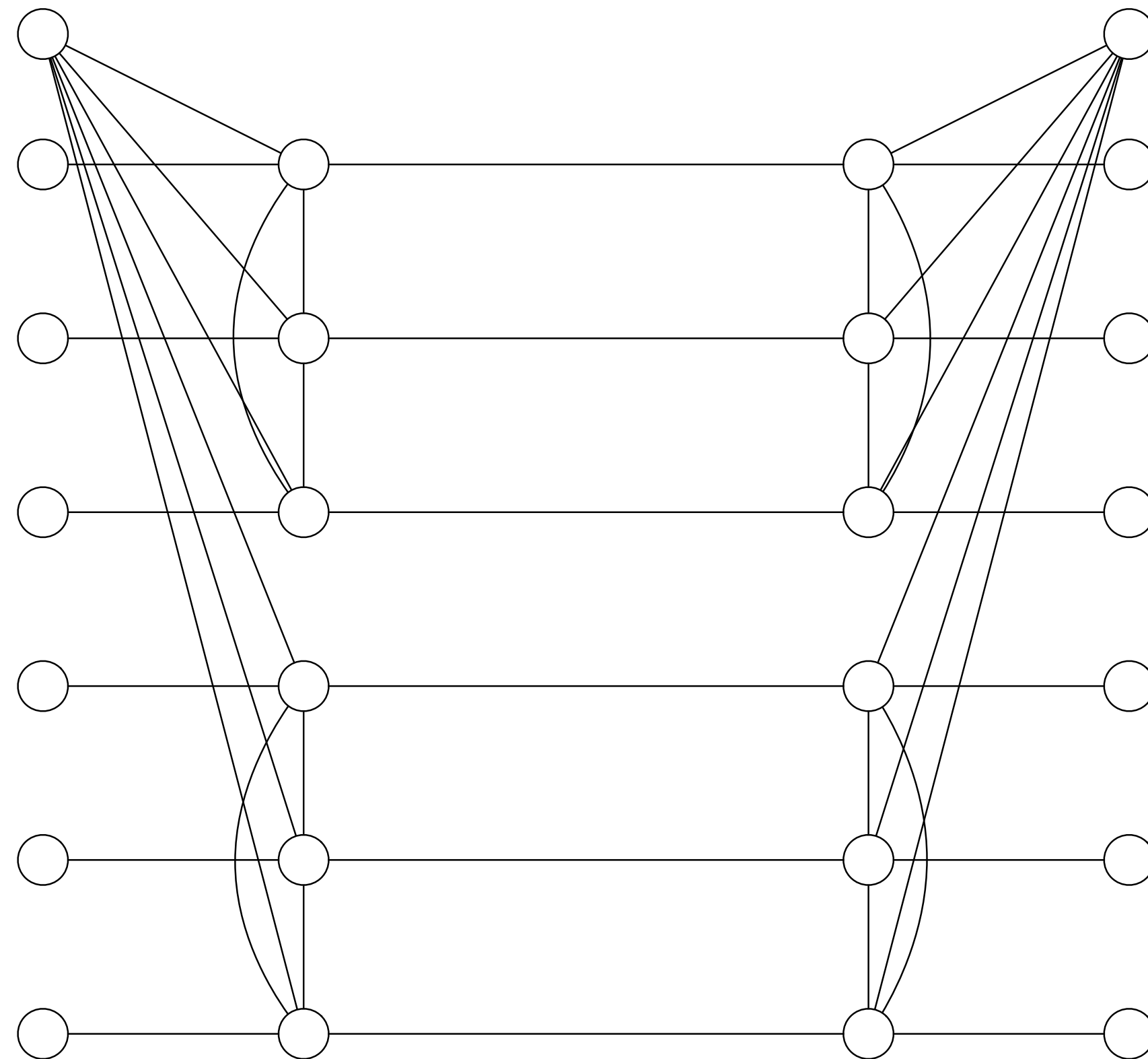


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
0
0
0
1
1
0
1
0



0
1
1
0
0
0
1
0
0

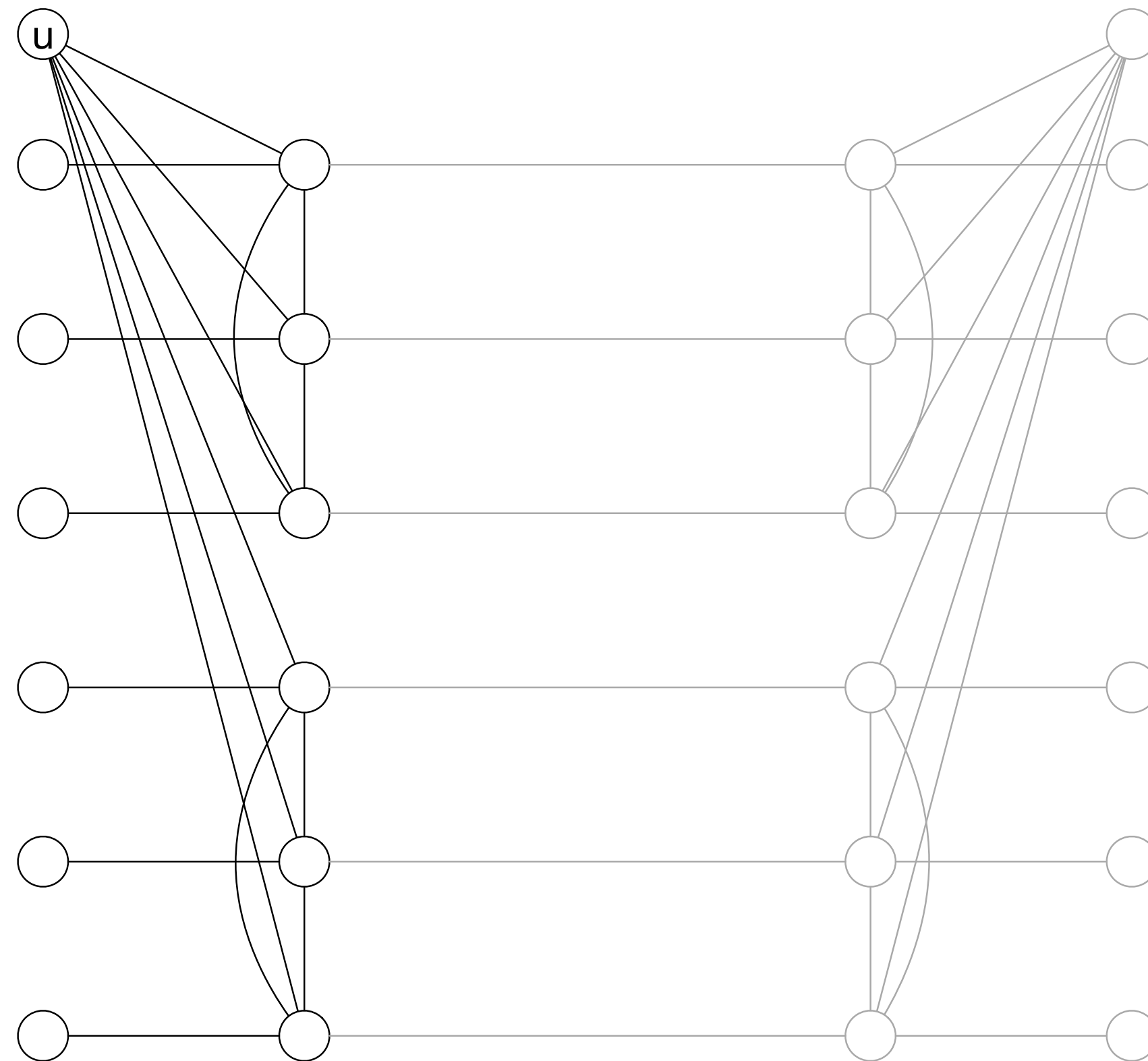


Lower bound for computing the diameter

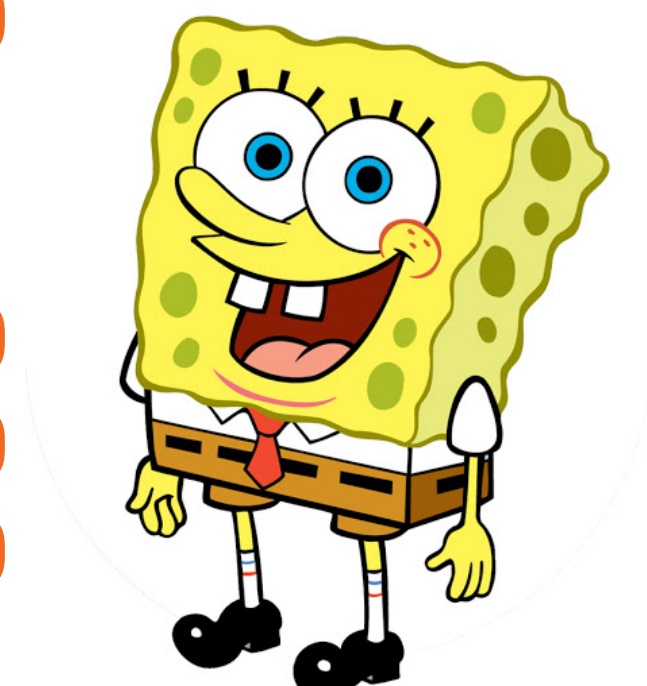
Sets are disjoint, **diameter = 4**



1
0
0
0
1
1
0
1
0

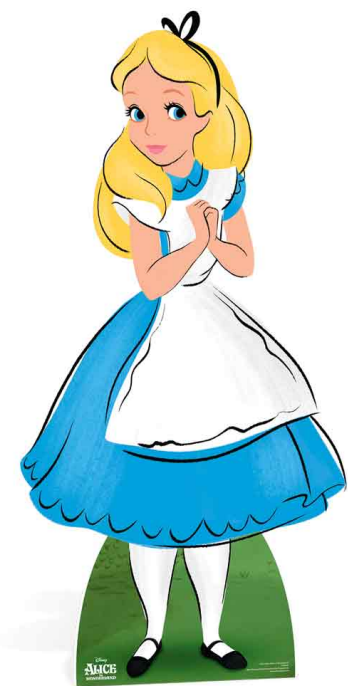


0
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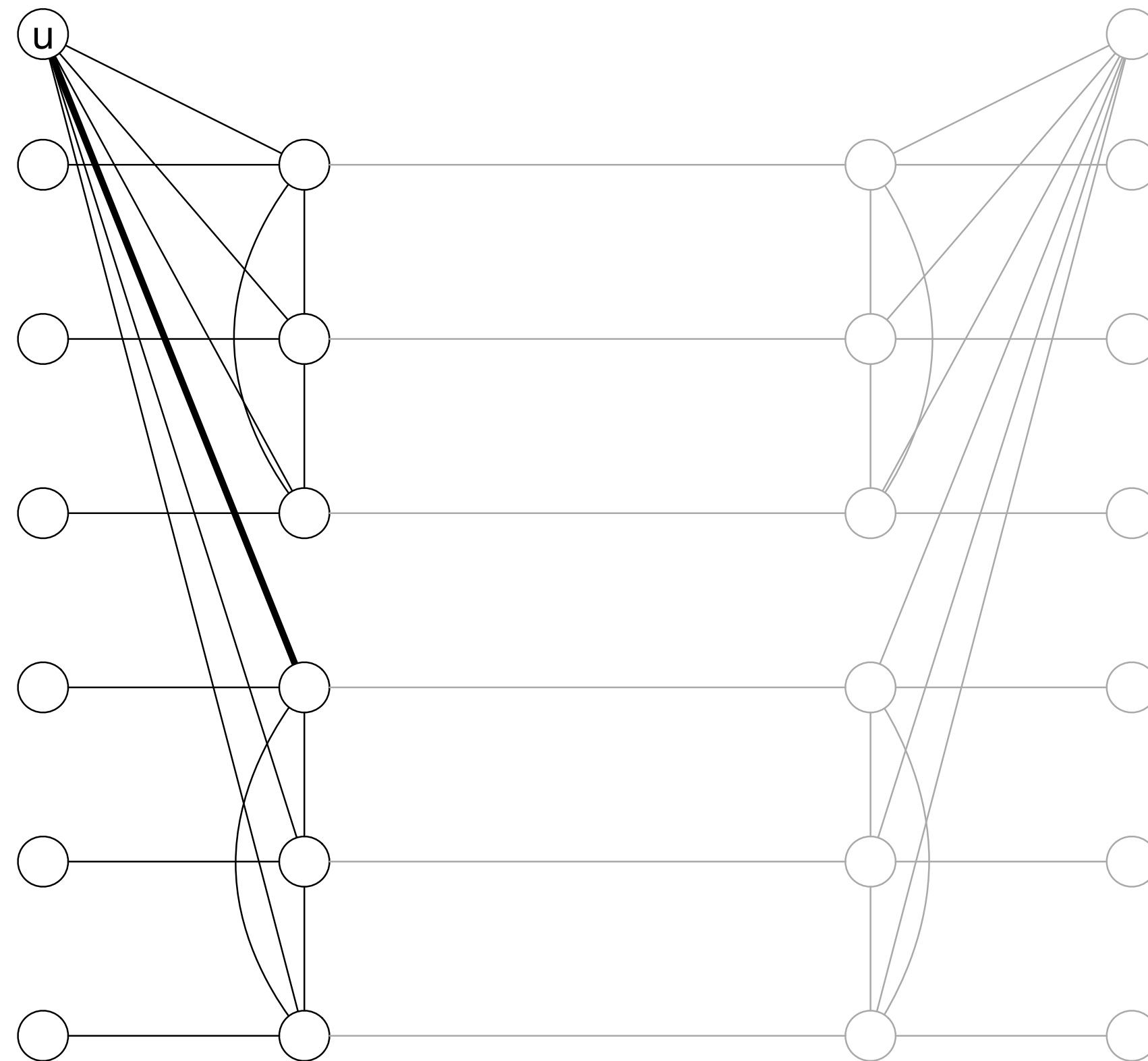


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
0
0
0
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0
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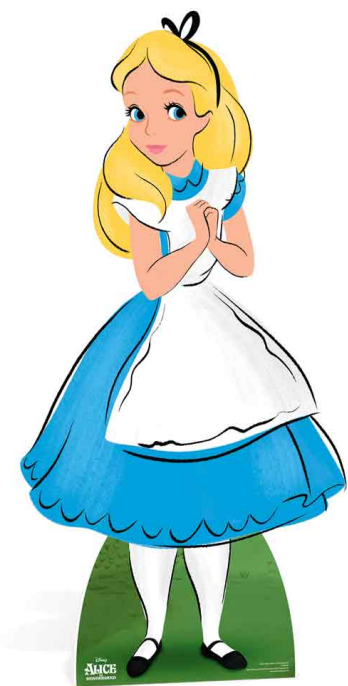


0
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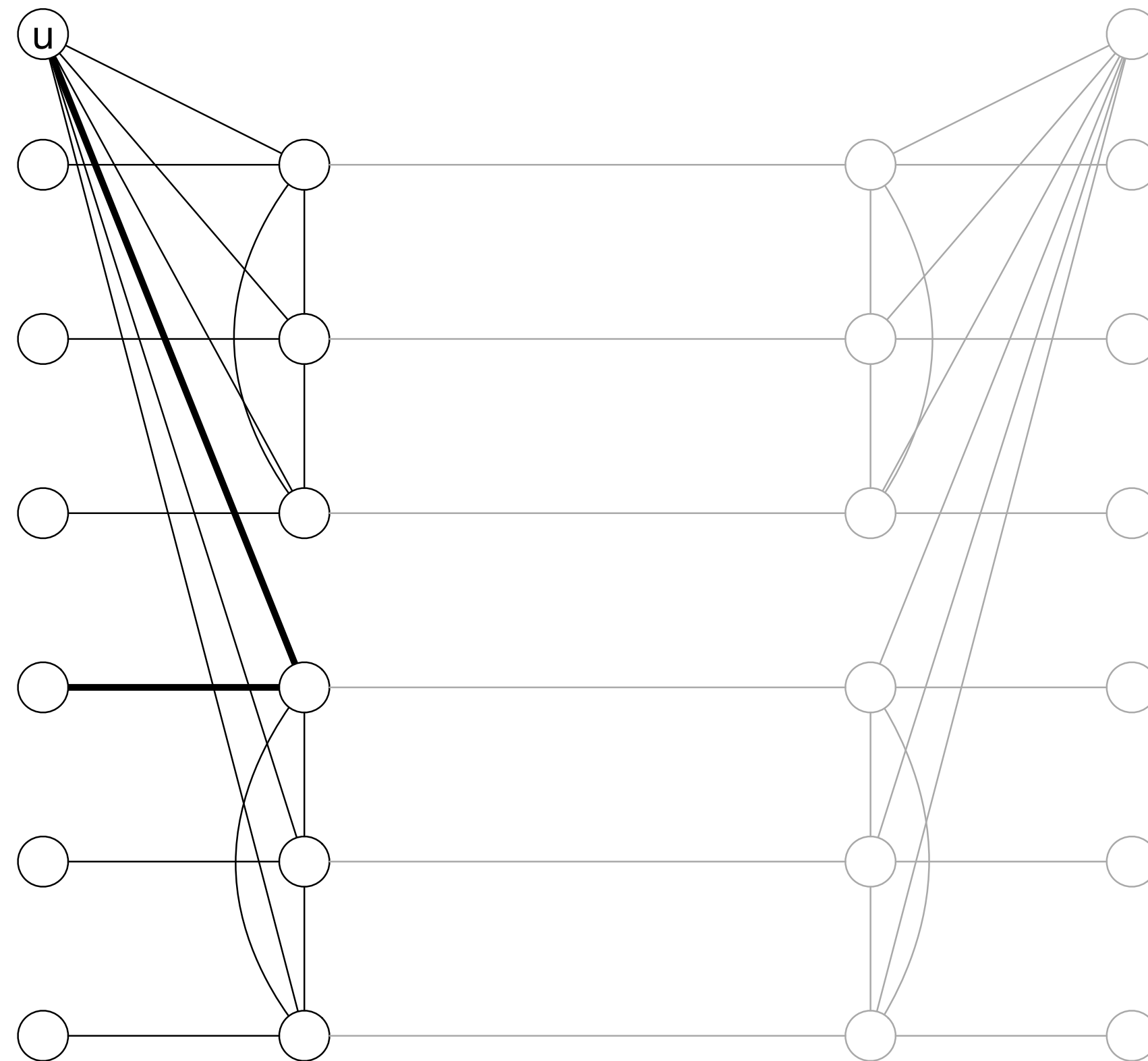


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
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0
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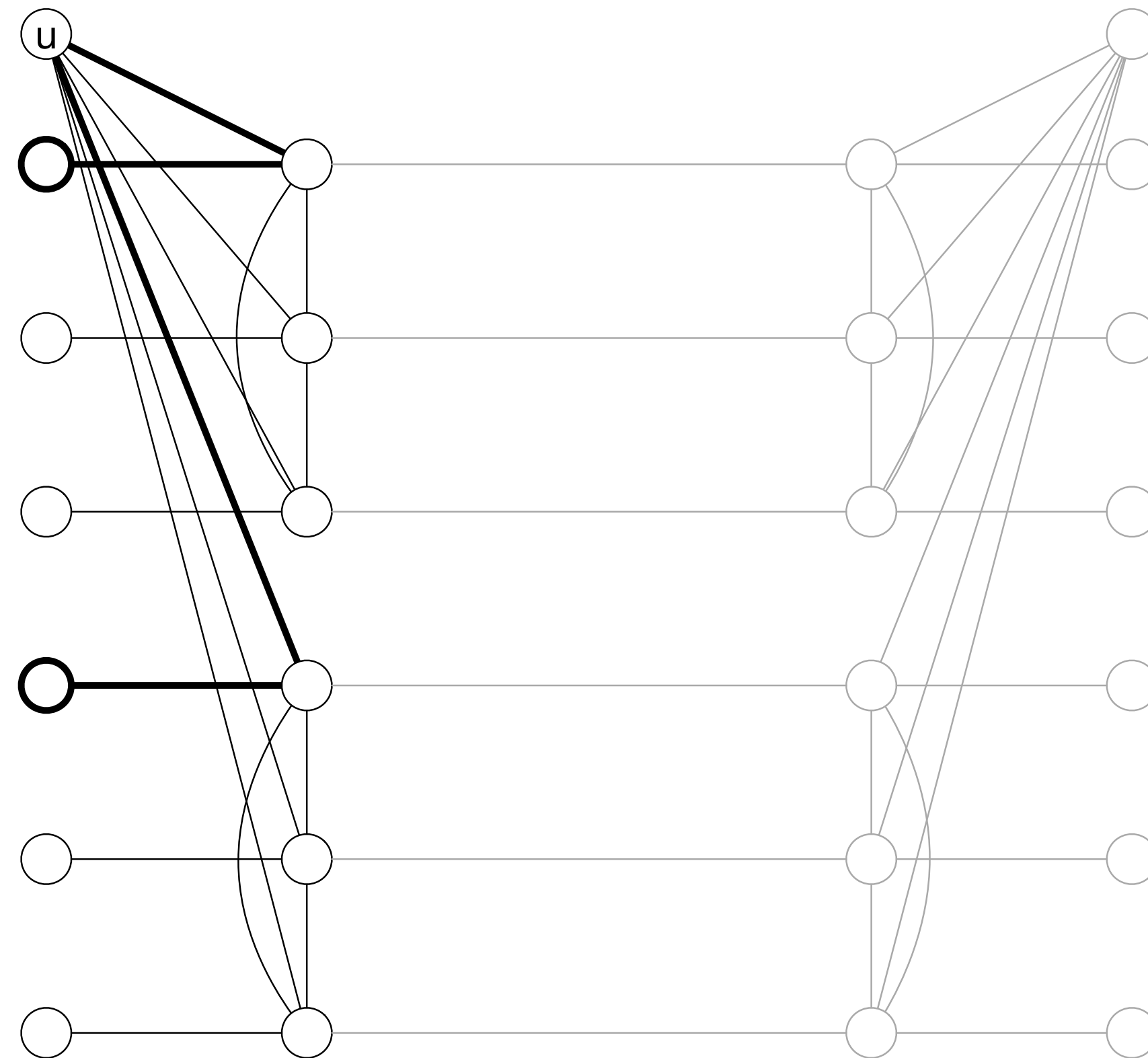


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
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1
0



0
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0
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0

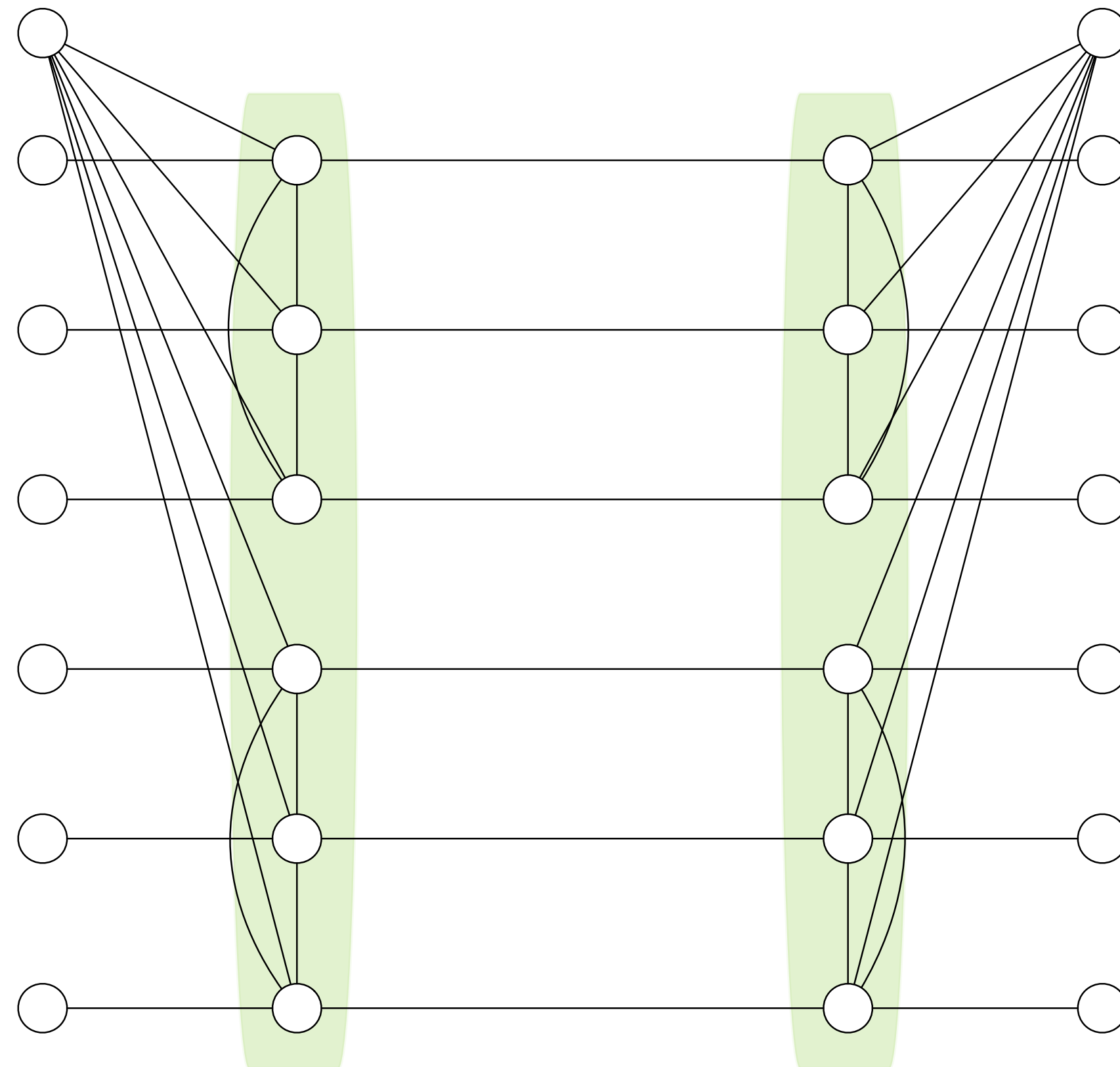


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
0
0
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0
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0



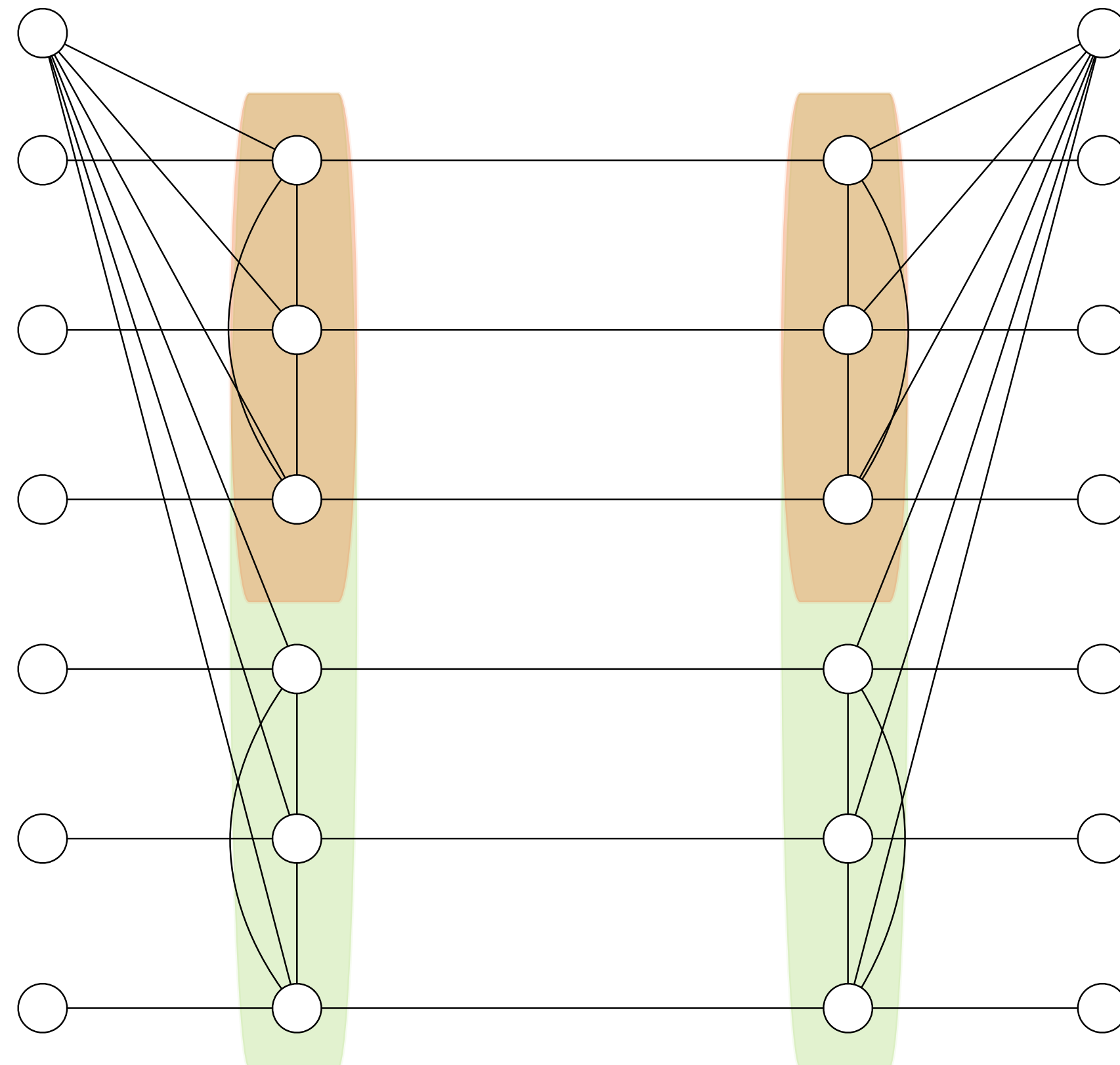
0
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Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
0
0
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0
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0

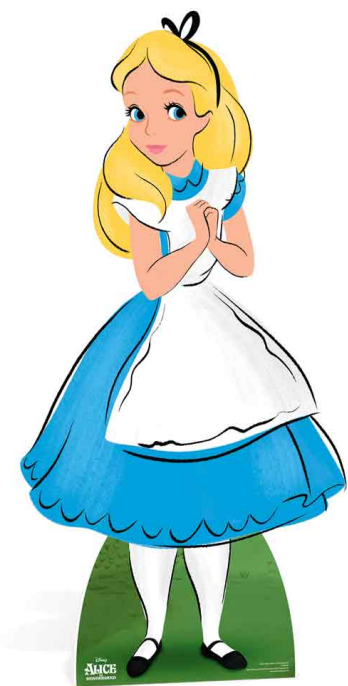


0
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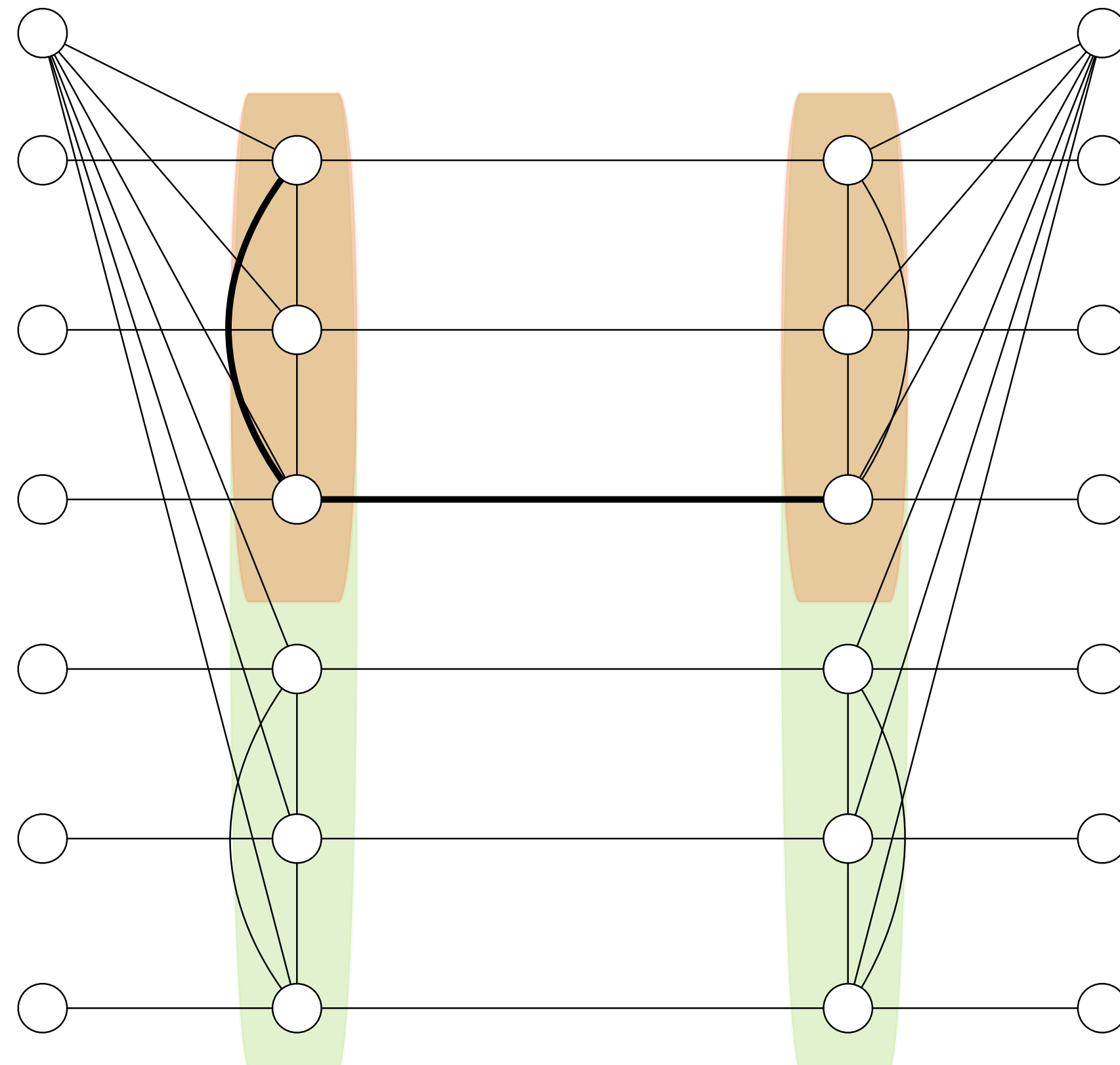


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
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0
1
0



0
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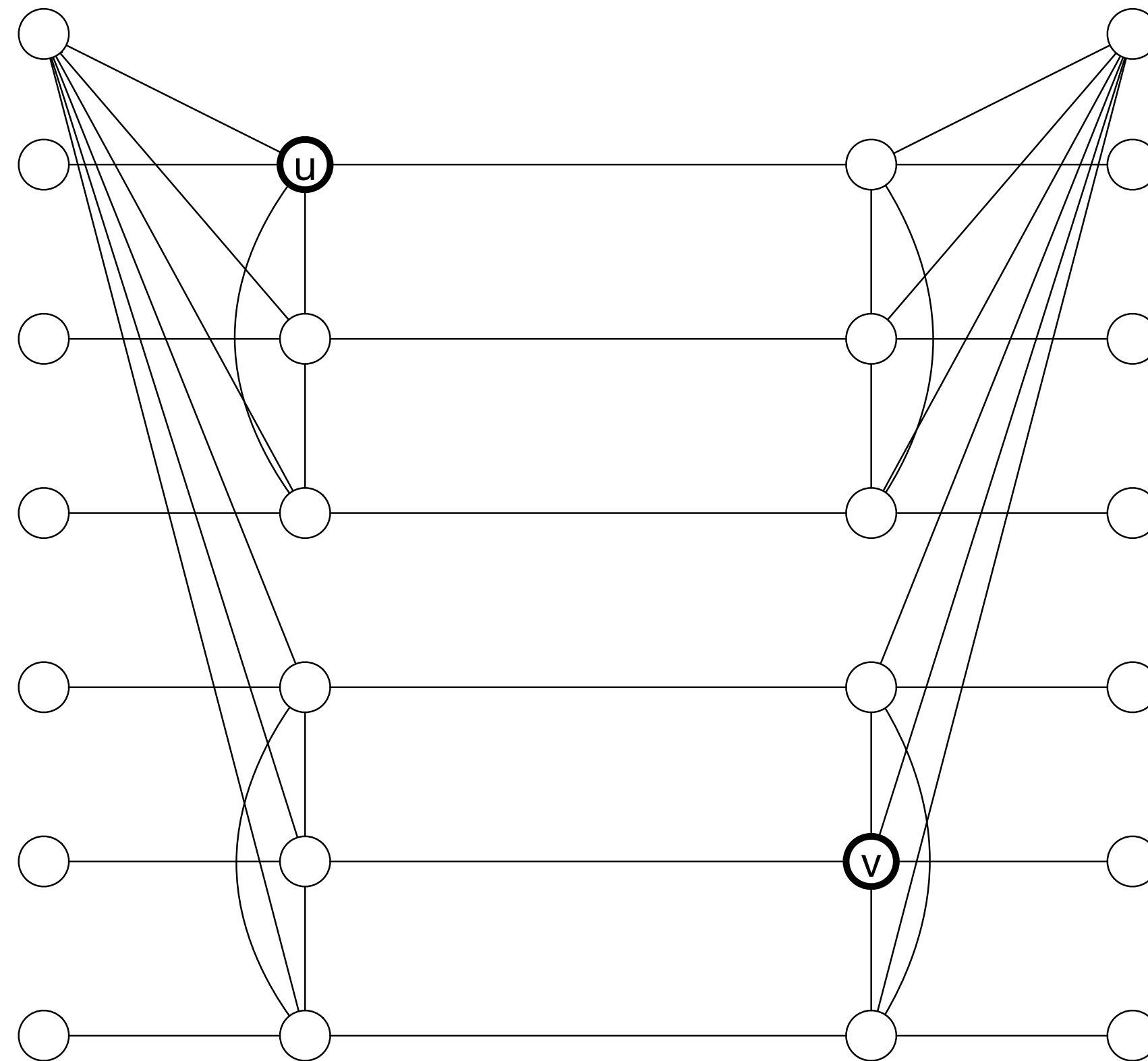


Lower bound for computing the diameter

Sets are disjoint, **diameter = 4**



1
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0



0
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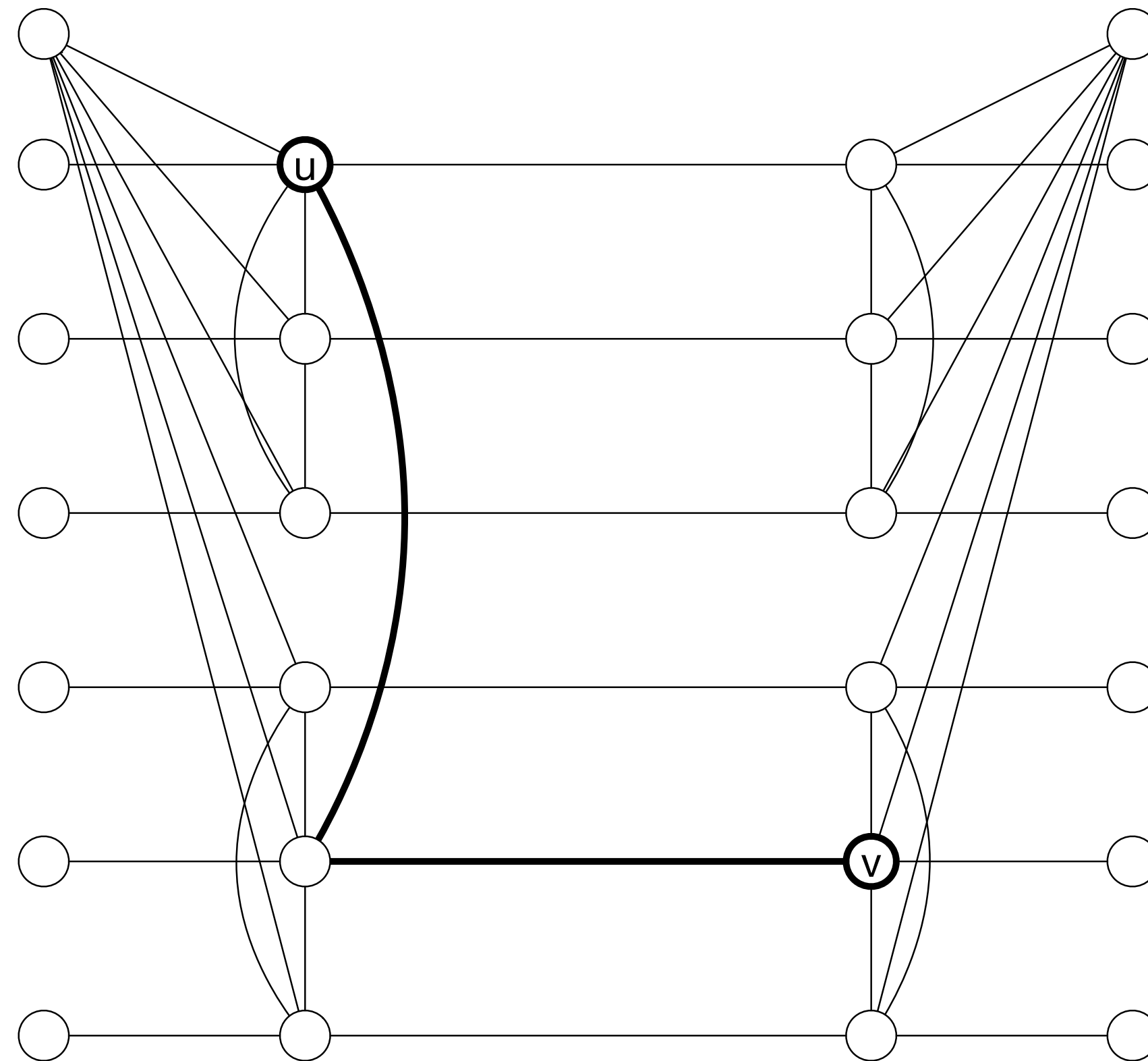


Lower bound for computing the diameter

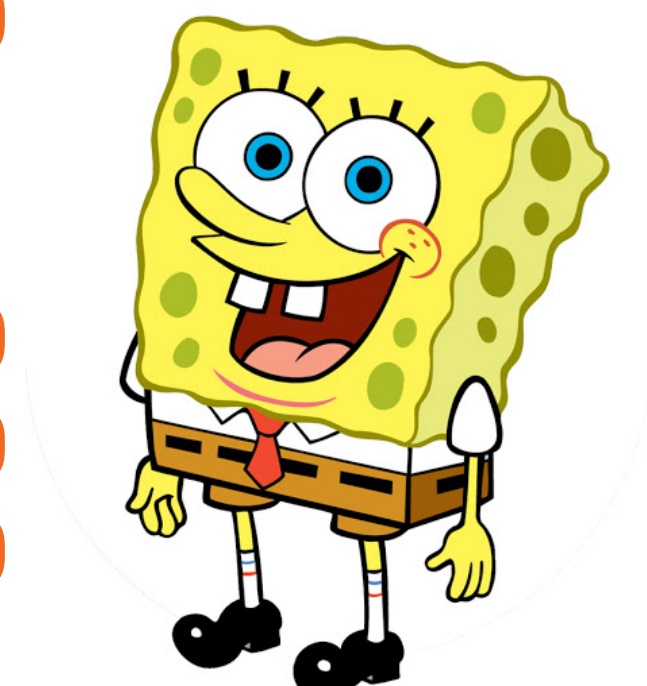
Sets are disjoint, **diameter = 4**



1
0
0
0
1
1
0
1
0



0
1
1
0
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1
0
0

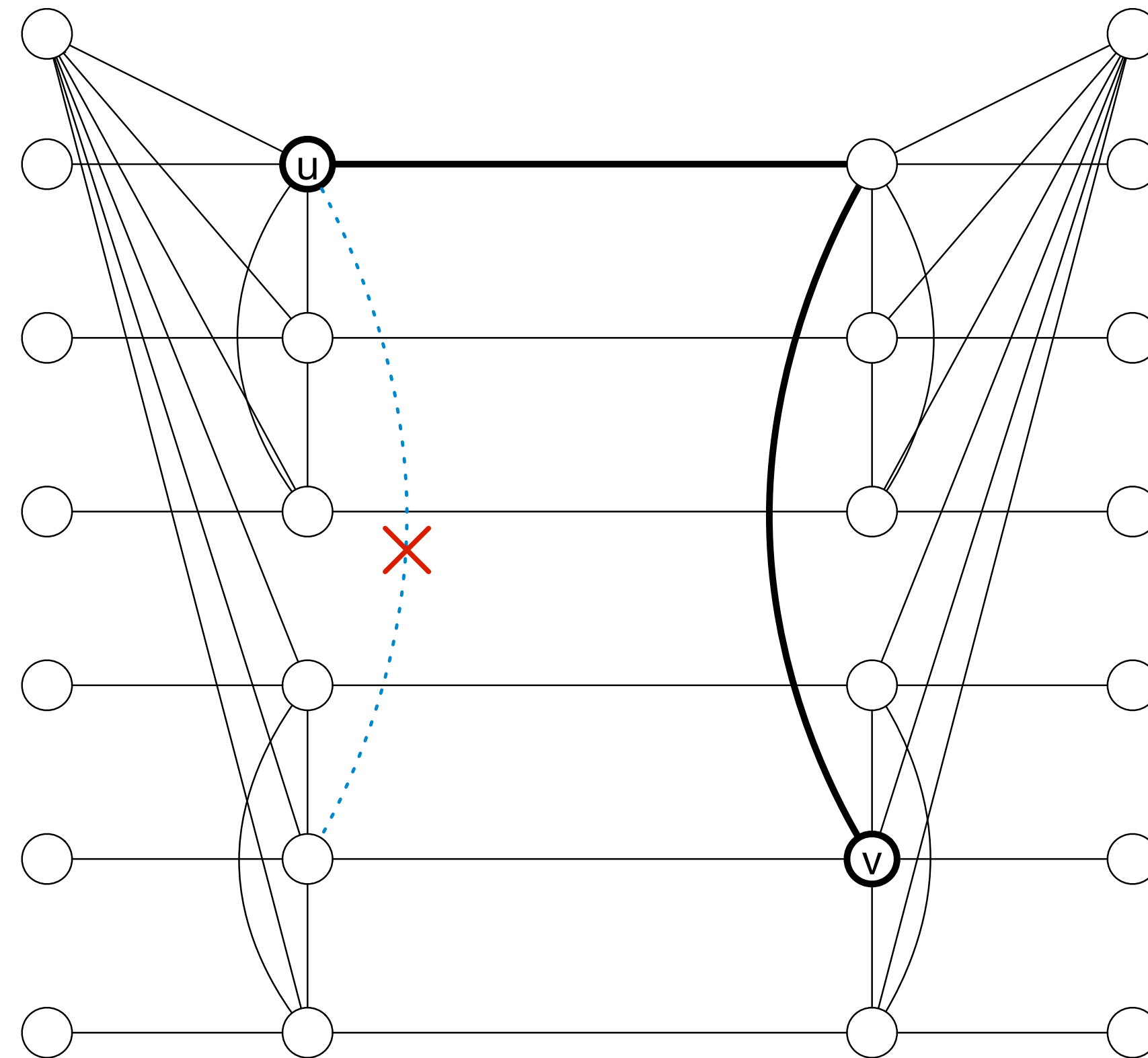


Lower bound for computing the diameter

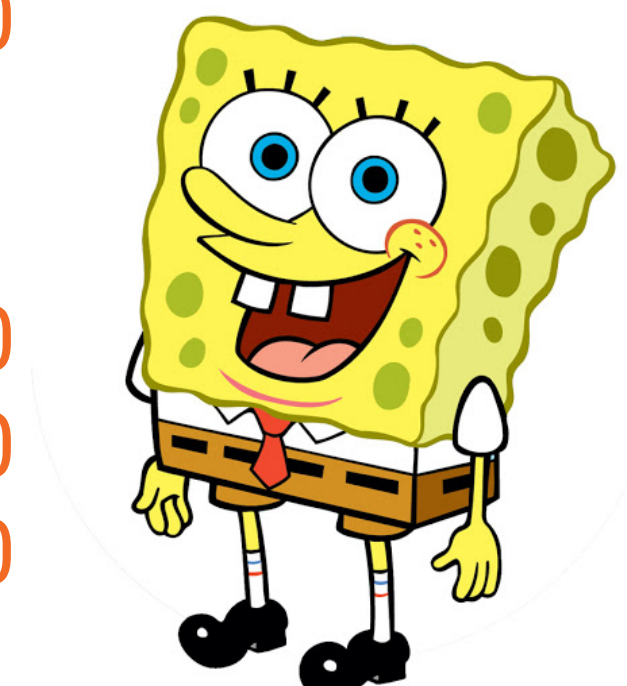
Sets are disjoint, **diameter = 4**



1
0
0
0
1
1
0
1
0

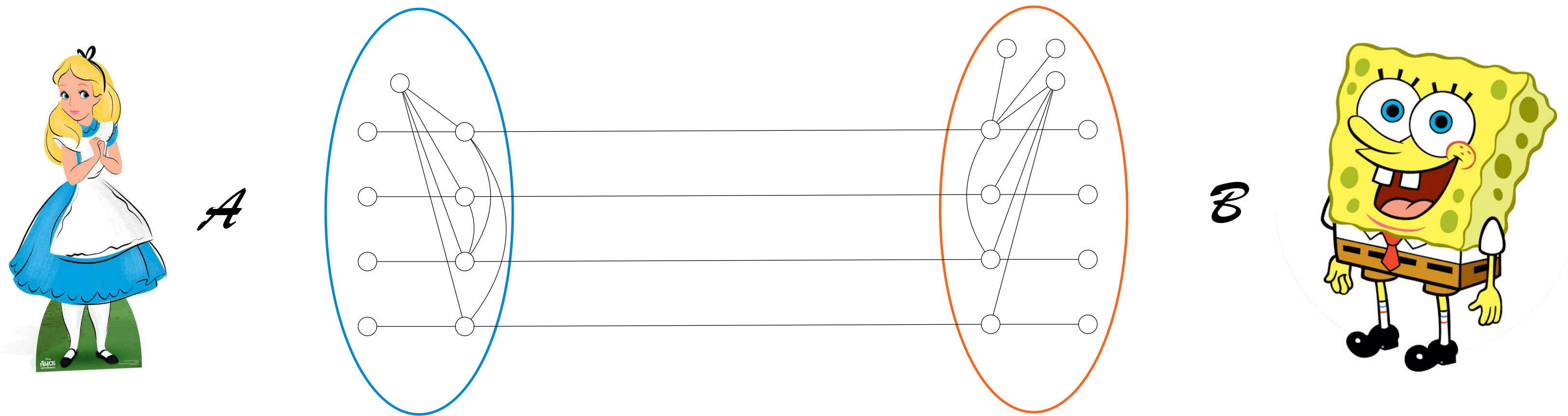


0
1
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0



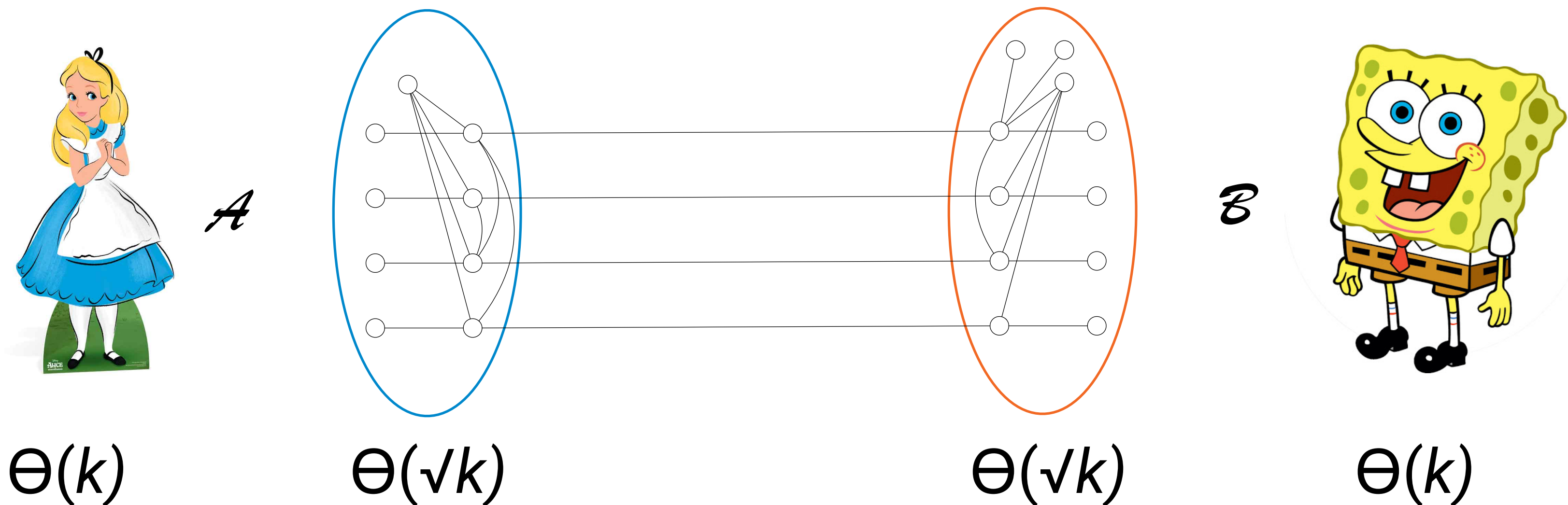
Lower bound for computing the diameter

- Diameter = **4** \Rightarrow sets are **disjoint**
- Diameter \geq **5** \Rightarrow are **not** disjoint



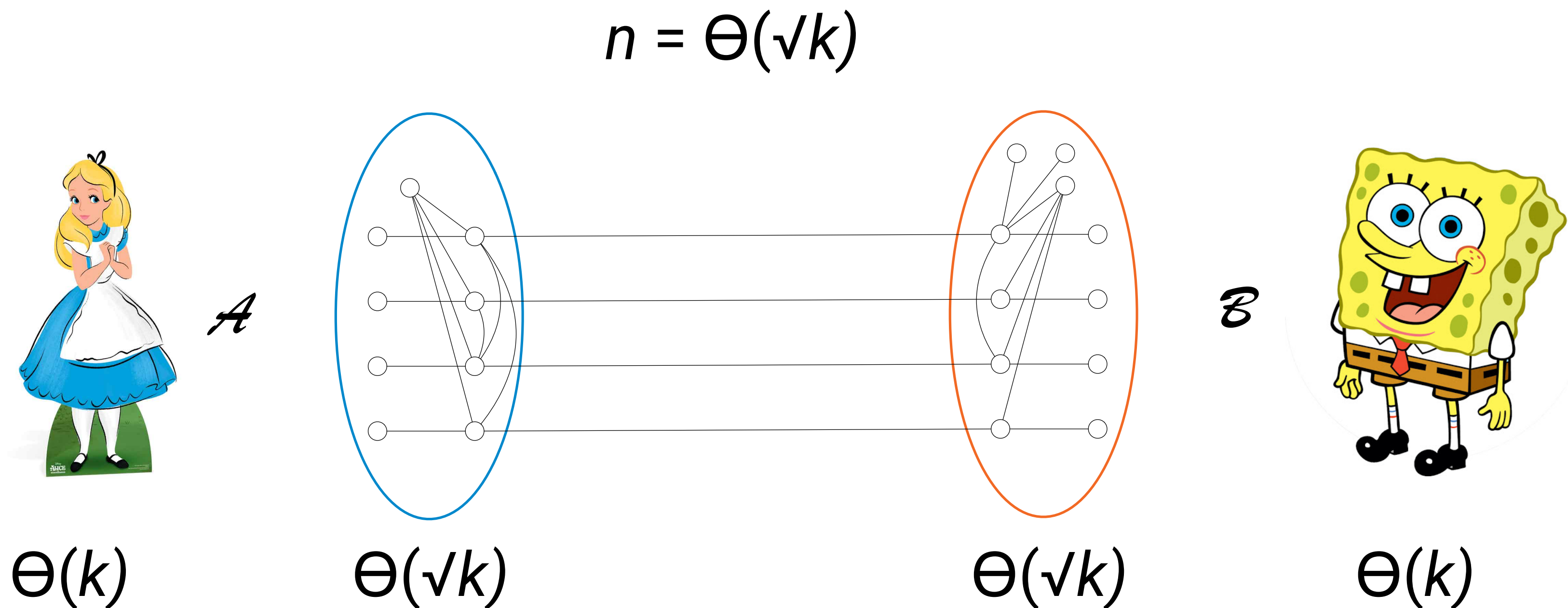
Lower bound for computing the diameter

- Diameter = **4** \Rightarrow sets are **disjoint**
- Diameter \geq **5** \Rightarrow are **not** disjoint



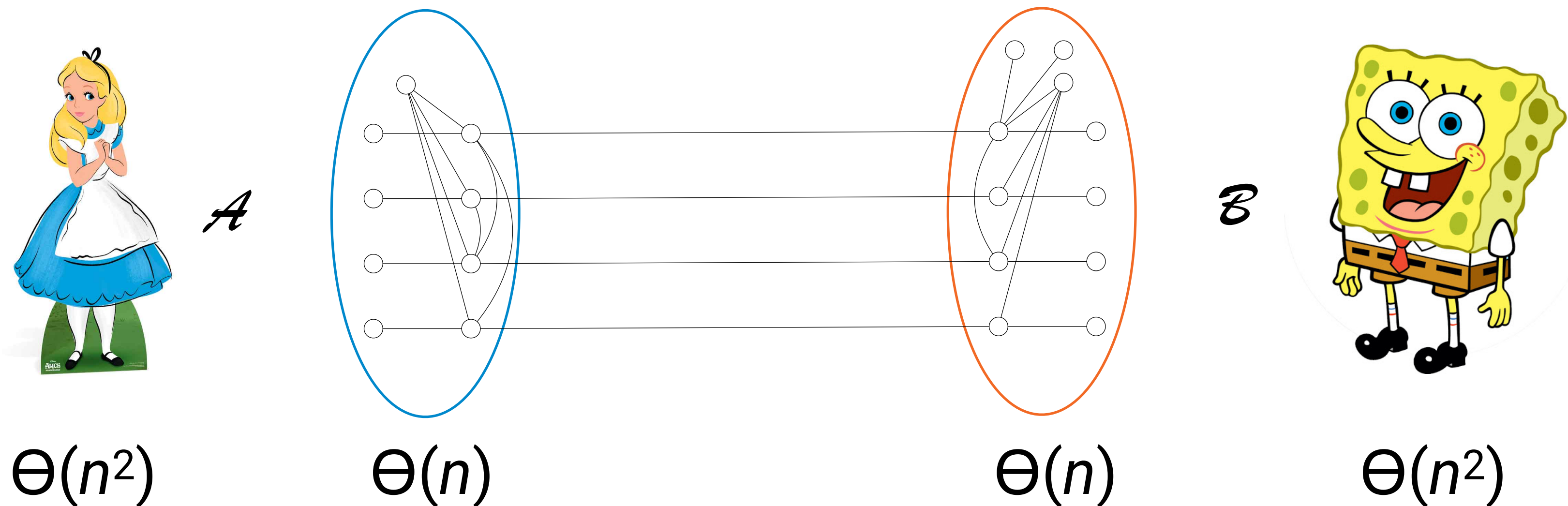
Lower bound for computing the diameter

- Diameter = **4** \Rightarrow sets are **disjoint**
- Diameter \geq **5** \Rightarrow are **not** disjoint



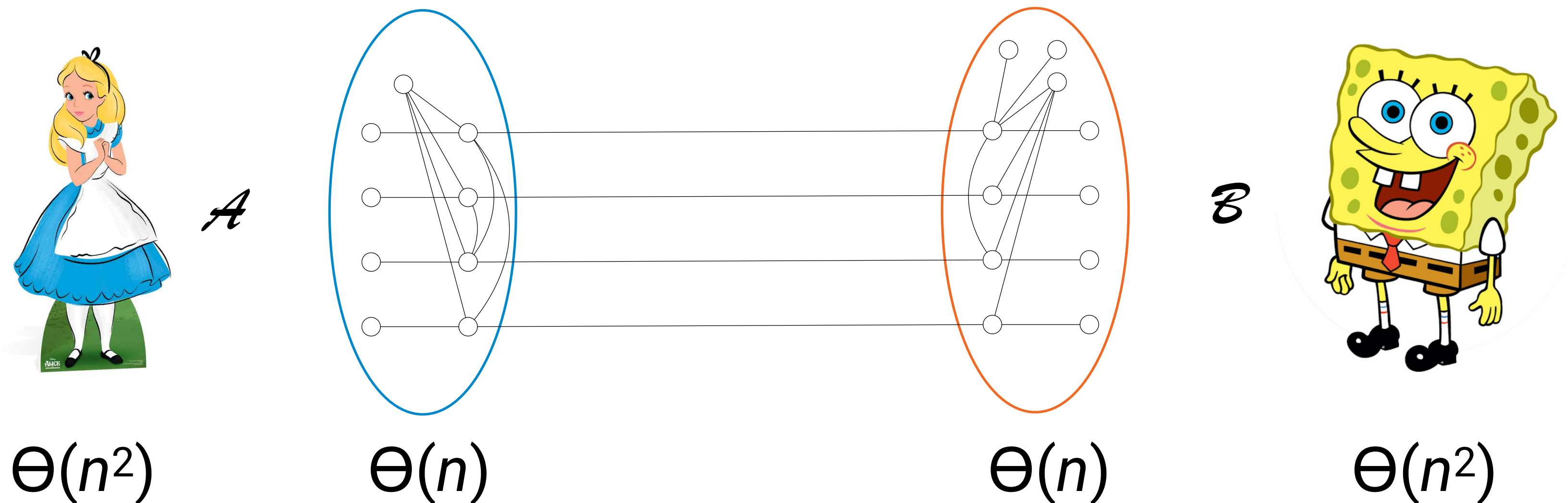
Lower bound for computing the diameter

- Diameter = **4** \Rightarrow sets are **disjoint**
- Diameter \geq **5** \Rightarrow are **not** disjoint



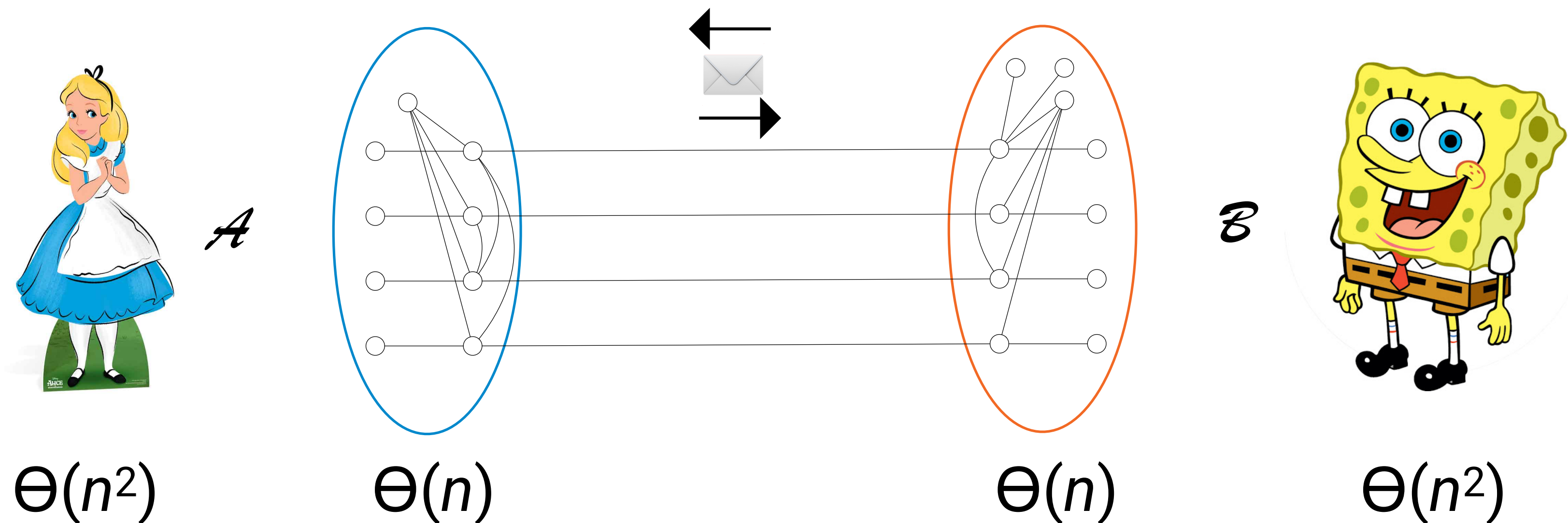
Lower bound for computing the diameter

Suppose we have an **algorithm A** for computing the diameter in time $T(A, n)$



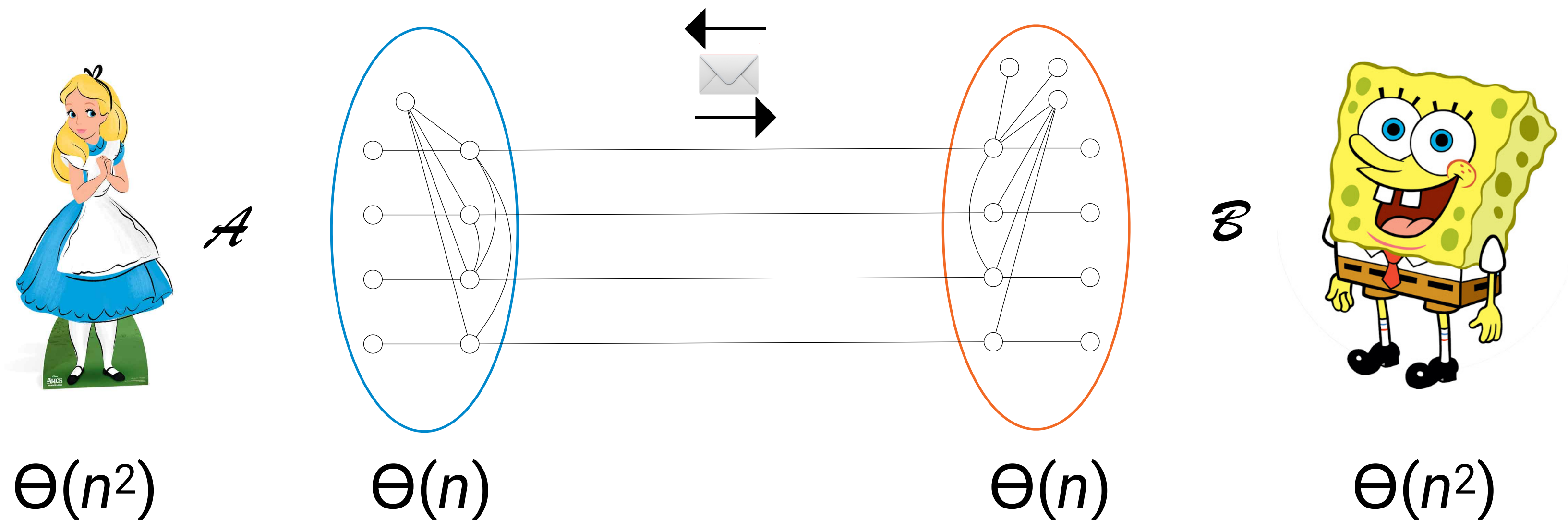
Lower bound for computing the diameter

- Simulate **A** \Rightarrow solve **set disjointness**



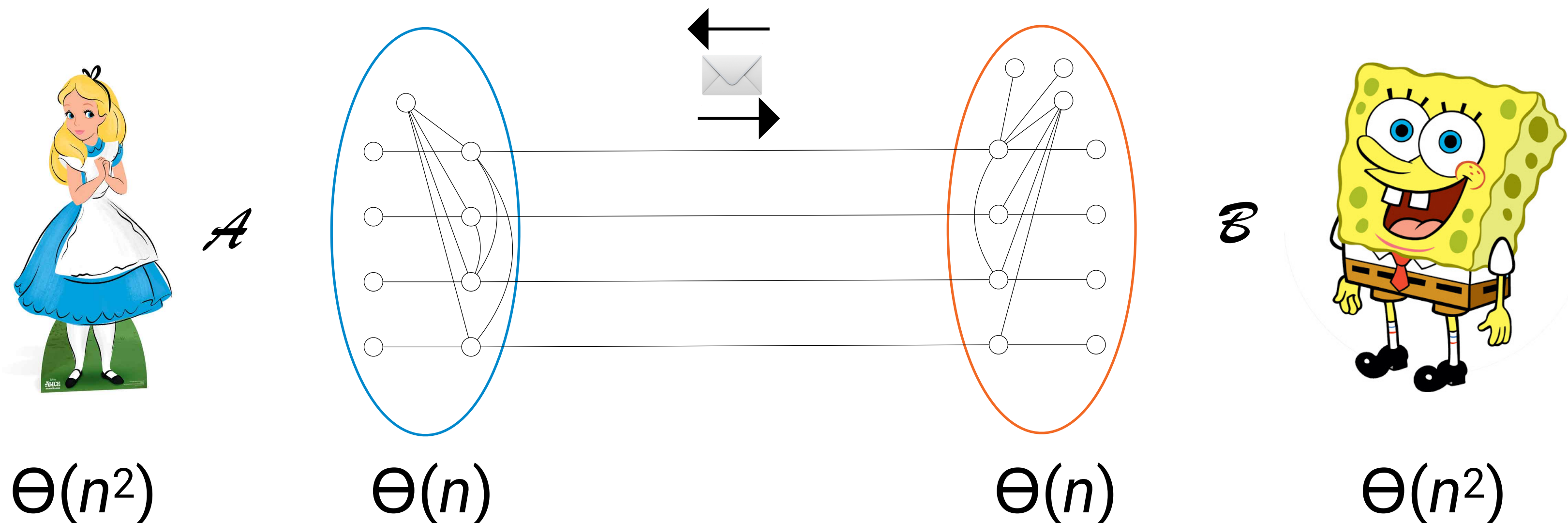
Lower bound for computing the diameter

- Simulate **A** \Rightarrow solve **set disjointness**
- **1 round** of simulation of A: exchange $\Theta(n \log n)$ bits



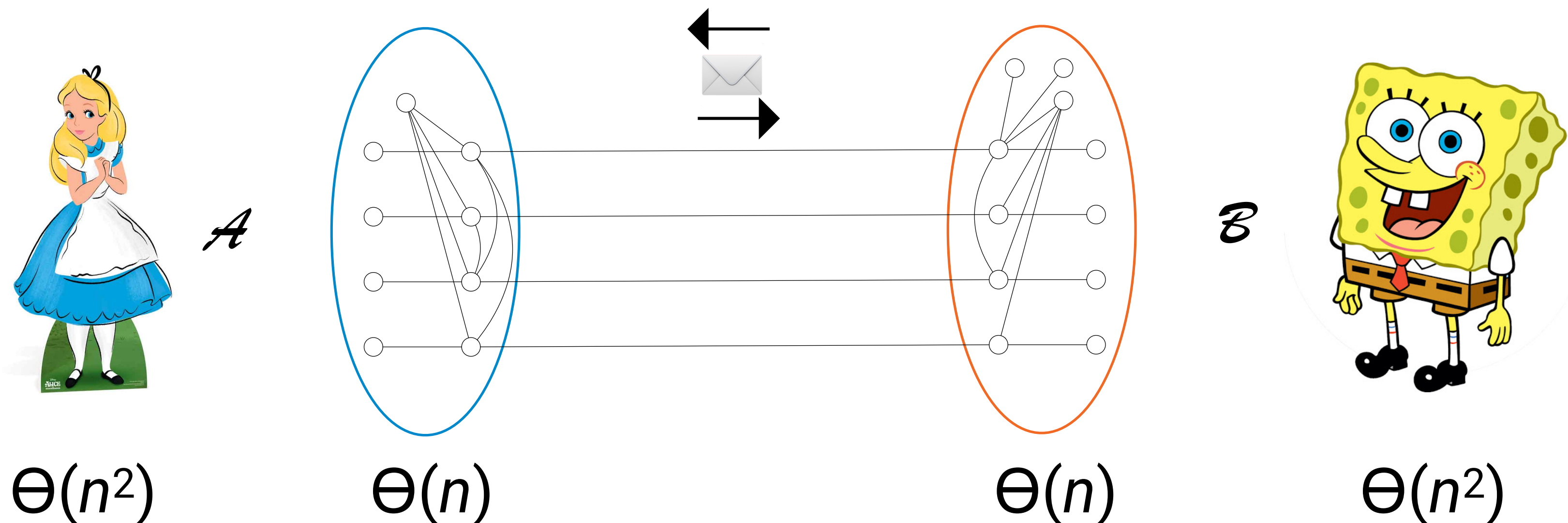
Lower bound for computing the diameter

- Simulate **A** \Rightarrow solve **set disjointness**
- **1 round** of simulation of A : exchange $\Theta(n \log n)$ bits
- Total: $T(A, n) \times \Theta(n \log n)$



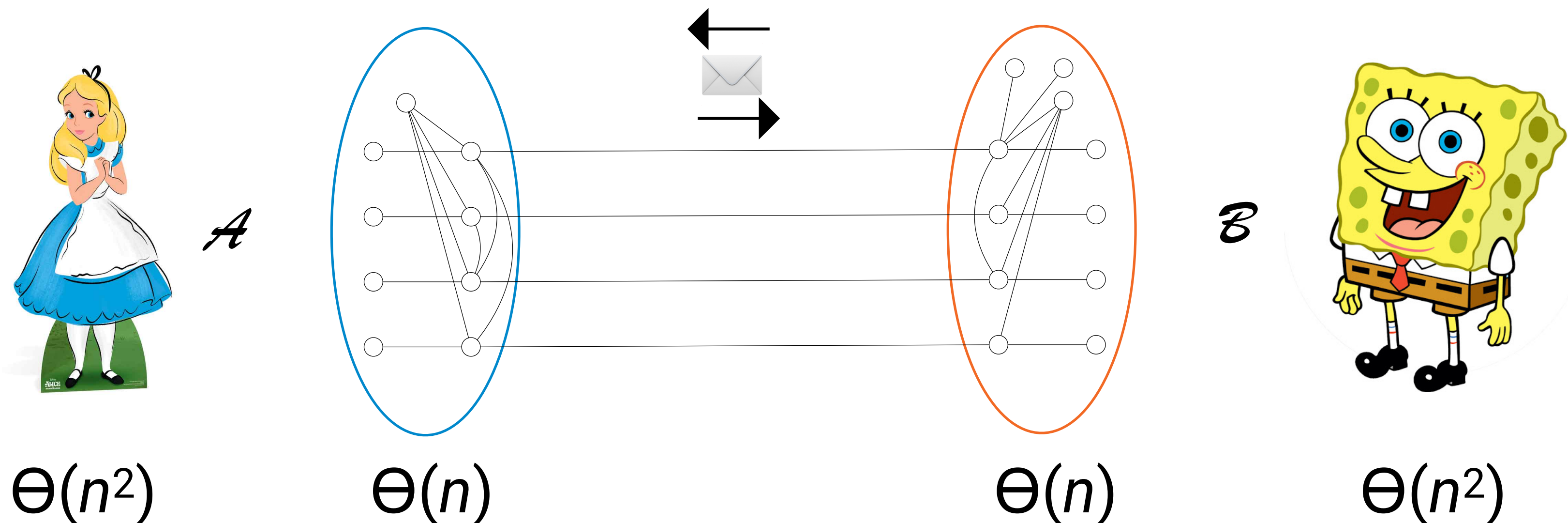
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Summary

- **LOCAL** model: unlimited bandwidth
- **CONGEST** model: $O(\log n)$ bandwidth
- $O(n)$ or $O(\text{diam}(G))$ time is no longer trivial
- Example:
 - **APSP** in time $O(n)$, pipelining helps
 - **APSP** requires $\Omega(n / \log n)$ rounds