



Theoretical Computer Science - Bridging Course

Exercise Sheet 3

Due: Tuesday, 12th of May 2026, 12:00 pm

Exercise 1: REs

(8 Points)

- (a) Let $\Sigma = \{a, b\}$. Let L_1 be the language defined by the regular expression $a^*b^*a^*$ and L_2 the language defined by a^*b^*b . Draw a DFA and NFA for L_1 and L_2 . (2 Points)
- (b) Let $\Sigma = \{a, b, c\}$. What language does the following regular expression describe $((a \cup c)^*b(a \cup c)^*b(a \cup c)^*b(a \cup c)^*)^*$? (2 Points)
- (c) Let $\Sigma = \{a, b\}$. Provide a regular expression that recognizes the following two languages.
- Let language L_3 contain all strings in which at least one of the symbols a or b occurs an even number of times. (2 Points)
 - Let language L_4 contain all strings of length at least 2 such that a and b are alternating. (2 Points)

Exercise 2: Limits of the Pumping Lemma

(4 Points)

Consider the language $L = \{c^m a^n b^n \mid m, n \geq 0\} \cup \{a, b\}^*$ over the alphabet set $\Sigma = \{a, b, c\}$.

- (a) Describe in words (not using the pumping lemma), why L can not be a regular language. (1 Point)
- (b) Show that, while the property described in the Pumping Lemma is a necessary condition for regularity, it is *not* sufficient for regularity.

Hint: Use L as counter example, i.e., show that it can be 'pumped' (in the sense of the pumping lemma), but is still not regular. (3 Points)

Exercise 3: Proving Non-regularity

(5 Points)

Use the Pumping Lemma to show that the following languages over the alphabet set $\Sigma = \{a, b, c\}$ are not regular.

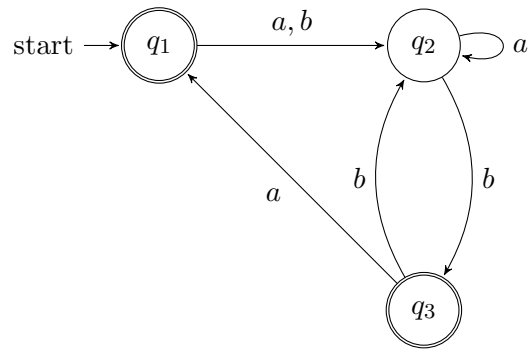
- (a) $L := \{a^n c b^{n+2} \mid n \geq 0\}$. (2 Points)
- (b) $L = \{a^m \mid m = n^2 \text{ for some } n \geq 0\}$. (3 Points)

Bonus: $L = \{a^n b w a^n \mid n \geq 1 \text{ and } w \in \Sigma^\}$.*

Exercise 4: NFA-GNFA-RE

(3 Points)

Consider the following NFA:



Give the regular expression defining the language recognized by this NFA by converting it *stepwise* into an equivalent GNFA with only two nodes. Document your steps.