

(7 Points)

## Algorithms and Datastructures Exercise Sheet 3

## Exercise 1: Bucket Sort

Bucketsort is an algorithm to stably sort an array A[0..n-1] of n elements where the sorting keys of the elements take values in  $\{0, \ldots, k\}$ . That is, we have a function key assigning a key  $\text{key}(x) \in \{0, \ldots, k\}$  to each  $x \in A$ .

The algorithm works as follows. First we construct an array B[0..k] consisting of (initially empty) FIFO queues. That is, for each  $i \in \{0, ..., k\}$ , B[i] is a FIFO queue. Then we iterate through A and for each  $j \in \{0, ..., n-1\}$  we attach A[j] to the queue B[key(A[j])] using the function enqueue.

Finally we empty all queues B[0], ..., B[k] using dequeue and write the returned values back to A, one after the other. After that, A is sorted with respect to key and elements  $x, y \in A$  with key(x) = key(y) are in the same order as before.

Implement *Bucketsort* based on this description<sup>1</sup>. You can use the template BucketSort.py which uses an implementation of FIFO queues that are available in Queue.py und ListElement.py.<sup>2</sup>

## Exercise 2: Radix Sort

## (13 Points)

Assume we want to sort an array A[0..n-1] of size *n* containing integer values from  $\{0,...,k\}$  for some  $k \in \mathbb{N}$ . We describe the algorithm *Radixsort* which uses *Bucketsort* as a subroutine.

Let  $m = \lfloor \log_b k \rfloor$ . We assume each key  $x \in A$  is given in base-*b* representation, i.e.,  $x = \sum_{i=0}^{m} c_i \cdot b^i$  for some  $c_i \in \{0, \ldots, b-1\}$ . First we sort the keys according to  $c_0$  using *Bucketsort*, afterwards we sort according to  $c_1$  and so on.<sup>3</sup>

- (a) Implement *Radixsort* based on this description. You may assume b = 10, i.e., your algorithm should work for arrays containing numbers in base-10 representation. Use *Bucketsort* as a subroutine. If you did not solve task 1, you may use a library function (e.g., **sorted**) as alternative to *Bucketsort*. (7 Points)
- (b) Compare the runtimes of *Bucketsort* and *Radixsort*. For both algorithms and each  $k \in \{2 \cdot i \cdot 10^4 \mid i = 1, ..., 60\}$ , use an array of fixed size  $n = 10^4$  with randomly chosen keys from  $\{0, ..., k\}$  as input and plot the runtimes. Shortly discuss your results in experiences.txt. (3 Points)
- (c) Explain the asymptotic runtime of your implementations of *Bucketsort* und *Radixsort* depending on n and k. (3 Points)

<sup>&</sup>lt;sup>1</sup>Remember to make unit-tests and to add comments to your source code.

 $<sup>^2 {\</sup>rm You}$  are allowed to use librarys, but note that the names of the methods may differ.

<sup>&</sup>lt;sup>3</sup>The *i*-th digit  $c_i$  of a number  $x \in \mathbb{N}$  in base-*b* representation (i.e.,  $x = c_0 \cdot b^0 + c_1 \cdot b^1 + c_2 \cdot b^2 + \ldots$ ), can be obtained via the formula  $c_i = (x \mod b^{i+1}) \operatorname{div} b^i$ , where mod is the modulo operation and  $\operatorname{div}$  the integer division.