# Theory of Distributed Systems <br> Exercise Sheet 8 

Due: Wednesday, 19th of June 2024, 12:00 noon

## Exercise 1: 2-coloring in paths

Show that there is no randomized distributed algorithm that finds a 2-coloring in paths in $o(n)$ rounds with probability at least 0.9 . Assume that $n$ is known and IDs are from $\{1, \ldots, n\}$.
Hint: use that the principle of locality can be extended to randomized algorithms i.e. for any possible output, nodes with the same view have the same probability of giving that output.

## Exercise 2: Independent sets in paths

An independent set (IS) is a subset of nodes such that no two neighboring nodes are in the independent set. A maximal independent set (MIS) is an independent set that cannot be extended. Assume that $n$ is known and IDs are from $\{1, \ldots, n\}$. Show that, in paths:

1. it is trivial to find some IS in $O(1)$ time with a deterministic distributed algorithm.
2. there exists an IS with at least $n / 2$ nodes.
3. it is not possible to find an IS of size at least $n / 2$ in $o(n)$ rounds.
4. there is no deterministic distributed algorithm that finds an MIS in $o\left(\log ^{*} n\right)$ rounds.

Bonus: a vertex cover of a graph is a subset of nodes that includes at least one endpoint of every edge of the graph. Deduce from number 3 above that it is not possible in paths to find a vertex cover of size at most $n / 2$ in $o(n)$ rounds.

## Exercise 3: Counting

Assume we are given a path of size $n$ where nodes know an upper bound on the size of the network in $\left\{n, \ldots, c n+c^{\prime}\right\}$ for some constants $c, c^{\prime}$ (i.e., nodes do not necessarily know the exact value of $n$ but only e.g. a constant approximation). Show that there is no deterministic distributed algorithm that counts the number of nodes in paths in $o(n)$ rounds, where at the end of the algorithm every node needs to output this count. Assume that IDs are from $\{1, \ldots, n\}$.

