



Algorithms and Datastructures

Winter Term 2022

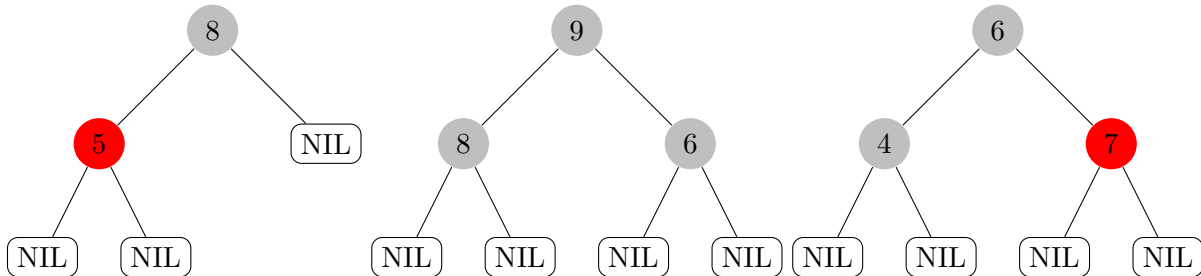
Sample Solution Exercise Sheet 6

Due: Wednesday, December 7th, 2pm

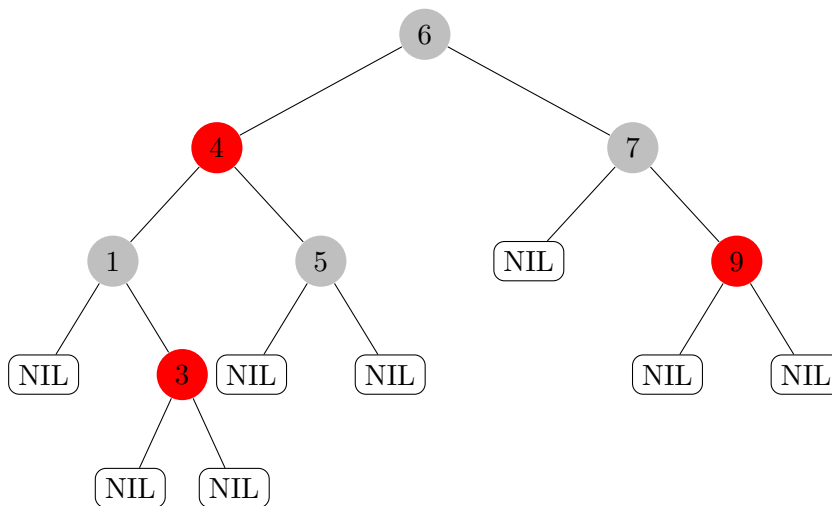
Exercise 1: Red-Black Trees

(10 Points)

(a) Decide for each of the following trees if it is a red-black tree and if not, which property is violated:



(b) On the following red-black tree, first execute the operation `insert(8)` and afterwards `delete(5)`. Draw the resulting tree and document intermediate steps.

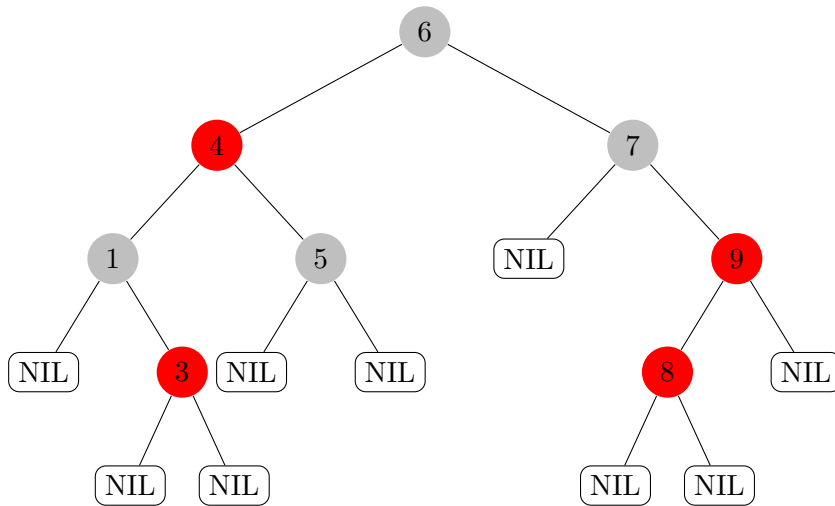


Sample Solution

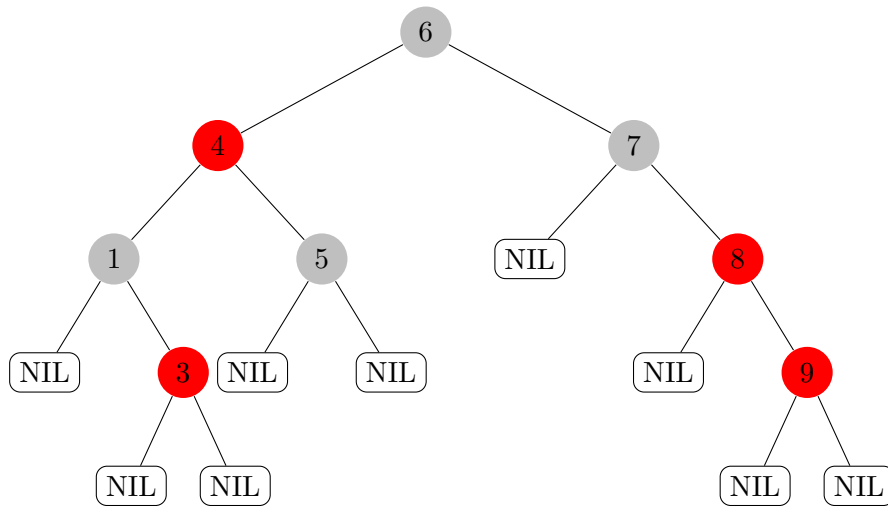
(a) From left to right:

- 1) Red-black-tree
- 2) No red-black-tree, because it is no binary search tree (the root's right child has a smaller key).
- 3) No red-black-tree, because the number of black nodes on a path from the root to a leaf is larger if you go through the left subtree.

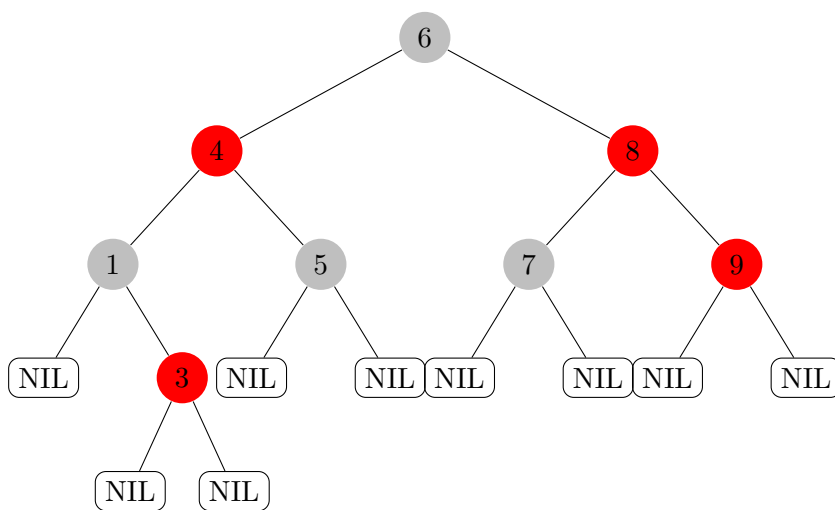
(b) We insert a red node with key 8 according to the rule of inserting into binary search trees.



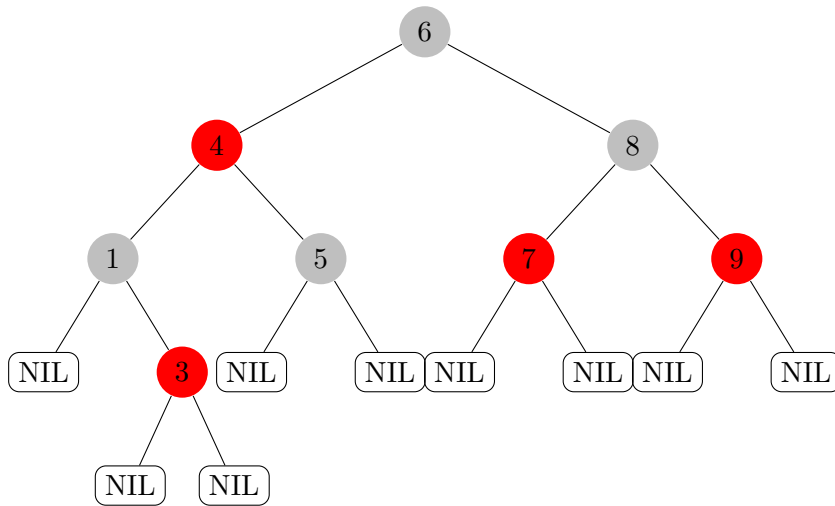
We are in case 1b from the lecture. We do a right-rotate(9,8),



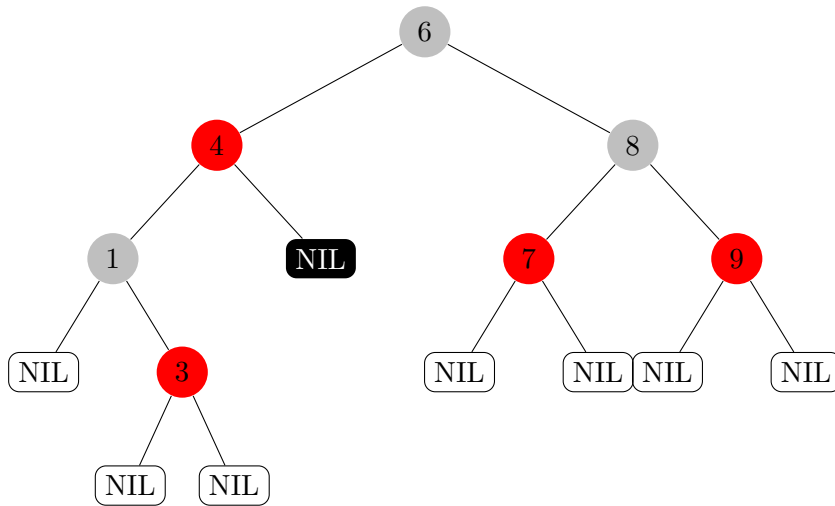
a left-rotate(7,8)



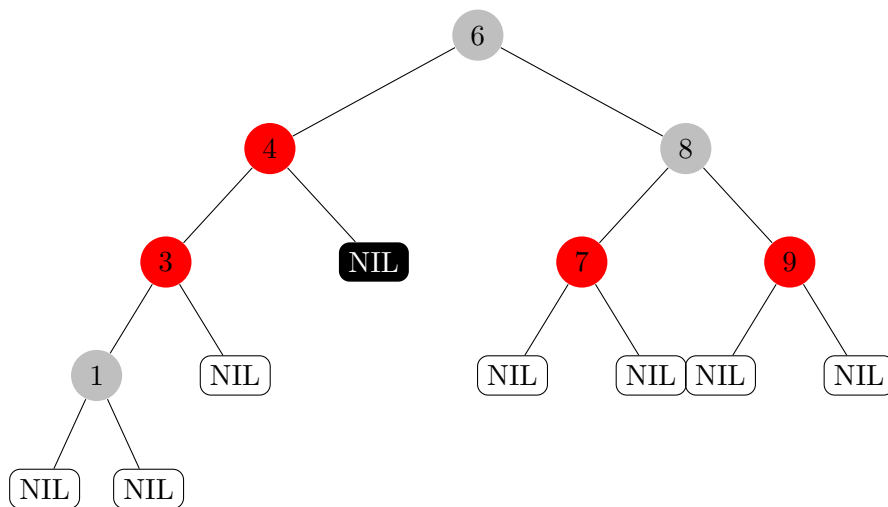
and recolor nodes 7 and 8.



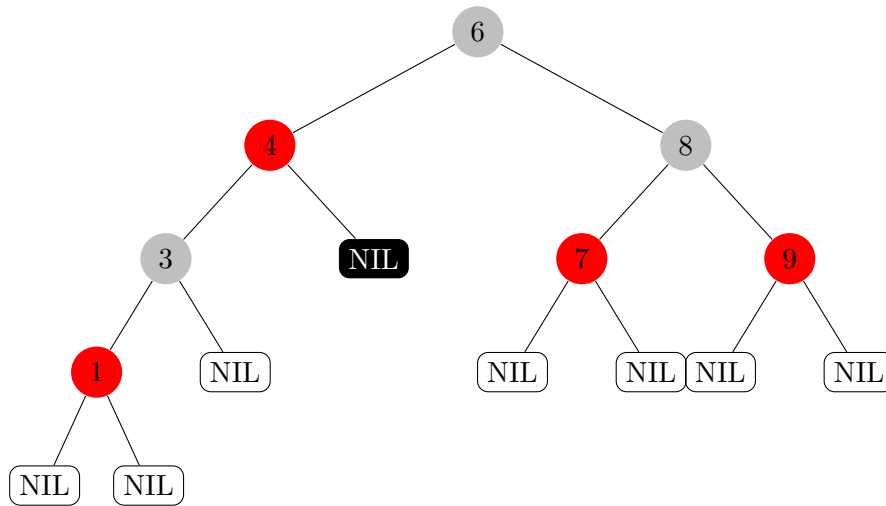
Now we execute `delete(5)`. We are in case 2b from the lecture (deleting a black node with two NIL-children). First we remove node 5 from the tree and color the right NIL-child of node 4 double black to correct the black height.



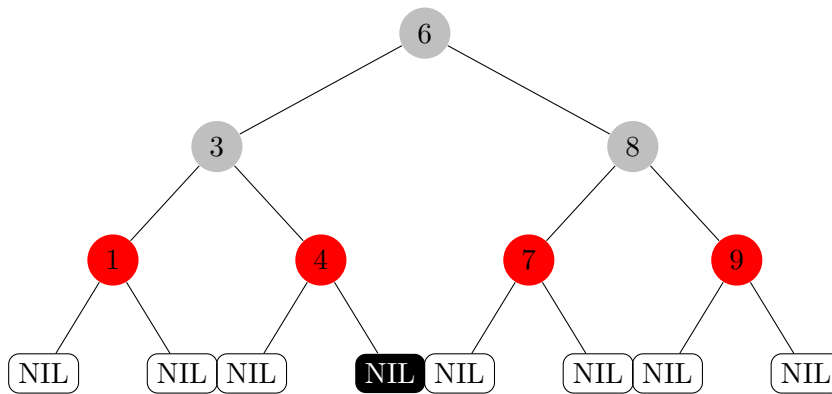
We are in case A.2 from the lecture. We do a `left-rotate(1,3)`



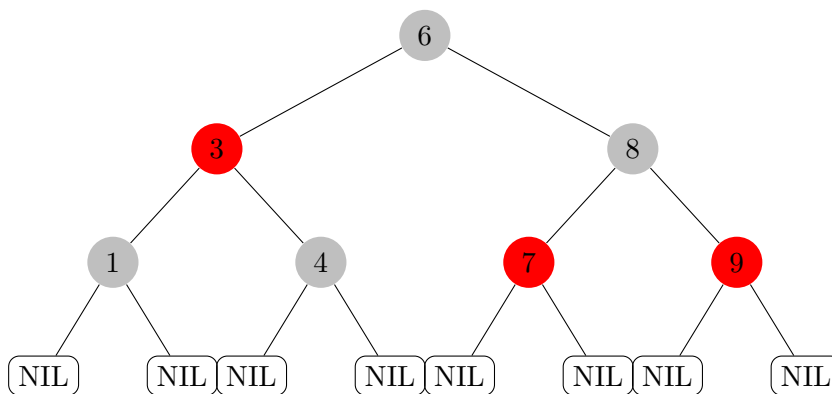
and recolor nodes 1 and 3.



Now we are in case A.1. We do a right-rotate(4,3)



and recolor. Finally, the tree looks like this.



Exercise 2: AVL-Trees ¹

(10 Points)

An AVL-tree is a binary search tree with the additional property that for each node v , the depth of its left and its right subtree differ by at most 1.

- (a) Show via induction that an AVL-tree of depth d is filled completely up to depth $\lfloor \frac{d}{2} \rfloor$. (3 Points)
A binary tree is filled completely up to depth d' if it contains for all $x \leq d'$ exactly 2^x nodes of depth x .

¹AVL-Trees are not part of the lecture. To solve this exercise the definition given below is sufficient.

- (b) Give a recursion relation that describes the minimum number of nodes of an AVL-tree as a function of d . (3 Points)
- (c) Show that an AVL-tree with n nodes has depth $\mathcal{O}(\log n)$. (4 Points)
You can either use part (a) or part (b).

Sample Solution

- (a) **Induktion start:** Each non-empty tree has a root and is hence completely filled up to depth 0. Hence the statement is true for $d = 0$ and $d = 1$ (as $\lfloor d/2 \rfloor = 0$ for $d = 0$ and $d = 1$).

Induktion step: Assume the statement holds for all AVL-trees up to depth d . We show that it also holds for AVL-trees of depth $d + 1$.

Let T be an AVL-tree of depth $d + 1$ with r as root and T_ℓ and T_r as left and right subtree. One of these subtrees must have depth d (lets say T_ℓ). As T is an AVL-tree, it follows that T_r has depth at least $d - 1$. By the induction hypothesis, T_ℓ is completely filled up to depth $\lfloor d/2 \rfloor$ and T_r is completely filled up to depth $\lfloor \frac{d-1}{2} \rfloor$. So both subtrees are completely filled up to depth $\lfloor \frac{d-1}{2} \rfloor = \lfloor \frac{d+1}{2} - 1 \rfloor = \lfloor \frac{d+1}{2} \rfloor - 1$ and hence T is filled completely up to depth $\lfloor \frac{d+1}{2} \rfloor$.

- (b) Let n_d be the minimum number of nodes in an AVL-tree of depth d . As every tree of depth d has at least $d + 1$ nodes (as it contains a path of length d), we obtain as base cases $n_0 = 1$ and $n_1 = 2$. Now let $d \geq 2$. An AVL-tree T of depth d consists of a root r , a left subtree T_ℓ and a right subtree T_r . One of them, lets say T_ℓ , has depth $d - 1$ and hence at least n_{d-1} nodes. As T is an AVL-tree, it follows that T_r has depth at least $d - 2$ and hence at least n_{d-2} nodes. Hence T has at least $n_d = n_{d-1} + n_{d-2} + 1$ nodes.

- (c) **Using (a):** And AVL-tree of depth d is filled completely up to depth $\lfloor \frac{d}{2} \rfloor$, so T has $n \geq 2^{\lfloor \frac{d}{2} \rfloor}$ nodes. We obtain

$$\begin{aligned}
 2^{\lfloor \frac{d}{2} \rfloor} &\leq n \\
 \iff \lfloor \frac{d}{2} \rfloor &\leq \log(n) \\
 \implies \frac{d}{2} - \frac{1}{2} &\leq \lfloor \frac{d}{2} \rfloor \leq \log(n) \\
 \implies d &\leq 2 \log n + 1 \\
 \implies d &\in \mathcal{O}(\log(n)).
 \end{aligned}$$

Using (b): Similar to the Fibonacci-series we have $n_d = n_{d-1} + n_{d-2} + 1 = 2n_{d-2} + n_{d-3} + 2 \geq 2n_{d-2}$. This means that increasing the depth by 2 doubles the number of nodes, so the number of nodes grows exponentially in the depth, or the depth grows logarithmically in the number of nodes. More formally, we have $n_d \geq 2n_{d-2} \geq 2^2 n_{d-4} \geq \dots \geq 2^{\lfloor d/2 \rfloor} n_{d-2\lfloor d/2 \rfloor} \geq 2^{\lfloor d/2 \rfloor} n_0 = 2^{\lfloor d/2 \rfloor}$. The rest follows as above.